

ALGEBRAIC DIFFERENTIATION USING
A TOP-DOWN SYNTAX ORIENTED TRANSLATOR

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ABSTRACT

This work employs a top-down syntax oriented translator to obtain a derivative of an algebraic expression. Tables of syntax and semantics to achieve this purpose are given. Simplification of the resulting expression is performed eliminating both redundant parentheses and terms consisting of 0 or 1. With a slightly extended version of the translator, partial derivatives can also be obtained, and any resulting expressions evaluated at assigned points.

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TABLE OF CONTENTS

	PAGE
ABSTRACT	ii
ACKNOWLEDGEMENTS	iii
TABLE OF CONTENTS	iv
CHAPTER	
I. INTRODUCTION	1
1.1. Introduction	1
1.2. Formal grammars	2
1.3. Works on algebraic differen- tiation and related fields	5
II. METHOD OF ANALYSIS	7
2.1. Introduction	7
2.2. Notation for syntax specifi- cation	8
2.2.1. Modification to standard BNF notation .	9
2.3. Entering the syntax table	13
2.4. The SEARCH routine and its analysis record	16
2.5. Interpretation of the analysis record	19
2.5.1. Introduction	19
2.5.2. Internal semantics	20
2.5.3. External semantics	23

	PAGE
2.6. Entering the external semantics	26
2.7. The COMPILE routine for external semantics	29
2.8. User's documents and other necessary parameters	30
III. ALGEBRAIC DIFFERENTIATION	35
3.1. Syntax and semantics for differentiation	35
3.2. Simplification of the form of the derivative	51
3.3. Evaluating the derivative at assigned points	60
IV. MODIFICATIONS MADE ON THE TRANSLATOR.	73
4.1. Introduction	73
4.2. Additional parameters for evaluation and order of differentiation	73
4.3. Internal defined syntax for dependent, independent variables and constants	73
4.4. Input routine for numeric data	77
4.5. Modification on the main program for co-ordination of the various subroutines.....	80

	PAGE
V. CONCLUSION	84
5.1. Introduction	84
5.2. Advantages of this system	84
5.3. Disadvantages	85
5.4. Suggestions for future works .	86
APPENDIX A. LISTING FOR DIFFERENTIATION OF RESTRICTED ARITHMETIC EXPRESSION	92
APPENDIX B. A COMPLETE LISTING OF THE TRANSLATOR	97
APPENDIX C. PROBLEM ON PRECEDENCE OF / AND *	118
APPENDIX D. COMPLETE OUTPUT LISTING FOR DIFFERENTIATIONS OF A EXPRES- SION.....	121
APPENDIX E. COMPLETE SYNTAX AND SEMANTICS LISTING FOR DIFFERENTIATION .	143
APPENDIX F. LISTINGS OF SIMPLIFIED EXPRES- SION	146
REFERENCES	150

CHAPTER I

INTRODUCTION AND REVIEW OF LITERATURE

1.1. Introduction.

This work employed a syntax-directed translator [30] to perform algebraic differentiation on a digital computer. Particular efforts are made on the provision of the necessary syntax specifications so that the input algebraic expressions would be recognized. With the necessary semantics the corresponding target string would be produced. This can be in any desired form depending on the interpretation rules (semantics) ; in this case it would be the result of algebraic differentiation.

The translator employed uses a top-down technique [9, 16, 30] as contrast to the bottom-up method used in other syntax-directed translator [26, 27, 31]. Some basic features of the translator are discussed and the precise rules of syntax and semantics are discussed next. This translator was originally written with a purpose of performing works on compiling but it was found that other applications can be done as well [20]. This work is one of the example of these applications.

The advantage of this method is that any changes and modifications necessary in the translation can be made by just changing the syntactic and semantic

specifications without regard to the logic of the translator. This would not be the case if a program were written specifically for analytic differentiation [23, 29, 35].

In contrast to Schorr's [39] work, three syntax and semantic tables are provided for translation. The first set deals with algebraic differentiation using the basic rules of differentiation whereas the target string produced is processed alternately by the other two sets of syntax and semantic tables to obtain a simplified output. With a separately written input routine and 'value' routine for the translator the following can be obtained:

1. the derivatives of higher order;
2. the partial derivatives;
3. the derivatives of implicit functions; and
4. the evaluation of the resulting derivative.

1.2. Formal grammars.

The importance of formal grammars was first apparent when in the report on ALGOL [33, 34] the language is defined by a set of syntactic rules. The studies were further spurred when the development of the syntax-directed and syntax-controlled processors required a formal definition of the language to be processed. The first important original papers on formal studies of

grammars were done by Chomsky [10, 11] and Bar-Hillel et al [2]. Since then formal studies on grammars have become one of the major fields in linguistic studies both in programming and natural languages. It is found that out of the formal grammars the context-free phrase-structure grammar is of particular interest. The definition of this grammar is perhaps best expressed with the notation used by Ginsburg [21]. Suppose \mathcal{A} denotes the set of terminal symbols (e.g. $\underline{T}_1, \underline{T}_2, \underline{T}_3$, etc.), η , the set of nonterminal symbols (denoted by Latin capitals e.g. $\underline{U}, \underline{V}$, and \underline{Z}) and the vocabulary \mathcal{V} is defined as the union of \mathcal{A} and η . The symbols $\underline{S}, \underline{S}_1, \underline{S}_2$ etc. are used to denote members of \mathcal{V} , while strings of symbols (including the empty string Λ) are denoted by lower case Latin letters $\underline{u}, \underline{v}, \underline{w}, \dots$. The context-free phrase-structure grammar \mathcal{G} then can be defined as a finite set of productions of the form $\underline{U}_i \rightarrow \underline{u}_i$ with the following properties :

1. each \underline{u}_i is a nonempty string whose symbols are in \mathcal{V}
2. each \underline{U}_i is a nonterminal symbol i.e. $\underline{U}_i \in \eta$
3. there is a unique \underline{U}_i which occurs in no \underline{u}_i .

A language that can be specified in terms of this grammar is called a context-free phrase-structure language. As an example the following is a context-free phrase-structure grammar

$$\begin{array}{lcl}
 S & \longrightarrow & E \\
 E & \longrightarrow & T + E \\
 E & \longrightarrow & T \\
 T & \longrightarrow & P * T \\
 T & \longrightarrow & P \\
 P & \longrightarrow & I \\
 P & \longrightarrow & (E)
 \end{array}$$

where S E T P are nonterminal symbols and + * () I are terminal symbols. This grammar specifies restricted arithmetic expressions like $I+I*I$ and $(I+I+I)*(I+I)$ etc. However, in order to express the grammar in a more compact form Backus Normal Form (BNF) [1] is used. In this notation the symbol ' \longrightarrow ' is replaced by ' $::=$ ' while production rules with the same left side are combined using the symbol ' $|$ ' with the meaning "or". Hence the above grammar can be rewritten as

$$\begin{array}{lcl}
 S & ::= & E \\
 E & ::= & T + E \mid T \\
 T & ::= & P * T \mid P \\
 P & ::= & I \mid (E)
 \end{array}$$

The BNF notation is in fact proved to be equivalent to the formal notation by Ginsburg [21] and is the notation employed for the syntax specification in this work. The particulars and details of these syntactic rules are discussed in Chapter II.

1.3. Works on algebraic differentiation and related fields.

There is much literature on the topic of formula manipulation using a digital computer [36, 37]. The most significant systems are perhaps the FORMAC [38] and ALPAK [6, 7, 25]. These systems are general systems that handle algebraic manipulations like multiplication, division, addition and subtraction of polynomials, differentiation and certain simplifications of algebraic expressions. Programs also have been written solely on algebraic differentiation. The earliest work on this was written independently by Nolan [35] and Kahramanian [29] in 1953. In both cases the mathematical expression to be differentiated has to be completely parenthesized and encoded into a three-address code. The derivative was output in tabular form and the resulting analytic expression for the derivative then developed by hand. A more recent work was done by Hanson et al [23] in which the algebraic expression can directly serve as input without any modifications, but is transformed automatically into a table of triples by Ershov's algorithm [14]. The required derivative of each triple can then be obtained by applying the basic rules of differentiation on these triples. The parenthesis-free tabular form of the derivative is then transformed into a parenthesized string or expression. Another approach to symbolic differentiation is that of Weissman [42] who makes use of the list process-

ing language LISP 1.5. The input algebraic expression is first transformed to its prefix form, then differentiated, simplified and converted back to its infix form for output.

All these approaches to algebraic differentiation are written in a specific language for symbol-manipulation. In this work an entirely different method is used. A top-down syntax oriented translator is employed with the proper syntax and semantic tables. Schorr [39] also used this method for algebraic differentiation, but the translator used is a bottom-up type developed by E. T. Irons [26, 27]. There has been much controversy about the relative efficiency of the top-down and bottom-up translators [4, 22], however, it is not intended here to present an analytic judgement on these two types.

CHAPTER II
METHOD OF ANALYSIS

2.1. Introduction.

The translator employed was first designed and written by King [30] in ALGOL 60. A pure top-down syntax analysis technique is used in this translator to analyze the input string. The input string is first read in together with its assumed phrase (syntactic) type, and is stored in an integer array called "source". It is then analysed according to the given syntax. If the search is successful, that is, if the source string conforms to the syntax, the corresponding target string is produced according to the semantics, and is stored in the integer array "targ" ready for output. The translator thus consists of four main subroutines :

- (a) READSYNTAX to read and store the syntax.
- (b) READSEMANTICS to read and store one or more sets of semantics.
- (c) SEARCH to attempt a syntactic analysis of the source string according to the syntax and to assemble a tabular analysis record if the search is successful.
- (d) COMPILE to apply the semantics to the source string using the analysis record, in order to produce the target string.

The program also contains a number of subsidiary routines for input output etc.; a complete listing of the program appear in Appendix B.

2.2. Notation for Syntax Specification.

The notation employed is basically Backus Normal Form (BNF). Backus Normal Form has become extremely well known since its first appearance in 1959 [1]; it is perhaps best to explain its principal features with examples.

Example 1

$$\langle \text{digit} \rangle ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

Here digit is the name of the phrase to be defined and is enclosed by the symbols \langle and \rangle to distinguish it from the terminal symbols 0, 1, 2, 9; the symbol $::=$ is a separation symbol separating the right and left side of the definition and it can be interpreted as "is defined as"; the symbol \mid separates the alternatives and has the meaning 'or'.

The above definition means that phrase type $\langle \text{digit} \rangle$ is defined as the symbol 0 or 1 or 2 or 3 or 9.

A phrase type may be defined in terms of other phrase types.

Example 2

$$\langle \text{integer} \rangle ::= \langle \text{digit} \rangle \mid \langle \text{integer} \rangle \langle \text{digit} \rangle$$

Note the recursive nature of the above definition; an integer is defined as one digit or a string of digits of arbitrary length.

A phrase type can be defined in terms of both terminal symbols and phrase types.

Example 3

$$\langle \text{u/s integer} \rangle ::= \langle \text{integer} \rangle \mid + \langle \text{integer} \rangle \mid - \langle \text{integer} \rangle$$
2.2.1. Modifications to standard BNF notation.

Because of the search mechanism employed in our translator the following two restrictions must be imposed on the phrase definitions.

1. Infinite recursion

A definition of the form

$$\begin{array}{l} \langle A \rangle ::= \langle A \rangle \langle B \rangle \\ \text{or } \langle A \rangle ::= \langle B \rangle \\ \langle B \rangle ::= \langle C \rangle \\ \langle C \rangle ::= \langle A \rangle \end{array}$$

will cause the search routine to search in an infinite loop. In order to avoid this the following restrictions are imposed on the grammar :

(i) No definition should start with the phrase type it is defining.

(ii) In a set of definitions ($S_1, S_2, S_3, \dots, S_n$), no cycle or loop should exist i.e. none of the alternatives of S_1 may begin with S_2 , S_2 with S_3, \dots, S_n with S_1 .

Thus in view of the above conditions

$$\langle \text{integer} \rangle ::= \langle \text{integer} \rangle \langle \text{digit} \rangle \mid \langle \text{digit} \rangle$$

is not a permissible definition. It may be rewritten as

$$\langle \text{integer} \rangle ::= \langle \text{digit} \rangle \langle \text{integer} \rangle \mid \langle \text{digit} \rangle$$

2. Non-maximum search.

Using the definition

$$\langle \text{integer} \rangle ::= \langle \text{digit} \rangle \mid \langle \text{digit} \rangle \langle \text{integer} \rangle$$

due to the method employed in the search routine, no integer of more than one digit will be recognized successfully. For example given the integer 456, the search routine will recognize the digit 4 and exit because the first alternative for phrase type $\langle \text{integer} \rangle$ is satisfied. Therefore it is necessary to rearrange all the alternatives so that they are in the order of decreasing length of phrase which may be found. A satisfactory definition of integer would be

$$\langle \text{integer} \rangle ::= \langle \text{digit} \rangle \langle \text{integer} \rangle \mid \langle \text{digit} \rangle$$

The above restrictions are in fact quite common for top-down syntax analyzers.

The following fairly trivial additional modifications are made for this particular translator :

1. The metasymbol $::=$ is replaced by the symbol $=$.

2. Each phrase type is assigned a unique positive integer.

3. The phrase type name referred on the right side of the metasymbol '=' can be replaced by its corresponding phrase type number.

4. Each phrase type definition is terminated by the metasymbol ';'.

With the above modifications the phrase definitions

$$\begin{aligned} \langle \text{digit} \rangle & ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 \\ \langle \text{integer} \rangle & ::= \langle \text{digit} \rangle \langle \text{integer} \rangle | \langle \text{digit} \rangle \end{aligned}$$

are rewritten as

$$\begin{aligned} 11 \langle \text{digit} \rangle & = 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 ; \\ 12 \langle \text{integer} \rangle & = \langle 11 \rangle \langle 12 \rangle | \langle 11 \rangle ; \end{aligned}$$

5. In order to increase the efficiency of searching the metasymbols '(' and ')', are introduced. These metasymbols may enclose alternatives that have common subphrases. <E> is used for the null string. Hence phrase type 12 above can be rewritten as

$$12 \langle \text{integer} \rangle = \langle 11 \rangle (\langle 12 \rangle | \langle E \rangle);$$

6. Some of the most common phrase definitions such as <digit>, <integer>, <number>, <letter>, <identifier> and <string> that are frequently referred to are built into the translator. The phrase type digit

may be referred to as $\langle D \rangle$, letter as $\langle L \rangle$, identifier as $\langle I \rangle$, integer as $\langle N \rangle$ and character string as $\langle S \rangle$. Hence if we want to define a phrase type $\langle \text{decimal number} \rangle$, with the above built-in phrase definitions we can define it as

$$31 \quad \langle \text{decimal number} \rangle = \langle N \rangle . \langle N \rangle | \langle N \rangle . | \langle N \rangle | . \langle N \rangle ;$$

or

$$31 \quad \langle \text{decimal number} \rangle = \langle N \rangle (. \langle N \rangle | . | \langle E \rangle) | . \langle N \rangle ;$$

In fact automatically assigned phrase type numbers are entered whenever the above built-in phrase type (bips) are used; 1 is assigned to phrase type letter $\langle L \rangle$, 2 to digit $\langle D \rangle$, 4 to identifier $\langle I \rangle$, 6 to integer $\langle N \rangle$, and 8 to character string $\langle S \rangle$. Actually, integers from 1 to 10 are reserved for bips.

7. The symbols $\langle \rangle$; () and = are reserved as metasymbols in phrase type definitions and therefore cannot be used as terminal characters. They are enclosed by cornered brackets if required as terminals. Thus $\langle \langle \rangle \rangle$ represents the terminal character \langle , $\langle \rangle \rangle$ the terminal character \rangle , $\langle | \rangle$ the terminal character $|$ and so on.

Example of a complete syntax table.

$$11 \quad \langle \text{primary} \rangle = \langle 12 \rangle | \langle 13 \rangle | \langle 14 \rangle | \langle \langle \rangle \rangle | \langle 16 \rangle | \langle \rangle \rangle ;$$

$$12 \quad \langle \text{indept. variable} \rangle = X;$$

$$13 \quad \langle \text{dependent variable} \rangle = Y;$$

$$14 \quad \langle \text{constant} \rangle = 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9;$$

$$15 \quad \langle \text{term} \rangle = \langle 11 \rangle (* \langle 15 \rangle | \langle E \rangle) ;$$

$$16 \quad \langle \text{expression} \rangle = \langle 15 \rangle (+ \langle 16 \rangle | \langle E \rangle) ;$$

Some valid phrases of type 16 would be

$4 + X * Y$
 $2 * X * (Y + 5)$
 $X + Y$
 $4 + 5 * X$
 X
 0
 $(3 + 9 + X)$

2.3. Entering the syntax table.

The syntax table must be first read and converted into an internal form that can be made use of by the SEARCH routine in the syntax analysis. This is done by the routine READSYNTAX with two integer arrays DSYN and ILFSYN used for storage purposes. The integer array DSYN is used to store the contents on the right of the metasymbol '=' of each phrase type while another integer array ILFSYN is used to store the start location of each phrase type in DSYN. Hence each element of the vector ILFSYN $[i]$ is, in fact, a pointer to the start of phrase type number i in the array DSYN. Note that the intervening material between the phrase type number and the metasymbol '=' is ignored by the READSYNTAX routine. The significant symbols following '=' are entered successively into the array DSYN with the following assumptions.

1. The metasymbols '=', '<' and '>' are not entered into the array DSYN but used as check points to evoke some appropriate instructions.

2. A unique positive input code number is assigned to each different terminal symbol read.

3. The phrase type number enclosed by the metasymbols '<' and '>' is entered into DSYN as -1 followed by the phrase type number. Hence <11> , would be entered as -1 11 in the array DSYN. Negative numbers are also assigned to some of the metasymbols. The input code and internal representation of each metasymbol is shown in the following list.

<u>metasymbol</u>	<u>Corresponding identifier used in readsyntax</u>	<u>input code</u>	<u>internal representation</u>
	<u>for input code</u>	<u>for internal representation</u>	
=	equals	nil	49 nil
<	ob	nil	44 nil
>	cb	nil	45 nil
(inin	in	46 -10
)	outin	out	47 -12 -11
;	termin	term	50 -12 -13
	orin	or	51 -12

4. The built-in phrase types would be entered by the corresponding built-in phrase type number with a preceding negative one. Hence

<L>	would be entered as	-1	1
<D>	as	-1	2
<I>	as	-1	4
<N>	as	-1	6
<S>	as	-1	8

For <E> nothing is entered since it represents a null string.

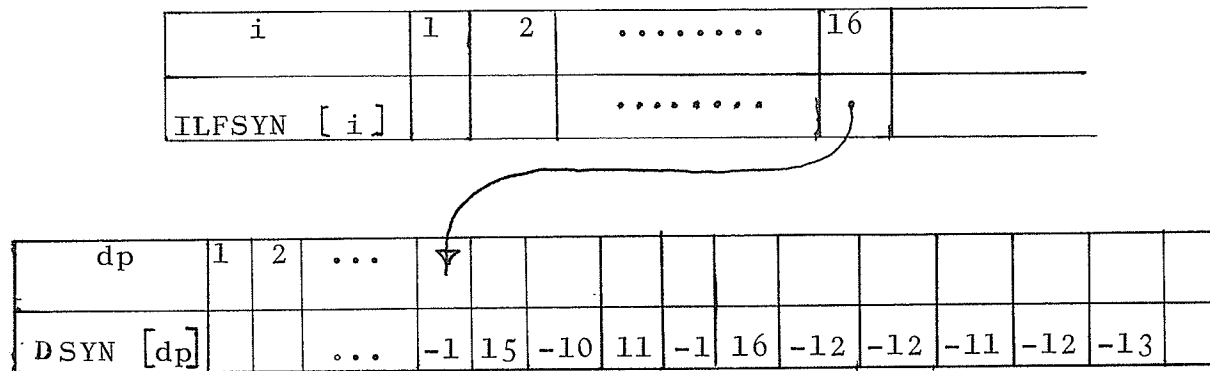
5. When the metasymbol ; is read and the corresponding internal representation is entered, start to read in the next phrase type number as above and perform a similar process. This will be repeated until a suitable terminator for the syntax table is reached such as the occurrence of a phrase of type 0.

Example 5

The locations of ILFSYN and D SYN that contains the phrase type

$$16 \text{ <expression> } = \text{ <15> (+ <16> | <E>)}$$

is shown schematically below.



where ILFSYN [16] points to the start location in the array DSYN for the phrase type considered. Note that the positive number 11 is the internal code for + and the terminator ; of the phrase type is entered as -12 -13.

2.4. The SEARCH routine and its analysis record.

When the syntax table is stored as above, a source string, given as of a certain phrase type can then be analysed by the SEARCH routine. The first category of the assumed phrase type is searched recursively until a definite answer, either a 'success' or 'failure', is returned. When the result of this search is failure, the next category is tried. The process is repeated until either a successful search is made in one of the alternatives of the assumed phrase type or the terminator of the phrase definition is reached. The details of the development and operation of the procedure SEARCH are fully described in [30]. The

important feature here is the information assembled by the routine in the analysis record. This comprises, for positive integer $j = 1, 2, \dots$, the following integer arrays.

- type $[j]$: the phrase type of entry no. j in the record.
- cat $[j]$: the category recognized in this phrase type.
- from $[j]$: the starting position of the source string analyzed.
- to $[j]$: the ending position of the source string analyzed.
- link $[i,j]$: the position in the record where details of the constituents are stored.

Example 6

The analysis record for the source string
 4 + X * Y according to the syntax table on page 12
 would be

entry	j	<u>type</u> [j]	<u>cat</u> [j]	<u>from</u> [j]	<u>to</u> [j]	<u>link</u> [i,j] i = 1	i = 2
1		16	1	1	5	2	5
2		15	2	1	1	3	0
3		11	3	1	1	4	0
4		14	5	1	1	0	0
5		16	2	3	5	6	0
6		15	1	3	5	7	9
7		11	1	3	3	8	0
8		12	1	1	1	0	0
9		15	2	5	5	10	0
10		11	2	5	5	11	0
11		13	1	5	5	0	0

The first entry of the above analysis record then reads that it is a phrase type 16 of category 1, occupying positions from 1 to 5 of the source string, and the records of its constituent parts are located in entry 2 (for type 15) and 5 (for type 16) respectively. The second entry simply shows that the phrase type 15 that was found is of category 2, occupying only one

position of the source string and further details are located in entry 3. When there is no more detail that can be traced, the value of link $[i,j]$ is automatically zero. The integer i indicates the number of links required and is dependent on the number of phrase types that appears as constituent parts of a definition.

Thus the definitions

$$\begin{array}{lcl} 1 & \langle A \rangle & = \quad \langle 2 \rangle \quad \langle 3 \rangle \quad \langle 3 \rangle ; \\ 2 & \langle B \rangle & = \quad \langle 3 \rangle \quad \langle 2 \rangle \quad ; \\ 3 & \langle C \rangle & = \quad \langle 2 \rangle \quad ; \end{array}$$

requires respectively 3, 2 and 1 links. Usually we find that 3 is sufficient.

2.5. Interpretation of the analysis record : Compilation.

2.5.1. Introduction.

In some compiler or translator systems the syntax and semantic analysis is so inextricably entwined that it is impossible to treat them separately and hazardous to change. However this translator has quite separate routines to treat the syntax and semantics. The advantages of this is that either the syntax or semantics or both can be modified without too much difficulty.

There are two possible and useful approaches for providing semantics. The first approach is internal : the semantic rules appear as subroutine in the translator and are written in the language of the translator. The second approach is to design a special purpose metasemantic language in terms of which the semantics can be expressed and which like the syntax form part of the primary data.

2.5.2. Internal semantics.

Here semantic interpretation is performed using a recursive ALGOL procedure (which we called COMPILE for the moment) which has the following procedure heading.

```
procedure compile (i,j); value i,j; integer i,j;
```

The procedure interprets the ith link of the phrase detailed by the jth entry of the analysis record. Details of this jth entry can be obtained by the following instructions.

```
j:=link [i,j] ; i:=type [j] ; k:=cat [j] ;
```

The variables i and k specify the type and category of the phrase to which the original ith link of the jth entry referred. The procedure has a set of instructions for interpreting every possible type and

category of a phrase, including recursive calls when interpretation of the constituent phrases is required. This procedure also make use of three other additional procedures for the semantic interpretation, namely :

writetext ('string') : to insert the string quoted to the target string.

copyphrase (j) : to copy the phrase given by the j entry of the analysis record to the target string, that is, that portion of source phrase specified by the pointers from [j] and to [j] of the jth entry of the analysis record.

copylink (i,j) : to copy the ith constituent of the phrase defined by the jth entry of the analysis record. This is equivalent to copyphrase (link [i,j]).

With the given syntax and analysis record the algebraic differentiation of the given restricted arithmetic expression $4 + X * Y$, can be performed by the procedure.

```

procedure diff (i,j) value i,j; integer i,j;
begin      integer k;
           switch s := primary, ind variable,
                    dep variable, constant,
                    term, expression ;
j := link [i,j] ; i := type [j] ; k := cat [j] ;

goto      s [i - 10] ;
primary :   if k = 4 then
           begin writetext ('('); diff (1,j); write-
                    text (')')
                    end

           else diff (1,j) ;
           goto end ;

ind variable : writetext ('1'); goto end ;
dep variable : copyphrase (j) ; writetext ('') ; goto end ;
constant :    writetext ('0'); goto end ;
term      :   if k = 1 then
           begin writetext ('('); diff (1,j);
                    writetext ('*'); copylink (2,j);
                    writetext ('+');
                    copylink (1,j); writetext ('*');
                    diff (2,j);      writetext (')');
           end

```

```

    else   diff (1,j);
    goto   end ;
expression : if  k = 1 then
    begin diff (1,j); writetext ('+'); diff (2,j);
    end
    else   diff (1,j);
end          : end diff ;

```

Hence, with information stored in the analysis record and with link [1,0] set to 1, the instruction

```
diff (1,0) ;
```

will cause the output of the derivative of the source string. In our particular example it would be

$$0 + (1 * Y + X * Y')$$

Therefore the output expression is dependent on the procedure which is written for the particular semantics. In this case it is the procedure DIFF .

2.5.3. External semantics.

The semantics described above has a serious drawback in that when the semantics have to be changed or modified for some reason particularly during testing, the entire procedure must be corrected, re-

written and recompiled many times. A metasemantic language has been designed for semantic rules each of which has a form similar to that of the syntax definition. Each rule begins with a phrase type number followed usually by the name of the phrase type enclosed by the cornered brackets then the symbol '=' followed by the semantics. The definition is terminated by the symbol ';'.

The semantics may include five basic instructions described as follows.

- (i) $W \langle \text{string} \rangle$: copy the string into the output. If any of the metasympols '=', '.', '<', '>' and '&' are required, they appear in the form '<;>', '<.>', '<◇>', '<◇◇>' and '&' respectively. When space or newline is required, the corresponding instruction would be $W \langle \langle S \rangle \rangle$ and $W \langle \langle N \rangle \rangle$ respectively.
- (ii) P : copy of the phrase of the current entry of the analysis record into the target string.

- (iii) $L \langle i \rangle$: copy the i th constituent of the phrase of the current entry of the record into the target string.
- (iv) $C \langle i \rangle$: compile, recursively, the i th constituent given by the current entry of the record.
- (v) $K \langle i_1, \dots, i_n \rangle$: 'K' specifies the categories i_1, \dots, i_n ($n > 0$). If the current category belongs to the set $\{i_1, \dots, i_n\}$ then the instructions following K are obeyed until either one of the two terminators $.$ (meaning end) or $\&$ (meaning else) is reached; otherwise these instructions are skipped. The terminator $\&$ is used only when followed by another $K \langle i'_1, \dots, i'_{n_1} \rangle$ where the sets $\{i_1, \dots, i_n\}$, $\{i'_1, \dots, i'_{n_1}\}$ are disjoint. Otherwise the terminator $.$ should be used instead.

Hence instead of the procedure DIFF (page 22) written for the interpretation of the syntax table of page 12 , the following corresponding external semantics for algebraic differentiation of the restricted arithmetic

expressions (e.g. $4 + X * Y$) can be provided instead.

- 11 <primary> = K <4> W<(> .
 K <1,2,3,4> C <1> .
 K <4> W <)> . ;
- 12 <indep variable> = W <1> ;
- 13 <dependent variable> = PW <'> ;
- 14 <constant> = W <0> ;
- 15 <term> = K <1> W <(> C <1>
 W <*> L <2> W <+>
 L <1> W <*> C <2>
 W <)> & K <2> C <1> . ;
- 16 <expression> = C <1> K <1> W <+>
 C <2> . ;

The entire input listing of the syntax and semantics and the analyzed output of $4 + X * Y$ can be found in Appendix A.

2.6. Entering the external semantics.

The semantic table is read in using the routine READSEMANTICS. Like READSYNTAX two integer arrays ILFSEM and DSEM are assembled and have similar functions to ILFSYN and DSYN in READSYNTAX. After the phrase type number is read, the intervening character

string is again skipped until the metasymbol '=' is reached. The symbols then entered by the READSEMANTICS routine with the following assumptions.

1. The metasymbols '=', '<' '>' and ',' are again not entered into the array DSEM but serves as check points to evoke some appropriate instructions.
2. A positive input code number is assigned to each different symbol entered.
3. Any character string between any one of the special characters (W, K, C, L) and the metasymbol '<' is ignored and treated as comment. Special characters (W, K, C, L) and the metasymbol '<' is ignored and treated as comment.
4. To facilitate the work of the compile routine later, the input codes of some metasymbols are changed to negative code numbers whenever necessary. Sometimes two internal code numbers are necessary for the representation of a metasymbol. For example K is represented by two internal codes, "kst" (-5) and "kterm" (-9). "kst" signifies the start of the categories entered following by the categories and terminated by kterm. Hence for K <1, 2, 3> , it will be entered internally as

-5	1	2	3	-9
----	---	---	---	----

Note that metasympols '<', ',', and '>' are not entered. The input code and internal representation of each metasympol is shown below.

<u>metasympol</u>	<u>Corresponding identifier used in readsemantics</u>		<u>input code</u>	<u>internal represent- ation</u>
	<u>for input code</u>	<u>for internal representation</u>		
=	equals	nil	49	nil
<	ob	nil	44	nil
>	cb	nil	45	nil
,	comma	nil	42	nil
.	endin	endd	43	-3
;	termin	term	50	-13
&	elsein	else	52	-2
*N	nn	nl	28	-1
*S	ss	sp	33	0
C	cc	cc	17	17
L	ll	ll	26	26
P	pp	pp	30	30
W	ww	ww, wterm	37	37, -4
K	kk	kst, kterm	25	-5, -9

* The metasymbol N and S must be always of the form $\langle N \rangle$ and $\langle S \rangle$ and when combined with the metasymbol W would give the desired function.

$W \langle \langle N \rangle \rangle$ would be interpreted by the compile routine as to skip a line, and $W \langle \langle S \rangle \rangle$ would mean to skip a space, whereas $W \langle N \rangle$ and $W \langle S \rangle$ would mean print out character N and S respectively.

5. Again, each semantic rule is terminated by the metasymbol ';'. The next semantic rule is entered similarly as above until the terminator '0' for the semantic table is reached.

Also two Boolean variables kf and elf are set so that any violations of the rules as specified on page 25, can be checked and appropriate action taken.

2.7. The COMPILE routine for External Semantics.

The COMPILE routine has the same initial statements as on page 20. It is a recursive procedure which will make use of the analysis record as well as the input semantic table to perform the desired interpretation. Appropriate instructions corresponding to the metasymbols W, P, L, C and K as described on page 24 are executed in this routine and the resulting target string is stored in the array TARG ready for output. The

the maximum length of the target string. The numeric values in card 3 can in fact be changed whenever necessary.

Immediately following the initial data are the arbitrary number of user's documents. A typical user's document is shown below.

+++

DIFFERENTIATION OF ALGEBRAIC EXPRESSION ;

16, 80, 160, 1 ;

11 <primary> = <12> | <13> | <14> | <() <16> <1> ;

12 <indept variable> = X;

13 <dependent variable> = Y;

14 <constant> = 0 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 ;

15 <term> = <11> (* <15> | <E>) ;

16 <expression> = <15> (+ <16> | <E>) ;

0;

1;

11 <primary> = K <4> W <() .
 K <1, 2, 3, 4> C <1> .
 K <4> W <1> .;

12 <indept variable> = W <1> ;

13 <dependent variable> = PW <'> ;

14 <constant> = W <0> ;

15 <term> = K <1> W <() C <1> W <*> L <2>
 W <+> L <1> W <*> C <2> W <1> &
 K <2> C <1> . ;

```

16 <expression>          =  C <1> K <1> W <+> C <2> . ;
0;
0;
16;
4 + X * Y ;
:
followed by any arbitrary number of input strings
preceded with an assumed phrase type.
:
0;
+++

```

The first card begins with three consecutive '+' symbols to signify the beginning of a user's document. Following this card is any comment (excluding the character ;) identifying this document terminated with a semi-colon. The third card consists of parameters which provide informations about the maximum phrase type number required in the syntax table, the maximum size of the integer array DSYN, the maximum size of the integer array DSEM, and the number of the semantic sets required. All these parameters are separated by a comma or semi-colon. The numeric values of these parameters can be changed and adjusted as required. Following this card is the complete syntax table which

is ended with a terminator 0. The semantic sets are then entered each begins with its semantic set number followed by the corresponding semantic set which is terminated with a 0. An extra terminator, 0, is added to signify the end of the semantic sets. This is immediately followed by the source statements to be analyzed each of which is preceded with its assumed phrase type. The complete set of the source statements is similarly terminated with 0. The document is then completed with the three consecutive '+' characters. Any number of user's documents can be run at the same time. That is, different or the same problems can be solved by various documents while any mistakes that occurs in one document do not have any effect on any other documents. This is particularly desirable in testing runs or when the translator is used by many persons.

The complete set of user's documents is finally terminated with three terminal cards (terminal data).

Card 1	+++
Card 2	END OF RUN ;
Card 3	0;

Card 1 consists of three consecutive '+' characters and card 2 contains any comment concerning about the documents and ended with a semi-colon. Card 3 consists of the terminator 0. The complete input data listing for algebraic differentiation of restricted arithmetic expressions can be seen in Appendix A.

CHAPTER III

ALGEBRAIC DIFFERENTIATION

3.1. Syntax and Semantics for Differentiation.

The syntax employed for algebraic differentiation is shown on the opposite page.

This syntax for algebraic expressions observes the following conventions.

1. X is the independent variable.
Y is the dependent variable.
2. No definition of <identifier> is given and it is assumed that all non-numeric constants and variables are represented by single letters.
3. The derivative of Y is indicated by a ' symbol. Successive uses of ' indicate higher order derivatives.
4. All the following trigonometric functions may be used in the algebraic expressions :
exp, ln, sinh, cosh, tanh, coth, sech, csch,
arcsin, arccos, arctan, arccot, arcsec, arccsc,
sin, cos, tan, cot, sec, csc, arsinh, artanh,
arcoth, arsech, arcsch.

The argument of the function is enclosed in parenthesis as in normal programming practice.

SYNTAX FOR DIFFERENTIATION

```

11 <INDEPENDENT VARIABLE>=X;
12 <DEPENDENT VARIABLE> =Y;
13 <CONSTANT> =A|B|C|D|E|F|G|H|I|J|K|L|M|N|O|P|Q|R|S|T|U|V|W|Z;
16 <QUOTE> ='(<16>|<E>);
17 <DIGITS> =0|1|2|3|4|5|6|7|8|9;
18 <INTEGERS> =<17>(<18>|<E>);
19 <FRACTION> =<17>(<19>|<E>);
20 <UNSIGNED DECIMAL NUMBER> =<18>.<19>|. <19>|<18>;
21 <PRIMARY> =(<(><26><)>|<28>|<20>|<13>|<11>|<12>(<16>|<E>));
22 <SIGNED PRIMARY>=(+|-|<E>)<21>;
23 <FACTOR> =<22>(a<23>|<E>);
24 <TERM1> =<23>(/<24>|<E>);
25 <TERM2> =<24>(*<25>|<E>);
26 <S.A.E.> =<25>(+<26>|-<26>|<E>);
27 <STATEMENT> =<26>(<=><26>|<E>);
28 <FUNCTION> =(EXP|LN|SINH|COSH|TANH|COTH|SECH|CSCH|
  ARCSIN|ARCCOS|ARCTAN|ARCCOT|ARCSEC|ARCCSC|
  SIN|COS|TAN|COT|SEC|CSC|
  ARSINH|ARCOSH|ARTANH|ARCOTH|ARSECH|ARCSCH)<(><26><)>;

```

5. There are five binary operators in the following (descending) order of precedence :

@ is used for exponentiation
 / is used for division
 * is used for multiplication
 +, - are used respectively for addition and subtraction
 + - can also be used as unary operators.

A problem arises in connection with the associativity of the operators - , / and @. In ordinary mathematical convention operators of equal precedence are assumed to work from left to right (with the exception of @ which is from right to left), so that

$$A / B / C / D$$

is interpreted as

$$((A / B) / C) / D$$

The above syntax interprets the expression as

$$A / (B / (C / D))$$

since a left to right scanning syntax analyser is used. Parentheses are used in such cases to achieve the required associativity, and it is recommended that parentheses should always be inserted in all cases of this type.

With the syntax table so formed, all the following arithmetic expressions would be valid phrase types of 27.

$$Y = (1-X)/(1+X)$$

$$Y = (X^3 + 2X^2 - 3X - 2)/(X^2 - 1)$$

$$Y = \text{LN} (2 * \text{TAN} (X) + 1) / (\text{TAN} (X) + 2)$$

$$Y = (\text{SIN} (2 * X)) @ B$$

$$X * Y + X^2 - 1$$

$$Y * \text{SIN}(X) - X$$

$$Y * \text{SIN}(X) + Y * (1 * \text{COS}(X)) - 1$$

$$X^3 + Y^3 - 6 * X * Y = 0$$

$$Y + \text{COSH}(X) / (1 + \text{SINH}(X) @ 2)$$

$$Y = 5 * X / 3$$

$$A + B + C / D$$

$$45 - 3.6 - E$$

As an example take $Y = 5 * X / 3$ as a phrase type 27. The analysis record after 'search' would be as shown on the opposite page.

The target string is produced by application of the semantics to the analysis record.

The semantic rules in turn make use of the mathematical rules for algebraic differentiation. The table that follows indicates the mathematical rules

NEW SOURCE STATEMENT
 TYPE 27
 LENGTH 7

SOURCE STATEMENT :

Y=5*X/3

ANALYSED SUCCESSFULLY AS TYPE 27
 TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	7	2	9
2	26	3	1	1	3	0
3	25	2	1	1	4	0
4	24	2	1	1	5	0
5	23	2	1	1	6	0
6	22	3	1	1	7	0
7	21	7	1	1	8	9
8	12	1	1	1	9	0
9	26	3	3	7	10	0
10	25	1	3	7	11	18
11	24	2	3	3	12	0
12	23	2	3	3	13	0
13	22	3	3	3	14	0
14	21	3	3	3	15	0
15	20	3	3	3	16	0
16	18	2	3	3	17	18
17	17	6	3	3	0	0
18	25	2	5	7	19	0
19	24	1	5	7	20	24
20	23	2	5	5	21	0
21	22	3	5	5	22	0
22	21	5	5	5	23	0
23	11	1	5	5	24	0
24	24	2	7	7	25	0
25	23	2	7	7	26	0
26	22	3	7	7	27	0
27	21	3	7	7	28	0
28	20	3	7	7	29	0
29	18	2	7	7	30	31
30	17	4	7	7	0	0

TARGET STRING PRODUCED :

Y'=(0*X/3+5*(1*3-X*0)/3@2)

****END OUTPUT THIS DOCUMENT

required. The symbol p is used to represent <simple arithmetic expression> (i.e. <S.A.E.>) as defined in the syntax table for <S.A.E.>. The expression 'a' prior to differentiation is indicated by p_a and the derivative of p_a is indicated by $p_{a.b}$ (a and b are integers).

Original form of the Expression Differentiation rule for the Expression

$p_1 \pm p_2$	$p_{1.2} + p_{2.2}$
$p_1 * p_2$	$p_{1.2} * p_2 + p_1 * p_{2.2}$
p_2/p_1	$(p_{2.1}*p_1 - p_2*p_{1.2})/p_1^2$
$p_2 @ p_1$	$p_2 @ p_1 * ((p_1 * p_{2.2}) / p_2 + p_{1.2} * \ln(p_2))$
$\ln(p_1)$	$p_{1.2}/p_1$
$\exp(p_1)$	$p_{1.2} * \exp(p_1)$
$\sin(p_1)$	$p_{1.2} * \cos(p_1)$
$\cos(p_1)$	$-(p_{1.2} * \sin(p_1))$
$\tan(p_1)$	$p_{1.2} * \sec(p_1)^2$
$\cot(p_1)$	$-p_{1.2} * \csc(p_1)^2$
$\sec(p_1)$	$p_{1.2} * \sec(p_1) * \tan(p_1)$
$\csc(p_1)$	$-p_{1.2} * \csc(p_1) * \cot(p_1)$
$\arcsin(p_1)$	$p_{1.2} / (1 - p_1^2)^{(1/2)}$
$\arccos(p_1)$	$-p_{1.2} / (1 - p_1^2)^{(1/2)}$
$\arctan(p_1)$	$p_{1.2} / (1 + p_1^2)$

<u>Original form of the Expression</u>	<u>Differentiation rule for the Expression</u>
$\operatorname{arccot}(p_1)$	$-p_{1.2}/(1+p_1^2)$
$\operatorname{arcsec}(p_1)$	$p_{1.2}/(p_1^*(p_1^2-1)^{(1/2)})$
$\operatorname{arccsc}(p_1)$	$-p_{1.2}/(p_1^*(p_1^2-1)^{(1/2)})$
$\sinh(p_1)$	$p_{1.2}*\cosh(p_1)$
$\cosh(p_1)$	$p_{1.2}*\sinh(p_1)$
$\tanh(p_1)$	$p_{1.2}*\operatorname{sech}(p_1)^2$
$\operatorname{coth}(p_1)$	$-p_{1.2}*\operatorname{csch}(p_1)^2$
$\operatorname{sech}(p_1)$	$-p_{1.2}*\operatorname{sech}(p_1)*\tanh(p_1)$
$\operatorname{csch}(p_1)$	$-p_{1.2}*\operatorname{csch}(p_1)*\operatorname{coth}(p_1)$
$\operatorname{arsinh}(p_1)$	$p_{1.2}/(1+p_1^2)^{(1/2)}$
$\operatorname{arcosh}(p_1)$	$p_{1.2}/(p_1^2-1)^{(1/2)}$
$\operatorname{artanh}(p_1)$	$p_{1.2}/(1-p_1^2)$
$\operatorname{arcoth}(p_1)$	$p_{1.2}/(1-p_1^2)$
$\operatorname{arsech}(p_1)$	$-p_{1.2}/(p_1^*(1-p_1^2)^{(1/2)})$
$\operatorname{arcsh}(p_1)$	$-p_{1.2}/(p_1^*(1+p_1^2)^{(1/2)})$

In addition to the list of rules above, it is understood that the following additional rules must be observed:

1. The derivative of independent variable is 1.
2. The first derivative of dependent variable (Y) is indicated by the symbol ' after the variable (Y'). Succeeding derivatives of Y would be indicated by additional ' on left of Y.

3. The derivatives of all constants and numbers are zero.

The semantics corresponding to the syntax table on page 36, can then be shown on the opposite page.

With these semantic rules the generation of the target string from $Y = 5*X/3$ according to the analysis record on page 39 would be $Y^{\#} = (0*X/3+5*(1*3-X*0)/3@2)$. The mechanism that the compile routine has gone through the analysis record can be illustrated on the following table on page 44.

SEMANTIC TABLE FOR DIFFERENTIATION

11	<INDEPENDENT VARIABLE>=W<1>;
12	<DEPENDENT VARIABLE>=PW<'>;
13	<CONSTANT> =W<0>;
16	<QUOTE> =P;
20	<UNSIGNED DECIMAL NUMBER>=W<0>;
21	<PRIMARY> =K<1>W<(>C<1>W<1>). K<2,3,4,5,6,7>C<1>. K<6>C<2>.;
22	<SIGNED PRIMARY>=K<2>W<->. K<1,2,3>C<1>.;
23	<FACTOR> =K<1>L<1>W<@>L<2>W<*((>L<2>W<*>C<1>W<1>/>L<1>W<+>C<2> W<*LN(>L<1>W<1>))&K<2>C<1>.;
24	<TERM1> =K<1>W<(>C<1>W<*>L<2>W<->L<1>W<*>C<2>W<1>/>L<2>W<@2>& K<2>C<1>.;
25	<TERM2> =K<1>W<(>C<1>W<*>L<2>W<+>L<1>W<*>C<2>W<1>)& K<2>C<1>.;
26	<S.A.E.> =K<1,2,3>C<1>. K<1>W<+>. K<2>W<->. K<1,2>C<2>.;
27	<STATEMENT> =C<1>K<1>W<=>C<2>.;
28	<FUNCTION> =K<1>W<(>C<1>W<1>)*EXP(>L<1>W<1>)>& K<2>W<(>C<1>W<1>/(>L<1>W<1>)>& K<3>W<(>C<1>W<1>)*COSH(>L<1>W<1>)>& K<4>W<(>C<1>W<1>)*SINH(>L<1>W<1>)>& K<5>W<(>C<1>W<1>)*(SECH(>L<1>W<1>))@2>& K<6>W<(-(>C<1>W<1>))*(CSCH(>L<1>W<1>))@2>& K<7>W<(-(>C<1>W<1>))*SECH(>L<1>W<1>)*TANH(>L<1>W<1>)>& K<8>W<(-(>C<1>W<1>))*CSCH(>L<1>W<1>)*COTH(>L<1>W<1>)>& K<9>W<(>C<1>W<1>)/(1-(>L<1>W<1>@2)@(<1/2>)>& K<10>W<(-(>C<1>W<1>))/(1-(>L<1>W<1>@2)@(<1/2>)>& K<11>W<(>C<1>W<1>)/(1+(>L<1>W<1>@2)@& K<12>W<(-(>C<1>W<1>))/(1+(>L<1>W<1>@2)@& K<13>W<(>C<1>W<1>/((>L<1>W<1>)*((>L<1>W<1>@2-1)@(<1/2>))>& K<14>W<(-(>C<1>W<1>)/((>L<1>W<1>)*((>L<1>W<1>@2-1)@(<1/2>))>& K<15>W<(>C<1>W<1>)*COS(>L<1>W<1>)>& K<16>W<(-(>C<1>W<1>))*SIN(>L<1>W<1>)>& K<17>W<(>C<1>W<1>)*(SEC(>L<1>W<1>))@2>& K<18>W<(-(>C<1>W<1>))*((CSC(>L<1>W<1>))@2>& K<19>W<(>C<1>W<1>)*SEC(>L<1>W<1>)*TAN(>L<1>W<1>)>& K<20>W<(-(>C<1>W<1>))*CSC(>L<1>W<1>)*COT(>L<1>W<1>)>& K<21>W<(>C<1>W<1>)/(1+(>L<1>W<1>@2)@(<1/2>)>& K<22>W<(>C<1>W<1>/((>L<1>W<1>@2-1)@(<1/2>)>& K<23>W<(>C<1>W<1>)/(1-(>L<1>W<1>@2)@& K<24>W<(>C<1>W<1>)/(1-(>L<1>W<1>@2)@(<1/2>)>& K<25>W<(-(>C<1>W<1>)/((>L<1>W<1>)*(1-(>L<1>W<1>@2)@(<1/2>))>& K<26>W<(-(>C<1>W<1>)/((>L<1>W<1>)*(1+(>L<1>W<1>@2)@(<1/2>))>.;

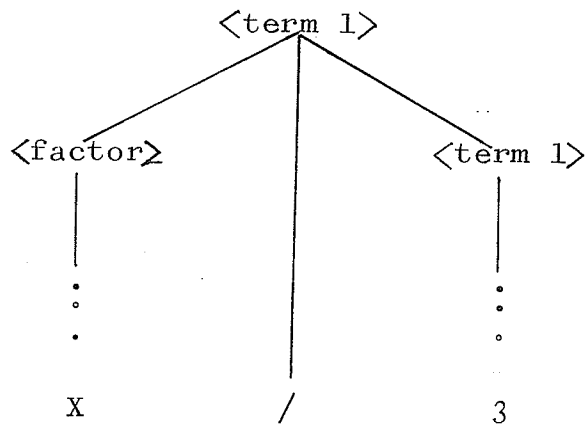
<u>entry j</u>	<u>Source string interpreted</u>	<u>target string produced</u>	<u>Instructions executed</u>	<u>Compile stack</u>
1			C <1>	0
2			C <1>	1
3			C <1>	2,1
4			C <1>	3,2,1
5			C <1>	4,3,2,1
6			C <1>	5,4,3,2,1
7			C <1>	6,5,4,3,2,1
8	Y	Y'	PW <'>	7,6,5,4,3,2,1
1	Y=	Y'='	W <=> C <2>	0
9	Y=	Y'='	C <1>	1
10	Y=	Y'=(W <() C <1>	9,1
11	Y=	Y'=(C <1>	10,9,1
12	Y=	Y'=(C <1>	11,10,9,1
13	Y=	Y'=(C <1>	12,11,10,9,1
14	Y=	Y'=(C <1>	13,12,11,10,9,1
15	Y=	Y'=(0	W <0>	14,13,12,11,10,9,1

<u>entry_j</u>	<u>Source string interpreted</u>	<u>target string produced</u>	<u>Instructions executed</u>	<u>Compile stack</u>
*16	Y=	Y=(0	-	-
*17	Y=5	Y'=(0	-	-
10	Y=5*	Y'=(0*X/3+5*	W<*>L<2>W<+>L<1>W<*>C<2>	9,1
18	Y=5*	Y'=(0*X/3+5*	C<1>	10,9,1
19	Y=5*	Y'=(0*X/3+5*(W<(>C<1>	18,10,9,1
20	Y=5*	Y'=(0*X/3+5*(C<1>	19,18,10,9,1
21	Y=5*	Y'=(0*X/3+5*(C<1>	20,19,18,10,9,1
22	Y=5*	Y'=(0*X/3+5*(C<1>	21,20,19,18,10,9,1
23	Y=5*X	Y'=(0*X/3+5*(1	W<1>	22,21,20,19,18,10,9,1
19	Y=5*X/	Y'=(0*X/3+5*(1*3-X*	W<*>L<2>W<->L<1>W<*>C<2>	18,10,9,1
24	Y=5*X/	Y'=(0*X/3+5*(1*3-X*	C<1>	19,18,10,9,1
25	Y=5*X/	Y'=(0*X/3+5*(1*3-X*	C<1>	24,19,18,10,9,1
26	Y=5*X/	Y'=(0*X/3+5*(1*3-X*	C<1>	25,24,19,18,10,9,1
27	Y=5*X/	Y'=(0*X/3+5*(1*3-X*	C<1>	26,25,24,19,18,10,9,1
28	Y=5*X/	Y'=(0*X/3+5*(1*3-X*0	W<0>	27,26,25,24,19,18,10,9,1
*29	Y=5*X/	Y'=(0*X/3+5*(1*3-X*0	-	-

<u>entry j</u>	<u>Source string interpreted</u>	<u>target string produced</u>	<u>Instructions executed</u>	<u>Compile stack</u>
*30	Y=5*X/3	Y=(0*X/3+5*(1*3-X*0	-	-
19	Y=5*X/3	Y=(0*X/3+5*(1*3-X*0)/3@2	W<1/>L<2>W<@2>	18,10,9,1
10	Y=5*X/3	Y=(0*X/3+5*(1*3-X*0)/3@2)	W<1>	9,1
1	Y=5*X/3	Y=(0*X/3+5*(1*3-X*0)/3@2)	-	0

* These entries do not have corresponding semantic interpretations, however, this is no handicap since the interpretation is already handled by some preceding entries.

The sub-tree formed by entry 19 and its succeeding entries are responsible for the formation of the portion of the target string $(1*3-X*0)/3@2$ from the substring $X/3$ of the source string.



The complete instructions for $\langle \text{term 1} \rangle$ at entry 19 is

$$W\langle \rangle C\langle 1 \rangle W\langle * \rangle L\langle 2 \rangle W\langle - \rangle L\langle 1 \rangle W\langle * \rangle C\langle 2 \rangle W\langle \rangle / \rangle L\langle 2 \rangle W\langle @ \ 2 \rangle$$

where C $\langle 1 \rangle$ produces a 1.

C $\langle 2 \rangle$ produces a 0.

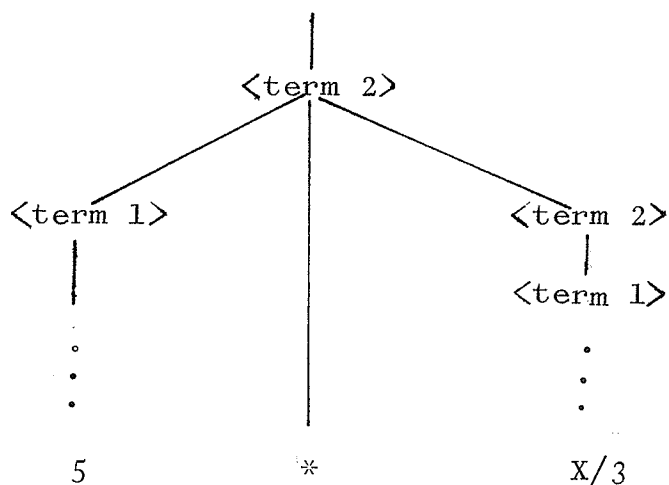
L $\langle 1 \rangle$ gives the first constituent of entry 19
i.e. X.

L $\langle 2 \rangle$ gives 3.

W $\langle \text{literal} \rangle$ gives the literal inside the
cornered brackets.

Therefore the resulting output from entry 19 is $(1*3-X*0)/3@2$ which is, in fact, the derivative of $X/3$.

Similarly the <term 2> at entry 10 corresponding to the formation of $(0*X/3+5*(1*3-X*0))/3@2$ from the substring $5*X/3$.



The complete instructions for <term 2> at this entry 10 is

$W\langle()C\langle 1 \rangle W\langle * \rangle L\langle 2 \rangle W\langle + \rangle L\langle 1 \rangle W\langle * \rangle C\langle 2 \rangle W\langle () \rangle$

where $C\langle 1 \rangle$ will finally produce a 5

$C\langle 2 \rangle$ compile recursively and finally goes

down to entry 19 which gives the resulting differentiation on $X/3$ i.e. $(1*3-X*0)/3@2$

$L\langle 1 \rangle$ and $L\langle 2 \rangle$ gives respectively 5 and $X/3$.

The complete differentiation of $5*X/3$ is $(0*X/3+5*(1*3-X*0)/3@2)$. Note the derivative of a quotient rule takes precedence over the product rule which is valid mathematically. This is due to the fact that though the source string is searched from left to right, the order for operators of equal precedence works from right to left in our approach in contrast to the ordinary convention (Appendix C). However, this is no handicap in differentiation, since the order of working on strings with series of multiplication or addition and subtraction (e.g. $A*B*C*D$ or $A+B-C-D+E$) would produce the same result though it may be in a different but equivalent form. The equal precedence for $*$ and $/$ assumed in conventional mathematics has been changed so that $/$ takes precedence over $*$.

For operator $/$ parentheses can always be inserted to ensure the proper order of operation. The derivative of $5*X/3$ is in fact trivial, and is solely for illustrative purpose. In fact, the syntax and semantic table we used above can work on quite complicated expressions including implicit functions. A few solutions have been worked out from the algebraic expressions as shown. Note that the last three solutions give examples on implicit functions.

OUTPUT LISTING ON DIFFERENTIATION

SOURCE STATEMENT :

 $Y = \sin(2X) @ B$

TARGET STRING PRODUCED :

 $Y' = \sin(2X) @ B * ((B * ((0 * X + 2 * 1)) * \cos(2 * X)) / \sin(2 * X) + 0 * \ln(\sin(2 * X)))$

SOURCE STATEMENT :

 $Y = \text{ARCOSH}(\text{SEC}(X))$

TARGET STRING PRODUCED :

 $Y' = ((1) * \text{SEC}(X) * \text{TAN}(X)) / ((\text{SEC}(X)) @ 2 - 1) @ (1/2)$

SOURCE STATEMENT :

 $Y = X @ 2 * \text{EXP}(A/X)$

TARGET STRING PRODUCED :

 $Y' = (X @ 2 * ((2 * 1) / X + 0 * \ln(X)) * \text{EXP}(A/X) + X @ 2 * ((0 * X - A * 1) / X @ 2) * \text{EXP}(A/X))$

SOURCE STATEMENT :

 $Y = (X @ N - X @ (-N)) / X$

TARGET STRING PRODUCED :

 $Y' = ((X @ N * ((N * 1) / X + 0 * \ln(X)) - X @ (-N) * (((-N) * 1) / X + (-0) * \ln(X))) * X - (X @ N - X @ (-N)) * 1) / X @ 2$

SOURCE STATEMENT :

 $X * Y + X @ 2 - 1$

TARGET STRING PRODUCED :

 $(1 * Y + X * Y') + X @ 2 * ((2 * 1) / X + 0 * \ln(X)) - 0$

SOURCE STATEMENT :

 $Y * \sin(X) - X$

TARGET STRING PRODUCED :

 $(Y' * \sin(X) + Y * (1) * \cos(X)) - 1$

SOURCE STATEMENT :

 $X @ 3 + Y @ 3 - 6 * X * Y = 0$

TARGET STRING PRODUCED :

 $X @ 3 * ((3 * 1) / X + 0 * \ln(X)) + Y @ 3 * ((3 * Y') / Y + 0 * \ln(Y)) - (0 * X * Y + 6 * (1 * Y + X * Y')) = 0$

*****END OUTPUT THIS DOCUMENT

3.2. Simplification of the Form of the Derivative.

It is obvious, from the examples given, that the derivatives obtained are mathematically inelegant and consist of some superfluous terms and redundant parentheses. In order to avoid these, two alternatives can be considered.

1. Redefine the syntax and semantic tables with provisions to handle constants or constant expressions.
2. Supply another syntax and semantic tables which will simplify the target string.

The second alternative is chosen for the reasons that it can demonstrate the otherwise obscure basic principles of differentiation and simplification.

In this simplification, no attempt is made to combine like terms or factors or to simplify trigonometric expressions, so that the expressions

$$A + B - 4 * A$$

$$A * C - A * D * \text{SIN} (X)$$

$$(X+1) * (X-1)/(X+1)$$

and $X * (1+\text{TAN} (X) @2)$

would remain unchanged and not simplify to

$$B - 3 * A$$

$$A * (C - D * \text{SIN } (X))$$

$$X - 1$$

$$X * \text{SEC } (X) @2$$

since this would be in the more involved field of symbolic or polynomial manipulation [6, 7, 25, 36, 38] , however, appropriate simplifications on terms with 1's and/or 0's are made. That is, the following basic transformations are always made.

$$A + 0 \text{ or } 0 + A \quad \longrightarrow \quad A$$

$$A * 0 \text{ or } 0 * A \quad \longrightarrow \quad 0$$

$$1 * A \text{ or } A * 1 \quad \longrightarrow \quad A$$

$$A @ 0 \quad \longrightarrow \quad 1$$

$$A @ 1 \quad \longrightarrow \quad A$$

$$1 @ A \quad \longrightarrow \quad 1$$

$$0 @ A \quad \longrightarrow \quad 0$$

Where A may be in any form from a single non-zero symbol A to an algebraic expression enclosed by parentheses, whereas 0(1) may be a digit 0(1) or any expression that after simplification becomes a 0(1).

In the syntax for simplification, we must define separately the non-zero expression, the zero expression, the various form of 1 and the various form of -1. In the corresponding semantic rules for the class 1 and -1 would be respectively $W\langle 1 \rangle$ and $W\langle -1 \rangle$. The various

combinations of the non-zero and zero expressions are also taken care of in the syntax definitions with the appropriate interpretation in the corresponding semantic rules.

The complete simplification syntax and semantic table designed for this purpose is on pages 54 and 55.

With this syntax and semantics the output strings of the differentiation on the expressions on page 50 can be fed in as source strings and produce simplified target strings as shown on page 56.

The appearance of some parentheses on some of these terms is in fact due to the omission of some of the zero and one terms. A complete omission of the parentheses is not intended for this syntax and semantic table (let's call it SIM1), but handled by a second syntax and semantic tables (call SIM2). This SIM2 will simplify the general expressions like

$$\begin{array}{ll} A/((X)@2) & A/X@2 \\ A*(B/C) & A*B/C \\ \text{or } A*((A+B*(C)))+C & A*(A+B*C)+C \end{array}$$

The resulting expressions from SIM1 on page 54, when passed through SIM2 (page 57) would give outputs as shown on page 58.

SIML

```

11 <ADD OP>          =+|-;
12 <MULP OP>        =*;
13 <DIV OP>         =/;
14 <ZERO>           =( <11>|<E> )0;
15 <ZERO TYPE>      =<11><( ><15>< )>|<( ><15>< )>|<14>(< * <15>|<E> );
16 <ONE>            =1;
17 <DIGITS>         =2|3|4|5|6|7|8|9;
18 <NONZERO INT>    =<17>(<N>|<E>)|<16><N>|<14><18>;
19 <NONZERO UNSIGNED DEC NO.>=<N>.|.|<E>|<18>;
20 <NONZERO DEC NO> =(<11>|<E>)<19>;
21 <CONSTANT>      =<L>(<34>|<E>)|<20>;
22 <FUNCTION>      =<44><( ><27>< )>;
23 <PRIMARY>       =<22>|<21>|<( ><37>< )>|<( ><27>< )>;
24 <FACTOR>        =1@<24>|<23>(@<24>|@<15>|@<41>|<E>);
25 <TERM2>         =1/<25>|<24>(/<25>|/<41>|<E>);
26 <TERM1>         =(<41>|<43>)*<26>|
                   <41>*<41>|
                   <25>(< * <26>|< * <41>|< * <43>|<E>);
27 <S.A.E>         =<41><11><27>|<26>(<11><27>|<E>)|<41>|<42>;
28 <ZERO FACTOR>   =<31>(@<24>|<E>)|<15>;
29 <ZERO TERM2>    =<28>(/<26>|<E>);
30 <ZERO TERM1>    =<29>*<26>|(<24>*|<E>)<29>;
31 <ZERO EXP>      =<( ><32>< )>;
32 <ZERO ARITH EXP>=<30>(<11><32>|<E>);
33 <EXPRESSION>    =<32>(< + <27>|< - <27>|<E> )|<26>(<11><32>|
                   <11><27>|<E>);
34 <QUOTE>         ='(<34>|<E>);
36 <STATEMENT>     =<37>(<=><37>|<E>);
37 <SIGNED EXPRESSION>=(<11>|<E>)<39>;
39 <COMBINE EXP.>  =<33>(<11><39>|<E>);
40 <POSITIVE ONE>  =0+1|1*1|1<11>0|1|+<40>|<( ><40>< )>;
41 <BRACKETED POS. ONE>=<( ><40>< )>|<( >-<43>< )>|<-<43>|1;
42 <NEGATIVE ONE>  =0-1|<-<41>;
43 <BRACKETED NEG ONE>=<( ><43>< )>|<( ><42>< )>;
44 <FUNCTION NAME> =EXP|LN|SINH|COSH|TANH|COTH|SECH|CSCH|
                   ARCSIN|ARCCOS|ARCTAN|ARCCOT|ARCSEC|ARCCSC|
                   SIN|COS|TAN|COT|SEC|CSC|
                   ARSINH|ARCOSH|ARTANH|ARCOATH|ARSECH|ARCSCH;

0;
1;
11 <ADD OP>          =P;
12 <MULP OP>        =P;
13 <DIV OP>         =P;
14 <ZERO>           =W<0>;
15 <ZERO>           =W<0>;
16 <ONE>            =P;
17 <DIGITS>         =P;
18 <NONZERO INT>    =K<1,2,3,4>C<1>.
                   K<1,3>L<2>.
                   K<4>C<2>.;
19 <NONZERO UNSIGNED DEC NO.>=K<1>L<1>W<.>C<2>&
                   K<2>W<.>C<1>&
                   K<3>C<1>.;
20 <NONZERO DEC NO>  =C<1>K<1>C<2>.;
21 <CONSTANT>      =K<1,2>L<1>.
                   K<1>C<2>.
                   K<3>C<1>.;
22 <FUNCTION>      =C<1>W<( >C<2>W< )>;
23 <PRIMARY>       =K<3,4>W<( >.

```

24	<FACTOR>	K<1,2,3,4>C<1>. K<3,4>W<1>.; =K<1,3>W<1>. K<2,4,5>C<1>. K<2>W<@>C<2>.;
25	<TERM2>	=K<1>W<1/>C<1>. K<2,3,4>C<1>. K<2>W</>C<2>.;
26	<TERM1>	=K<2,6>W<->. K<4,5,6,7>C<1>. K<4>W<*>. K<1,2,4>C<2>. K<2,6>W<1>. K<3>W<1>.;
27	<S.A.E.>	=K<1,2,3,4,5>C<1>. K<1,2>C<2>C<3>.;
33	<EXPRESSION>	=K<4,5,6>C<1>. K<2>W<->. K<3>W<0>. K<5>C<2>C<3>. K<1,2>C<2>.;
34	<QUOTE>	=P;
36	<STATEMENT>	=C<1>K<1>W<=>C<2>.;
37	<SIGNED EXPRESSION>	=C<1>K<1>C<2>.;
39	<COMBINE EXP.>	=C<1>;
41	<BRACKETED POS. ONE>	=W<1>;
42	<NEGATIVE ONE>	=W<-1>;
44	<FUNCTION NAME>	=P;

OUTPUT LISTING AFTER SIM1

SOURCE STATEMENT :

$$Y' = \sin(2*X) @ B * ((B * ((0*X + 2*1)) * \cos(2*X)) / \sin(2*X) + 0 * \ln(\sin(2*X)))$$

TARGET STRING PRODUCED :

$$Y' = \sin(2*X) @ B * ((B * ((2)) * \cos(2*X)) / \sin(2*X))$$

SOURCE STATEMENT :

$$Y' = ((1) * \sec(X) * \tan(X)) / ((\sec(X)) @ 2 - 1) @ (1/2)$$

TARGET STRING PRODUCED :

$$Y' = (\sec(X) * \tan(X)) / ((\sec(X)) @ 2 - 1) @ (1/2)$$

SOURCE STATEMENT :

$$Y' = (X @ 2 * ((2*1) / X + 0 * \ln(X)) * \exp(A/X) + X @ 2 * ((0*X - A * 1) / X @ 2) * \exp(A/X))$$

TARGET STRING PRODUCED :

$$Y' = (X @ 2 * ((2) / X) * \exp(A/X) + X @ 2 * ((-A) / X @ 2) * \exp(A/X))$$

SOURCE STATEMENT :

$$Y' = ((X @ N * ((N * 1) / X + 0 * \ln(X)) - X @ (-N) * (((-N) * 1) / X + (-0) * \ln(X))) * X - (X @ N - X @ (-N)) * 1) / X @ 2$$

TARGET STRING PRODUCED :

$$Y' = ((X @ N * ((N) / X) - X @ (-N) * (((-N)) / X)) * X - (X @ N - X @ (-N))) / X @ 2$$

SOURCE STATEMENT :

$$(1 * Y + X * Y') + X @ 2 * ((2*1) / X + 0 * \ln(X)) - 0$$

TARGET STRING PRODUCED :

$$(Y + X * Y') + X @ 2 * ((2) / X)$$

SOURCE STATEMENT :

$$(Y' * \sin(X) + Y * (1) * \cos(X)) - 1$$

TARGET STRING PRODUCED :

$$(Y' * \sin(X) + Y * \cos(X)) - 1$$

SOURCE STATEMENT :

$$X @ 3 * ((3*1) / X + 0 * \ln(X)) + Y @ 3 * ((3 * Y') / Y + 0 * \ln(Y)) - (0 * X * Y + 6 * (1 * Y + X * Y')) = 0$$

TARGET STRING PRODUCED :

$$X @ 3 * ((3) / X) + Y @ 3 * ((3 * Y') / Y) - (6 * (Y + X * Y')) = 0$$

****END OUTPUT THIS DOCUMENT

SIM2

11	<ADD OP>	=+ -;
13	<QUOTE>	'(<13> <E>);
14	<DECIMAL NUMBER>	=<N>.<N> . <N> <N>;
15	<CONSTANT>	=<L>(<13> <E>) <14>;
16	<POSITIVE CONSTANT>	=<15> <(><16><)>;
17	<NEGATIVE CONSTANT>	=<(>-<15><)> <(><17><)>;
18	<SYMBOL>	=<16> <17>;
19	<PRIMARY>	=<31> <29> <30> <18> <(><(><24><)><)> <(><24><)>;
20	<SIGNED PRIMARY>	=(<+ -> <E>)<19>;
21	<FACTOR>	=<(><20>@<20><)> <20>@<30> @<20> <E>;
22	<TERM1>	=<21>(</><22> <E>);
23	<TERM2>	=<22>(<*><23> <E>);
24	<S.A.E.>	=<23>(<11><24> <E>);
25	<FUNCTION NAME>	=EXP LN SINH COSH TANH COth SECH CSCH ARCSIN ARCCOS ARCTAN ARCCOT ARCSEC ARCCSC SIN COS TAN COT SEC CSC ARSINH ARCOSH ARTANH ARCOth ARSECH ARCSCH;
27	<STATEMENT>	=<24>(<=><24> <E>);
28	<PROGRAM>=<27>;	
29	<FUNCTION>	=<25><(><24><)>;
30	<BRAC CONS>	=<(><18>/<18><)>;
31	<BRACKETED TERM>	=<(><29><)>;
0;		
1;		
11	<ADD OP>	=P;
13	<QUOTE>	=P;
14	<DECIMAL NUMBER>	=K<1>L<1>W<.>L<2>& K<2>W<.>L<1>& K<3>L<1>.;
15	<CONSTANT>	=K<1>L<1>C<2>& K<2>L<1>& K<3>C<1>.;
16	<POSITIVE CONSTANT>	=C<1>;
17	<NEGATIVE CONSTANT>	=K<1>W<(->C<1>W<)>. K<2>C<1>.;
18	<SYMBOL>	=C<1>;
19	<PRIMARY>	=K<1,2,3,4>C<1>. K<5,6>W<(>C<1>W<)>.;
20	<SIGNED PRIMARY>	=K<2>W<->. K<1,2,3>C<1>.;
21	<FACTOR>	=K<1,2,3,4>C<1>. K<1,3>W<@>C<2>. K<2>W<@(>C<2>W<)>.;
22	<TERM1>	=C<1>K<1>W</>C<2>.;
23	<TERM2>	=C<1>K<1>W<*>C<2>.;
24	<S.A.E.>	=C<1>K<1>C<2>C<3>.;
25	<FUNCTION NAME>	=P;
27	<STATEMENT>	=C<1>K<1>W<=>C<2>.;
28	<PROGRAM>=C<1>;	
29	<FUNCTION>	=C<1>W<(>C<2>W<)>;
30	<BRAC CONS>	=C<1>W</>C<2>;
31	<BRACKETED TERM>	=C<1>;

OUTPUT LISTING AFTER SIM2

SOURCE STATEMENT :

$$Y' = \sin(2 * X) @ B * ((B * ((2)) * \cos(2 * X)) / \sin(2 * X))$$

TARGET STRING PRODUCED :

$$Y' = \sin(2 * X) @ B * ((B * 2 * \cos(2 * X)) / \sin(2 * X))$$

SOURCE STATEMENT :

$$Y' = (\sec(X) * \tan(X)) / ((\sec(X)) @ 2 - 1) @ (1/2)$$

TARGET STRING PRODUCED :

$$Y' = (\sec(X) * \tan(X)) / (\sec(X) @ 2 - 1) @ (1/2)$$

SOURCE STATEMENT :

$$Y' = (X @ 2 * ((2) / X) * \exp(A / X) + X @ 2 * ((-A) / X @ 2) * \exp(A / X))$$

TARGET STRING PRODUCED :

$$Y' = (X @ 2 * 2 / X * \exp(A / X) + X @ 2 * ((-A) / X @ 2) * \exp(A / X))$$

SOURCE STATEMENT :

$$Y' = ((X @ N * ((N) / X) - X @ (-N) * (((-N)) / X)) * X - (X @ N - X @ (-N))) / X @ 2$$

TARGET STRING PRODUCED :

$$Y' = ((X @ N * N / X - X @ (-N) * (-N) / X) * X - (X @ N - X @ (-N))) / X @ 2$$

SOURCE STATEMENT :

$$(Y + X * Y') + X @ 2 * ((2) / X)$$

TARGET STRING PRODUCED :

$$(Y + X * Y') + X @ 2 * 2 / X$$

SOURCE STATEMENT :

$$(Y' * \sin(X) + Y * \cos(X)) - 1$$

TARGET STRING PRODUCED :

$$(Y' * \sin(X) + Y * \cos(X)) - 1$$

SOURCE STATEMENT :

$$X @ 3 * ((3) / X) + Y @ 3 * ((3 * Y') / Y) - (6 * (Y + X * Y')) = 0$$

TARGET STRING PRODUCED :

$$X @ 3 * 3 / X + Y @ 3 * ((3 * Y') / Y) - (6 * (Y + X * Y')) = 0$$

****END OUTPUT THIS DOCUMENT

In some long algebraic expressions a single pass through SIM1 and SIM2 may not result in its simplest form. For example, the target string

$$W' = ((0 * \cos(X-Y) + 0 * (- (1-0)) * \sin(X-Y)) + (1 * (- (0-1)) * \sin(X-Y) + X * ((- (0-0)) * \sin(X-Y) + (- (0-1)) * (1-0) * \cos(X-Y)))));$$

on passing through SIM1 and SIM2 would give

After SIM1

$$W' = ((\sin(X-Y) + X * ((-0) * \sin(X-Y) + \cos(X-Y)))) ;$$

After SIM2

$$W' = (\sin(X-Y) + X * ((-0) * \sin(X-Y) + \cos(X-Y)));$$

which is not the simplest result and if again passed through SIM1 and SIM2 the immediate and final results are

After SIM1

$$W' = (\sin(X-Y) + X * (\cos(X-Y)))$$

After SIM2

$$W' = (\sin(X-Y) + X * \cos(X-Y));$$

which is in fact the simplest form that can be achieved, since there is no more zero or one terms or presence of any superfluous parentheses. Actually we can pass the output on differentiation through SIM1 and SIM2 alternately until the target string that produced is of constant length.

3.3. Evaluating the Derivative at Assigned Points.

The above description is concerned with finding the derivative of a function and presenting the result as an algebraic expression. It may be required to calculate the numerical value of such a derivative. The system allows for this by rewriting `COMPILE` as a type procedure in which the resulting value of a component can be assigned. This process, however, cannot be described in the external semantics.

A type procedure is written to handle the mathematical calculations. The analysis record from the `SEARCH` routine can be made use of by this procedure via

```
j := link [i,j] ; i := type [j] ; k := cat [j] ;
```

as has been discussed on section 2.5.2. The syntax table used to define the source string is exactly the same as that for algebraic differentiation, but the interpretation of the syntax rules (semantics) is entirely different. For operators `+`, `-`, `*`, `/` they carry the ordinary mathematical function. Thus, if $A = 5$, $B = 3$, the expressions

```
C = A + B would set C = 8 after compilation
```

```
C = A - B would set C = 2 after compilation
```

```
C = A * B would set C = 15 after compilation
```

```
C = A / B would set C = 1.666667 after compilation
```

however `C = A @ B` is interpreted as `C = A ** B` where `**`

is the unary operator of exponentiation in the ALGOL programming language, and would give the result $C = 125$ after compilation. Floating point mathematics is assumed in all cases.

On the evaluation of functions we take advantages of the built-in functions that are provided in the ALGOL programming language.

EXP(X)
 SIN(X)
 COS(X)
 ABS(X)
 SQRT(X)
 ARCTAN(X)
 LN(X)

where X is an argument that could be a simple variable, arithmetic expression or another function. With these given built-in functions evaluation of other functions can be done with the following given formulas.

$$\tan x = \frac{\sin x}{\cos x}$$

$$\cot x = \frac{\cos x}{\sin x}$$

$$\sec x = \frac{1}{\cos x}$$

$$\csc x = \frac{1}{\sin x}$$

$$\begin{aligned}
\arcsin x &= \arctan \left(\frac{x}{\sqrt{1-x^2}} \right) \\
\arccos x &= \frac{\pi}{2} - \arctan \left(\frac{x}{\sqrt{1-x^2}} \right) \\
\operatorname{arccot} x &= \frac{\pi}{2} - \arctan x \\
\operatorname{arcsec} x &= \arctan \left(\sqrt{x^2-1} \right) \quad \text{for } x \geq 0 \\
\operatorname{arcsec} x &= \arctan \left(\sqrt{x^2-1} - \pi \right) \quad \text{for } x < 0 \\
\operatorname{arccsc} x &= \arctan \left(\frac{1}{\sqrt{x^2-1}} \right) \quad \text{for } x \geq 0 \\
\operatorname{arccsc} x &= \arctan \left(\frac{1}{\sqrt{x^2-1}} \right) - \pi \quad \text{for } x < 0 \\
\sinh x &= \frac{e^x - e^{-x}}{2} \\
\cosh x &= \frac{e^x + e^{-x}}{2} \\
\tanh x &= \frac{e^x - e^{-x}}{e^x + e^{-x}} \\
\operatorname{coth} x &= \frac{e^x + e^{-x}}{e^x - e^{-x}} \\
\operatorname{sech} x &= \frac{2}{e^x + e^{-x}} \\
\operatorname{csch} x &= \frac{2}{e^x - e^{-x}} \\
\sinh^{-1} \frac{x}{a} &= \ln \left(\frac{x + \sqrt{x^2 + a^2}}{a} \right), \quad a > 0 \\
\cosh^{-1} \frac{x}{a} &= \ln \left(\frac{x + \sqrt{x^2 - a^2}}{a} \right), \quad a > 0
\end{aligned}$$

$$\begin{aligned} \tanh^{-1} \frac{x}{a} &= \frac{1}{2} \ln \left| \frac{a+x}{a-x} \right|, & |x| < |a| \\ \coth^{-1} \frac{x}{a} &= \frac{1}{2} \ln \left| \frac{a+x}{a-x} \right|, & |x| > |a| \\ \operatorname{sech}^{-1} \frac{x}{a} &= \ln \left(\frac{a}{x} + \sqrt{\frac{a^2}{x^2} - 1} \right), & a > 0 \\ \operatorname{csch}^{-1} \frac{x}{a} &= \ln \left(\frac{a}{x} + \sqrt{\frac{a^2}{x^2} + 1} \right), & a > 0 \end{aligned}$$

The value routine that handles all the evaluations can then be shown on pages 64, 65 and 66.

The derivative of $Y = \operatorname{ARCOSH}(\operatorname{SEC}(X))$ would be $Y' = ((1)*\operatorname{SEC}(X)*\operatorname{TAN}(X))/((\operatorname{SEC}(X))^2-1)^{(1/2)}$ and if $X = 1.000$ radians, can be evaluated through this 'value' routine to give the numeric result of 1.851. The complete analysed result is shown on pages 67, 68 and 69.

In the syntax the operator / remains to take the higher precedence than * since it always provides the correct result as if they are of equal precedence and works from left to right. Hence the expression $3*5/4*3$ would give the same result in both cases as illustrated below.

$\begin{array}{c} \underbrace{3*5/4*3}_{T_1} \\ \underbrace{\quad\quad\quad}_{T_2} \\ \underbrace{\quad\quad\quad}_{T_3} \end{array}$	$\begin{array}{c} \underbrace{3*5/4*3}_{T_1} \\ \underbrace{\quad\quad\quad}_{T_2} \\ \underbrace{\quad\quad\quad}_{T_3} \end{array}$
---	---

```

'REAL' 'PROCEDURE' VALUE(I,J); 'VALUE' I,J; 'INTEGER' I,J;
  'BEGIN' 'INTEGER' K; 'REAL' L,Y;
  'SWITCH' S:=INDEPENDENT VARIABLE,DEP VARIABLE,CONSTANT,ALPHABETS,
    DUMMY1,QUOTE,DIGITS,INTEGERS,FRACTION PART,UNSIGNED DECIMAL NO,
    PRIMARY,SIG PRIMARY,FACTOR,TERM1,TERM2,SAE;
  'SWITCH' SS:=STATEMENT,FUNCTION NAME;
  'BOOLEAN' LTAO,LTMO;
LTAO:=TAO;
LTMO:=TMO;
J:=LINK(/I,J/); I:=TYPE(/J/); K:=CAT(/J/);
'IF' I>26 'THEN' 'GOTO' SS(/I-26/); 'GOTO' S(/I-10/);
INDEPENDENT VARIABLE:
  VALUE:=IX; 'GOTO' END;
DEP VARIABLE:
  PRINTCH(VAR(/2/)); 'GOTO' END;
CONSTANT: VALUE:=A(/K/); 'GOTO' END;
ALPHABETS: 'GOTO' END;
DUMMY1: 'GOTO' END;
QUOTE:
  OUTSYMBOL(1,('''),1);
  'IF' K=1 'THEN' VALUE(1,J);
  'GOTO' END;
DIGITS: VALUE:=SA(/K/);
  'GOTO' END;
INTEGERS: 'IF' K=1 'THEN' VALUE:=10*VALUE(1,J)+VALUE(2,J)
  'ELSE' VALUE:=VALUE(1,J);
  'GOTO' END;
FRACTION PART:
  'IF' K=1 'THEN' VALUE:=VALUE(1,J)/10+VALUE(2,J)/100
  'ELSE' VALUE:=VALUE(1,J)/10;
  'GOTO' END;
UNSIGNED DECIMAL NO:
  'IF' K=1 'THEN' VALUE:=VALUE(1,J)+VALUE(2,J)
  'ELSE' VALUE:=VALUE(1,J);
  'GOTO' END;
PRIMARY: 'IF' K=1 'THEN' 'BEGIN'
  TAO:='TRUE';
  TMO:='TRUE';
  VALUE:=VALUE(1,J);
  TAO:=LTAO; TMO:=LTMO
  'END'
  'ELSE' VALUE:=VALUE(1,J);
  'IF' K=6 'THEN' VALUE(2,J);
  'GOTO' END;
SIG PRIMARY:
  'IF' K=2 'THEN' VALUE:=-VALUE(1,J)
  'ELSE' VALUE:=VALUE(1,J);
  'GOTO' END;
FACTOR: 'IF' K=1 'THEN' 'BEGIN' L:=VALUE(2,J);
  VALUE:=VALUE(1,J)**L
  'END'
  'ELSE' VALUE:=VALUE(1,J);
  'GOTO' END;
TERM1:
  'IF' K=1 'THEN' 'BEGIN'
  'IF' LTMO 'THEN' 'BEGIN'
  Y:=VALUE(1,J);
  TMO:='FALSE';
  VALUE:=Y/VALUE(2,J)

```

```

                                'END'
      'ELSE' VALUE:=VALUE(1,J)*VALUE(2,J)
                                'END'
      'ELSE' VALUE:=VALUE(1,J);
      'GOTO' END;
TERM2:
      'IF' K=1 'THEN' 'BEGIN'
      Y:=VALUE(1,J);
      TMO:='TRUE';
      VALUE:=Y*VALUE(2,J)
                                'END'
      'ELSE' VALUE:=VALUE(1,J);
      'GOTO' END;
SAE:   'IF' K=1 'THEN'
      'BEGIN'
      TMO:='TRUE';
      Y:=SIGN('IF' LTAO 'THEN' 1 'ELSE' -1)*VALUE(1,J);
      TAO:='TRUE' ;
      TMO:='TRUE';
      VALUE:=Y+VALUE(2,J);
      TMO:=LTMO
      'END'
      'ELSE' 'IF' K=2 'THEN'
      'BEGIN'
      TMO:='TRUE';
      Y:=SIGN('IF' LTAO 'THEN' 1 'ELSE' -1)*VALUE(1,J);
      TAO:='FALSE';
      TMO:='TRUE';
      VALUE:=Y+VALUE(2,J);
      TMO:=LTMO
      'END'
      'ELSE' VALUE:=SIGN('IF' LTAO 'THEN' 1 'ELSE'-1)*VALUE(1,J);
      'GOTO' END;
STATEMENT:
      VALUE(1,J); OUTST('('=')'); VALUE:=VALUE(2,J);
      'GOTO' END;
FUNCTION NAME:
      'BEGIN'
      'REAL' 'PROCEDURE' EXP1(X); 'REAL' X;
      EXP1:=EXP(X)-EXP(-X);
      'REAL' 'PROCEDURE' EXP2(X) ; 'REAL' X ;
      EXP2:=EXP(X)+EXP(-X);
      'SWITCH' S:=AEXP,ALN,SINH,COSH,TANH,COTH,SECH,CSCH,
      ARCSIN,ARCCOS,AARCTAN,ARCCOT,ARCSEC,ARCCSC,ASIN,ACOS;
      'SWITCH' SS:=TAN,COT,SEC,CSC,ARSINH,ARCOSH,ARTANH,ARCOTH,ARSECH,
      ARCSCH;
      'REAL' X,PI;
      TAO:='TRUE'; TMO:='TRUE' ;
      X:=VALUE(1,J); PI:=3.141593;
      'IF' K<=16 'THEN' 'GOTO' S(/K/) 'ELSE' 'GOTO' SS(/K-16/);
      AEXP: VALUE:=EXP(X); 'GOTO' END;
      ALN: VALUE:=LN(X); 'GOTO' END;
      SINH: VALUE:=EXP1(X)/2; 'GOTO' END;
      COSH: VALUE:=EXP2(X)/2; 'GOTO' END;
      TANH: VALUE:=EXP1(X)/EXP2(X) ; 'GOTO' END;
      COTH: 'IF' EXP1(X)=0 'THEN' 'GOTO' OVERFLO;
      VALUE:=EXP2(X)/EXP1(X);
      'GOTO' END;
      SECH: VALUE:=2/EXP2(X) ; 'GOTO' END;
      CSCH: 'IF' EXP1(X)=0 'THEN' 'GOTO' OVERFLO;

```

```

                                VALUE:=2/EXP1(X); 'GOTO' END;
ARCSIN: 'IF' X=1 'THEN' 'GOTO' OVERFLO;
                                VALUE:=ARCTAN(X/SQRT(1-X**2)); 'GOTO' END;
ARCCOS: 'IF' X=1 'THEN' 'GOTO' OVERFLO;
                                VALUE:=PI/2-ARCTAN(X/SQRT(1-X**2)); 'GOTO' END;
AARCTAN: VALUE:=ARCTAN(X); 'GOTO' END;
ARCCOT: VALUE:=PI/2-ARCTAN(X); 'GOTO' END;
ARCSEC: 'IF' X>=0 'THEN' VALUE:=ARCTAN(SQRT(X**2-1))
        'ELSE' VALUE:=ARCTAN(SQRT(X**2-1)-PI); 'GOTO' END;
ARCCSC: 'IF' X>=0 'THEN' VALUE:=ARCTAN(1/SQRT(X**2-1))
        'ELSE' VALUE:=ARCTAN(1/SQRT(X**2-1))-PI; 'GOTO' END;
ASIN: VALUE:=SIN(X); 'GOTO' END;
ACOS: VALUE:=COS(X); 'GOTO' END;
TAN: 'IF' COS(X)=0 'THEN' 'GOTO' OVERFLO;
      VALUE:=SIN(X)/COS(X); 'GOTO' END;
COT: 'IF' SIN(X)=0 'THEN' 'GOTO' OVERFLO;
      VALUE:=COS(X)/SIN(X); 'GOTO' END;
SEC: 'IF' COS(X)=0 'THEN' 'GOTO' OVERFLO;
      VALUE:=1/COS(X); 'GOTO' END;
CSC: 'IF' SIN(X)=0 'THEN' 'GOTO' OVERFLO;
      VALUE:=1/SIN(X); 'GOTO' END;
ARSINH: VALUE:=LN(X+SQRT(X**2+1)); 'GOTO' END;
ARCOSSH: VALUE:=LN(X+SQRT(X**2-1)); 'GOTO' END;
ARTANH: 'IF' ABS(X)>=1 'THEN' 'GOTO' OVERFLO;
        VALUE:=0.5*LN(ABS((1+X)/(1-X))); 'GOTO' END;
ARCOOTH: 'IF' ABS(X)<=1 'THEN' 'GOTO' OVERFLO;
        VALUE:=0.5*LN(ABS((1+X)/(X-1))); 'GOTO' END;
ARSECH: VALUE:=LN(1/X+SQRT(1/X**2-1)); 'GOTO' END;
ARCSCCH: VALUE:=LN(1/X+SQRT(1/X**2+1));
END: TAO:=LTAO; TMO:=LTMO;
      'END' ;
END: 'END' VALUE;

```

TYPE 27
 LENGTH 16
 HIGHEST ORDER OF DIFFERENTIATION ACHIEVED 1
 CODE FOR EVALUATION IS 1 ELSE 0, CODE= 1

SOURCE STATEMENT :

Y=ARCOSH(SEC(X))

INDEPENDENT VARIABLE IS X DEPENDENT VARIABLE IS Y
 ANALYSED SUCCESSFULLY AS TYPE 27

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	16	2	9
2	26	3	1	1	3	0
3	25	2	1	1	4	0
4	24	2	1	1	5	0
5	23	2	1	1	6	0
6	22	3	1	1	7	0
7	21	7	1	1	8	9
8	12	1	1	1	9	0
9	26	3	3	16	10	0
10	25	2	3	16	11	0
11	24	2	3	16	12	0
12	23	2	3	16	13	0
13	22	3	3	16	14	0
14	21	2	3	16	15	0
15	28	22	3	16	16	0
16	26	3	10	15	17	0
17	25	2	10	15	18	0
18	24	2	10	15	19	0
19	23	2	10	15	20	0
20	22	3	10	15	21	0
21	21	2	10	15	22	0
22	28	19	10	15	23	0
23	26	3	14	14	24	0
24	25	2	14	14	25	0
25	24	2	14	14	26	0
26	23	2	14	14	27	0
27	22	3	14	14	28	0
28	21	5	14	14	29	0
29	11	1	14	14	30	0

TARGET STRING PRODUCED :

$Y' = ((1) * \text{SEC}(X) * \text{TAN}(X)) / ((\text{SEC}(X))^2 - 1)^{(1/2)}$

X = 1.000

ANALYSED SUCCESSFULLY AS TYPE 27

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	43	2	10
2	26	3	1	2	3	0
3	25	2	1	2	4	0
4	24	2	1	2	5	0
5	23	2	1	2	6	0
6	22	3	1	2	7	0
7	21	6	1	2	8	9
8	12	1	1	1	9	0
9	16	2	2	2	10	0
10	26	3	4	43	11	0
11	25	2	4	43	12	0

12	24	1	4	43	13	57
13	23	2	4	22	14	0
14	22	3	4	22	15	0
15	21	1	4	22	16	0
16	26	3	5	21	17	0
17	25	1	5	21	18	31
18	24	2	5	7	19	0
19	23	2	5	7	20	0
20	22	3	5	7	21	0
21	21	1	5	7	22	0
22	26	3	6	6	23	0
23	25	2	6	6	24	0
24	24	2	6	6	25	0
25	23	2	6	6	26	0
26	22	3	6	6	27	0
27	21	3	6	6	28	0
28	20	3	6	6	29	0
29	18	2	6	6	30	31
30	17	2	6	6	0	0
31	25	1	9	21	32	44
32	24	2	9	14	33	0
33	23	2	9	14	34	0
34	22	3	9	14	35	0
35	21	2	9	14	36	0
36	28	19	9	14	37	0
37	26	3	13	13	38	0
38	25	2	13	13	39	0
39	24	2	13	13	40	0
40	23	2	13	13	41	0
41	22	3	13	13	42	0
42	21	5	13	13	43	0
43	11	1	13	13	44	0
44	25	2	16	21	45	0
45	24	2	16	21	46	0
46	23	2	16	21	47	0
47	22	3	16	21	48	0
48	21	2	16	21	49	0
49	28	17	16	21	50	0
50	26	3	20	20	51	0
51	25	2	20	20	52	0
52	24	2	20	20	53	0
53	23	2	20	20	54	0
54	22	3	20	20	55	0
55	21	5	20	20	56	0
56	11	1	20	20	57	0
57	24	2	24	43	58	0
58	23	1	24	43	59	96
59	22	3	24	37	60	0
60	21	1	24	37	61	0
61	26	2	25	36	62	87
62	25	2	25	34	63	0
63	24	2	25	34	64	0
64	23	1	25	34	65	81
65	22	3	25	32	66	0
66	21	1	25	32	67	0
67	26	3	26	31	68	0
68	25	2	26	31	69	0
69	24	2	26	31	70	0
70	23	2	26	31	71	0
71	22	3	26	31	72	0
72	21	2	26	31	73	0
73	28	19	26	31	74	0
74	26	3	30	30	75	0
75	25	2	30	30	76	0

76	24	2	30	30	77	0
77	23	2	30	30	78	0
78	22	3	30	30	79	0
79	21	5	30	30	80	0
80	11	1	30	30	81	0
81	23	2	34	34	82	0
82	22	3	34	34	83	0
83	21	3	34	34	84	0
84	20	3	34	34	85	0
85	18	2	34	34	86	87
86	17	3	34	34	0	0
87	26	3	36	36	88	0
88	25	2	36	36	89	0
89	24	2	36	36	90	0
90	23	2	36	36	91	0
91	22	3	36	36	92	0
92	21	3	36	36	93	0
93	20	3	36	36	94	0
94	18	2	36	36	95	96
95	17	2	36	36	0	0
96	23	2	39	43	97	0
97	22	3	39	43	98	0
98	21	1	39	43	99	0
99	26	3	40	42	100	0
100	25	2	40	42	101	0
101	24	1	40	42	102	108
102	23	2	40	40	103	0
103	22	3	40	40	104	0
104	21	3	40	40	105	0
105	20	3	40	40	106	0
106	18	2	40	40	107	108
107	17	2	40	40	0	0
108	24	2	42	42	109	0
109	23	2	42	42	110	0
110	22	3	42	42	111	0
111	21	3	42	42	112	0
112	20	3	42	42	113	0
113	18	2	42	42	114	115
114	17	3	42	42	0	0

TARGET STRING PRODUCED :
Y' = 1.851

***END OUTPUT THIS DOCUMENT

or	$T_1/4*3$	$T_1 = 15$		$3*T_1*3$	$T_1 = 1.25$
	T_2*3	$T_2 = 3.75$		$3*T_2$	$T_2 = 3.75$
	T_3	$T_3 = 11.25$		T_3	$T_3 = 11.25$

where T_i indicates the immediate result obtained at step i .

Note that for the operators of equal precedence, the operation proceeds from right to left according to the translator, flags must therefore be set on or off whenever necessary in order to carry the conventional order of operations. For example the expression $A-B+C-D-E$ if no control flags appended in the 'value' routine the order of operation would be as shown.

$$\begin{array}{c}
 A - B + C - D - E \\
 \quad \quad \quad \underbrace{\hspace{2cm}} \\
 \quad \quad \quad \quad T_1 \\
 \quad \quad \underbrace{\hspace{1.5cm}} \\
 \quad \quad \quad T_2 \\
 \quad \underbrace{\hspace{1cm}} \\
 \quad \quad T_3 \\
 \underbrace{\hspace{3cm}} \\
 \quad F
 \end{array}$$

It is obvious that such procedure of operation is wrong if it is assumed to work according to the conventional mathematics. Given $A=3$, $B=4$, $C=2$, $D=3$ and $E=5$ the intermediate values of T_1 , T_2 , T_3 are -2 , 4 , 8 respectively and the final value of the complete expression is -5 (i.e. $F = -5$) which is wrong. To achieve the correct result the above expression is interpreted as

ltao T F T F F

A + B + C + D + E

with the information for the sign of each variable stored in a logical flag ltao (true (T) status of which means positive whereas false (F) status means negative). The detection of the add and subtract operators is done by the global flag tao and the immediate information is passed to ltao to evaluate the sign of the current variable considered, while the status of tao may set to true or false depending on the subsequent add or subtract operators encountered.

The order of operation would be the same as before, but the result would be correct. The execution procedure is shown below.

ltao

A + B + C + D + E

A + B + C + T₁

$$T_1 = -D-E = -3-5 = -8$$

A + B + T₂

$$T_2 = +2-8 = -6$$

A + T₃

$$T_3 = -4-6 = -10$$

F

$$F = +3-10 = -7$$

Expressions with parentheses can be handled relatively easily with the following controls.

1. While the current status of tao is passed to ltao as before, the status of tao is set to

true status before the analysis of the bracketed expression.

2. After a similar analysis of the expression as before, the status of tao is restored by the instruction `tao := ltao`.

A similar control has been imposed on the division operator `/` for the reason that expressions of the form

$$A/B/C/D$$

will not be evaluated correctly without the control. Two flags, the global flag `tmo` and the local `ltmo` are similarly set up for the purpose. The flag `tmo` is initially set to true status whenever the operator `/` is encountered. Each time the information is passed to `ltmo` which determines the appropriate execution to be taken. Expressions with `+`, `-` and/or parentheses are similarly handled with the current status of `tmo` stored and restored via the local flag `ltmo`.

CHAPTER IV

MODIFICATIONS MADE ON THE TRANSLATOR

4.1. Introduction.

So far syntax and semantics tables have been given which will perform first order algebraic differentiation on an input expression with X, Y assumed to be independent and dependent variables respectively. This chapter discusses how, with little modification to the syntax and the translator, the syntax and semantics tables can be used to produce algebraic derivatives of any order (provided the length of the target string does not exceed the limit specified) and either the result can be in symbolic form or the numerical value of the result at assigned points can be evaluated through a routine VALUE. In the latter case the numeric values of all variables in a given input string must be supplied and will be handled by the input routine. The dependent and independent variables can be specified together with the input string and this information can be handled by the input routine which recognizes the dependent and independent variable.

4.2. Additional parameters for evaluation and order of differentiation.

In order to supply the information for the order of differentiation and numerical evaluation, two

parameters must be supplied.

The parameter for order of differentiation is an integer ORDER which can be any positive integer. A second parameter EVAL is used to indicate whether a numerical result (EVAL = 1) or simply the algebraic form of the result (EVAL = 0) is required. When a numerical result is required the values of the unknown variables should be given and the source string must be such that only the dependent variable is on the left side of the '=' symbol. These parameters are put to follow the phrase type number of each source string. Hence '27;1;0;' would mean the source string to be processed is of type 27, differentiation of first order is necessary and no evaluation for the result. The complete set of source statements is then terminated with '0;0;0;' instead of '0;' as discussed in section 2.8..

4.3. Internal Defined Syntax for dependent, independent variables and constants.

Since the independent and dependent variables are not necessarily X and Y, the phrase types <independent variable> , <dependent variable> , and <constant> in syntax table on page 36 must be deleted and handled by a special routine GETVARIABLE, while the phrase type <alphabets> defined to be any one of the upper case English letters is added to the syntax table. The variables chosen to be independent and dependent variables are provided immediately following the source string to be

differentiated. These variables are preceded by a '&' symbol which signifies that they are employed respectively as independent and dependent variables and the input routine will only take the character immediately following the mark '&' as the variable employed; and other characters preceding '&' or following the variable are ignored. An expression like $f(r,s) = e^{rs} \sin r \cos s$ when differentiation with respect to r is required is written in the input data as

```
27; 1; 0;
Y=EXP(R*S)*SIN(R)*COS(S);
INDEPENDENT VARIABLE IS &R DEPENDENT VARIABLE IS &Y;
```

The routine GETVARIABLE provides the internal codes for phrase types 11, 12, 13 which are equivalent to those specified by the syntax table on page 36, except that the independent and dependent variables may not be X and Y and the <constant> will consist of all the alphabets except the letters specified for independent and dependent variables. The GETVARIABLE routine listing is shown below.

```
procedure getvariable;
      begin integer i,j, indx ;
      boolean testor ;
      for i := 1 step 1 until 2 do
          begin for j := readch while j = 52 do ;
          var [i] := readch
          end ;
```



```

sysact (1, 14, 1) ;
outst ('( independent variable is ')');
printch (var [1] ) ;
outst ('(' dependent variable is ')') ;
printch (var [2] ) ;

testor := true
      indx := ilfsyn [13] := 7 ;
for i := 60 step 1 until 112 do
  begin if dsyn [i] = var [1] then
    begin ilfsyn [11] := 1 ;
    dsyn [1] := var [1] ; dsyn [2] := -12;
    dsyn [3] := -13;
    testor := false ;
    goto          endloop
    end ;
  if dsyn [i] = var [2] then
    begin ilfsyn [12] := 4 ;
    dsyn [4] := var [2] ; dsyn [5] := -12;
    dsyn [6] := -13;
    testor := false ;
    goto          endloop
  end ;
if testor then begin
  dsyn [indx] := dsyn [i] ;
  indx := indx + 1 ;
  end

```

```

        else testor := true ;
endloop : end ;
end getvariable ;

```

The starting points for the phrase types are indicated by the vector ILFSYN. The starting points for the phrase types concerned are shown as following

<u>phrase type number i</u>	<u>phrase type</u>	<u>ILFSYN[i]</u>
11	independent variable	1
12	dependent variable	4
13	constant	7
14	alphabets	60

4.4. Input routine for numeric data.

A routine VALUE has been designed to handle evaluation of derivatives and is discussed in section 3.3.. Here this input routine GETVALUES is discussed. The purpose of this routine is to read the numeric values for the necessary variables (including the numeric value for the independent variable). An array A [1 : 24] is reserved for storage of numeric values for the constants and is set to 0 initially. The real variable IX is also reserved for the numeric value of the independent variable. The input data for this routine is immediately following that specified for independent and dependent variables, and

it has the form

```
&variable 1 = numeric value ; &variable 2 = numeric
value ; ..... &variable n = numeric value ; #
```

where variable 1, variable n are variables consisting of one letter only. The mark & precedes the variable has the same purpose as discussed in the last section. The mark # following the last semi-colon indicates the end of input information for the numerical values of the variables in the input string. All the numeric values are assumed to be real floating point numbers and will be automatically assigned to the corresponding constants by the routine GETVALUES. The listing of the routine is shown below.

```
integer procedure readvalue ;
    begin integer i ;
start   : insymbol (0, ('ABCDEFHIJKLMNOPQRSTUVWXYZ
                &# '), i);
    if i = 29 then goto start ;
    readvalue := i
    end ;
procedure getvalues ;
    begin integer index, min, max ;
    if var [1] < var [2] then begin
        max := var [2] ; min := var [1]
    end
```

```

else begin
    max := var [1] ; min := var [2]
    end ;
    min := min - 14 ;
    max := max - 14 ;
reread : index := readvalue ;
    if index := 27 then goto readconst ;
    if index = 28 then goto end ;
    goto reread ;
readconst : index := readvalue ;
    r := read r ;
    sysact (1, 14, 1) ;
    outsymbol (1, ('ABCDEFGHIJKLMN O PQRSTU VWXYZ
    &# '),' , INDEX);
    oust ('(' = ')'); out r (R, true, 10, 3);
if index = var [1] - 14 then ix := r else
    A [if index < min then index
        else if index < max then index - 1
        else index - 2] := r ;
    goto reread ;
end : end getvalues ;

```

4.5. Modification on the main program for co-ordination of the various subroutines.

In order to implement the above subroutines, the main program of the translator must be slightly modified. As has been mentioned, ORDER and EVAL are parameters set up respectively for the order of differentiation and the code for evaluation. The variables LORDER and LEVAL are set up for similar purposes but will work logically inside the block structure of the for statement. A Boolean variable TESTEVAL is also set up and has initially the value false. It is set true when the EVAL has the value of 1. When the variable is true the subroutine GETVALUES is executed in addition to the routine VALUE. When the value of the parameter ORDER is greater than 1 or the code EVAL is 1 the previous target string produced is to be used as the source string for the next higher algebraic differentiation or evaluation, so that entire target string is copied back as a source string for further processing. This has also been done in the main program. The portion of the main program that has been modified is shown below.

```

          ·
          ·
          ·
readsyntax ;
readsemantics;
for i := 1 step 1 until 10 do
            sa [1] := i -1 ;

```

```

newss : readss (pt);
      testeval := false; lorder := order; leval := eval;
findval : for m := lorder step -1 until 1 do
  begin
    for i := 1 step 1 until 26 do A [ i ] := 0;
    for i := 1 step 1 until length do
      begin
        type [i] := cat [i] := from [i] := to [i] := 0;
        for j := 1 step 1 until nlinks do link [j,i] := 0
      end ;
    link [1, 0] := 1; rpoint := spoint := 1 ;
    nextout := 0 ;
    if testeval then
      begin
        tao := true ; tmo := true ;
        getvalues ;
      end
      else getvariable;
    if search (pt) then
      begin
        mi('(' analysed successfully as type ')', pt);
        sysact (1, 14, 1);
        outst ('(' tabular analysis record : ')');
        sysact (1, 14, 2);
        outst (
          '(' entry type cat from to links ')');
        sysact (1, 14, 2);
      end
    ;
  end

```

```

for i := 1 step 1 until rpoint -1 do
    begin outi (i, 10); outi (type [i] , 10);
outi (cat [i] , 10); outi (from [i] , 10);
outi (to [i] , 10);
for j := 1 step 1 until nlinks do
    outi (link [j, i] , 10);
    sysact (1, 14, 1)
    end ;
sysact (1, 14, 2); outst ('(' target string
    produced :')');
sysact (1, 14, 1);
if testeval then begin
    result := value (1, 0);
    out r (result, true, 15, 3);
    end
else begin
    compile (1, 0, 1); sysact (1, 14, 1);
target print ;
if nextout > = lss then
    begin
        fault (20, false); goto newss
    end ;
for i := 1 step 1 until nextout do
    source [i] := targ [i] ;
    nextin := nextout;
    insource (term);
    end
end

```

```
else
  begin
    mi ('(' absent as type ')', pt); sysact (1, 14, 1);
    outst ('(' no attempt at compilation possible ')');
    goto newss
  end ;
end ;

if leval = 1 then begin testeval := true ;
                                lorder := 1; leval := 0;
                                goto findval
                                end ;

  goto newss ;

ssfail : mi('('**** overlenght ss of type ')', pt);
          sysact (1, 14, 1);
```

•
•
•
•
•
•
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CHAPTER V
CONCLUSION

5.1. Introduction.

This work shows how a syntax-directed translator can be used for algebraic differentiation instead of its usual task of translating a high-level programming language into an assembly language. We now discuss advantages and disadvantages of this system and some suggestions of future works on it.

5.2. Advantages of this system.

The main advantage of using a syntax-directed translator for symbol manipulation is versatility. The same translator can be used for different symbol manipulation problems (e.g. the translation of an algebraic expression into its reverse polish notation form, to obtain a derivative of an algebraic expression, a programming language from another programming language). The only data to be supplied is the necessary syntax and semantic tables and the specimen source statement to be translated. Since the translator is already a tested program, there is no need to recompile or modify the logic of the program each time when a different translation is required, instead, effort is spent on designing a efficient syntax and semantic tables for a

particular translation. The fact that the writing and altering the syntax and semantic specification is a relatively easy task as compared to the approach when a program is written specifically for a particular translation makes this approach a feasible one. Furthermore the error-checking and self-recovering facilities built in the translator offer great assistance in providing the correct syntax and semantic specifications especially during test runs.

5.3. Disadvantages.

The main disadvantage for employing a top-down syntax-directed translator on symbolic manipulation is that it is relatively slow. This is especially so when the translator is written in ALGOL which is slow in IBM/360 in general. The main reason that it is slow as compared to some other methods is the time spent on searching. No matter how clever the syntax designer is, the analyzer will spend a certain amount of time in exploring the blind alleys. In fact it is found that the time spent on analyzing an expression and get the output increase exponentially with the length of the source string.

A second disadvantage is that the relatively large amount of memory space required for the translator. The translator employed here together with data areas

requires a memory space of approximately 180K. The reasons for this are that the translator itself is already a large program and requires a fairly amount of core. Furthermore, large amount of space must also be allocated to the arrays like DSYN, ILFSYN, DSEM, and ILFSEM to store internally the syntax and semantic tables. Memory space must also be given to the array SOURCE for the storage of the input string, the arrays TYPE, CAT, FROM, TO, LINK to store the analysis record and the array TARG for the resulting target string. Some of the main memory space could be save if some secondary storage devices is used for some of the informations mentioned above.

5.4. Suggestions for future work.

In this work, this method has been proved quite satisfactory as far as flexibility is concerned. The additional syntax rules and the corresponding semantic rules can be added quite easily whenever required. Also any alphabet can be used to stand for any constant or variable desired by simply supplying the information together with the source string in the input data. It can also do partial derivatives, derivatives of higher order and the derivatives of implicit functions. The evaluation of the resulting output can also be possible. Perhaps the algebraic differentiation can be made in the form of

a specific command (say DIFF). For example the third derivative of $y^3 - 4xy + 5$ with respect to x can be written as

DIFF (Y@3 - 4*X*Y + 5, X, 3)

Partial derivatives can also be found similarly

e.g.
$$\frac{\partial^3(x \sin(x-y))}{\partial x \partial y^2}$$

can be expressed as

DIFF (X*SIN (X-Y), X, 1, Y, 2)

The command for simplification can also be made similarly. For example, the above differentiation with simplification can be written as

SIM (DIFF (X*SIN (X-Y), X, 1, Y, 2)).

Similarly specific command can be made on evaluation (say EVAL). The value of $y \cos(x@3)$ can be found by this command when the values of x and y are given.

e.g. EVAL (Y*COS(X@3), X=4, Y=7)

The numeric result of the third derivative of $y^3 - 4xy + 3$ can be found similarly with specified values of x and y .

EVAL ((DIFF(Y@3 - 4*X*Y+3), X, 3), X=2, Y=5)

All the above specific commands are suggestions but there seem to be no reasons why they cannot be implemented.

This command language can be further extended so that they can be used in an on-line system. This is very desirable especially under a time-sharing environment which allows many users to have easy access to the computer simultaneously. When a user has gain access to the system the word 'MESSAGE;' will be displayed on his terminal. He can reply with any one of the following commands.

- (a) NEW.
- (b) SAME.
- (c) CONTINUE.
- (d) Any one of the following
 - DIFF.
 - EVAL.
 - SIM.
- (e) FINISH.

Each command in (d) corresponding to a set of syntax and semantics which has been previously feed in.

The functions of these commands can be explains as follows :

- (a) When the command NEW. is typed in, the system will responds with another
MESSAGE ;
and await a source statement.

The necessary source statement is then feed in accordingly.

- (b) Command SAME. is used when the current source statement is to be process the second time. Again the system will responds with

MESSAGE ;

and await next command instruction, which would, in fact, accept any of the above commands but presumably the user will reply with a command of type (d).

- (c) If the user wishes to process the compiled version of the current source statement, he can type in CONTINUE. , and the program will copy this target string to the source statement array and reply with

COPY BACK ;

MESSAGE ;

The user can then feed in the command DIFF. if he want to obtain a higher derivative, or SIM. , if he want to simplify the source statement, or simply type in EVAL. to obtain the numerical value of the source statement. As in (c), the user can in fact feed in any of the commands but the above three are the

one most like to be used.

- (d) When any one of these commands are fed in, the program will analyze the current source statement, and if the search is successful the message

SEARCH OK ;

is printed and start to assemble the target string via the COMPILE routine. The resulting target string is then printed out followed by the prompt

MESSAGE ;

If the search fails the message

SEARCH FAILS ;

is printed and followed by

MESSAGE ;

- (e) When the user wishes to terminate the program he can simply feeds in the command
FINISH.

This on-line system has been tried in the University of Nottingham and has proved quite successful. Although the terminal equipment they used was the KDF9 monitor flexowriter which is, in fact, slow and unsuitable for this purpose. In fact, instead of the normal keypunch as input and high-speed printer as output, if we can devise some good display equipment

for the purpose it would be highly desirable. The expression like $y = x^2 \exp\left(\frac{A}{x}\right)$ would be keypunched in the linear form like

$$Y = X@2*EXP(A/X)$$

and the derivative of this expression printed on the line-printer would be

$$Y' = (X@2*((2*1)/X+0*LN(X))*EXP(A/X) \\ + X@2*((0*X-A*1)/X@2)*EXP(A/X)}$$

which is almost incomprehensible without a lot of practice. If the special display device can accept the input as well as the resulting output expression in two-dimensional form like

$$Y = X^2 * EXP\left(\frac{A}{X}\right)$$

and

$$Y' = \left(X^2 * \left(\frac{2*1}{X} + 0 * LN(X) \right) * EXP\left(\frac{A}{X}\right) \right. \\ \left. + X^2 * \left(\frac{0*X - A*1}{X^2} \right) * EXP\left(\frac{A}{X}\right) \right)$$

the advantage of this is obvious.

APPENDIX A

LISTING FOR DIFFERENTIATION OF
RESTRICTED ARITHMETIC EXPRESSION

The following is a complete listing of the input data for the algebraic differentiation of restricted arithmetic expressions (e.g. $4 + X * Y$) and the complete output listing for $4 + X * Y$ is also shown.

```

09
<>()|E;&.+ =QU,WKLCPSNAZID
300;2;200;400;
+++
TEST DATA FOR ALGEBRAIC DIFFERENTIATION OF RESTRICTED ARITHMETIC EXPRESSIONS;
16,80,160,1;
11 <PRIMARY>      =<12>|<13>|<14>|<(><16><)>;
12 <INDEPT VARIABLE>=X;
13 <DEPENDENT VARIABLE>=Y;
14 <CONSTANT>     =0|1|2|3|4|5|6|7|8|9;
15 <TERM>         =<11>(*<15>|<E>);
16 <EXPRESSION>  =<15>(+<16>|<E>);
0;
1;
11 <PRIMARY>      =K<4>W<(>.
                  K<1,2,3,4>C<1>.
                  K<4>W<(>.;
12 <INDEPT VARIABLE>=W<1>;
13 <DEPENDENT VARIABLE>=PW<'>;
14 <CONSTANT>     =W<0>;
15 <TERM>         =K<1>W<(>C<1>W<*>L<2>
                  W<+>L<1>W<*>C<2>W<(>&
                  K<2>C<1>.;
16 <EXPRESSION>  =C<1>K<1>W<+>C<2>.;
0;
0;
16;
4+X*Y;
0;
+++
+++
END OF RUN;
0;

```

CODE FOR ZERO	1	
CODE FOR NINE	10	
CODE FOR OB	44	
CODE FOR CB	45	
CODE FOR ININ	46	
CODE FOR IN	-10	
CODE FOR OUTIN	47	
CODE FOR OUT	-11	
CODE FOR ORIN	51	
CODE FOR OR	-12	
CODE FOR EE	19	
CODE FOR TERMIN	50	
CODE FOR TERM	-13	
CODE FOR ELSEIN	52	
CODE FOR ELSE	-2	
CODE FOR ENDDIN	43	
CODE FOR ENDD	-3	
CODE FOR STAR	11	
CODE FOR EQUALS	49	
CODE FOR OQ	31	
CODE FOR CQ	35	
CODE FOR COMMA	42	
CODE FOR WW	37	
CODE FOR KK	25	
CODE FOR LL	26	
CODE FOR CC	17	
CODE FOR PP	30	
CODE FOR SS	33	
CODE FOR SP	0	
CODE FOR NN	28	
CODE FOR NL	-1	
CODE FOR AA	15	
CODE FOR ZZ	40	
CODE FOR II	23	
CODE FOR DD	18	
CODE FOR WTERM	-4	
CODE FOR KTERM	-9	
CODE FOR KST	-5	
MAX LENGTH RECORD	300	
MAX NUMBER LINKS	2	
MAX LENGTH SOURCE STRING		200
MAX LENGTH TARGET STRING		400

TEST DATA FOR ALGEBRAIC DIFFERENTIATION OF RESTRICTED ARITHMETIC EXPRESSIONS:

USER PARAMETERS FOR THIS DOCUMENT
 LENGTH OF ILIFFE VECTORS 16
 LENGTH OF SYNTAX VECTOR 80
 LENGTH OF SEMANTICS VECTOR 160
 NUMBER OF SEMANTIC SETS 1
 READSYNTAX ENTERED

SYNTAX OK TYPE 11
 SYNTAX OK TYPE 12
 SYNTAX OK TYPE 13
 SYNTAX OK TYPE 14
 SYNTAX OK TYPE 15
 SYNTAX OK TYPE 16

SYNTAX ENTERED STORES OCCUPIED 64

POINTERS FOR SYNTAX
 0 0 0 0 0 0 0 0 0 0 1 16
 19 22 43 54

CONTENTS OF SYNTAX
 -1 12 -12 -1 13 -12 -1 14 -12 46 -1 16
 47 -12 -13 38 -12 -13 39 -12 -13 1 -12 2
 -12 3 -12 4 -12 5 -12 6 -12 7 -12 8
 -12 9 -12 10 -12 -13 -1 11 -10 13 -1 15
 -12 -12 -11 -12 -13 -1 15 -10 11 -1 16 -12
 -12 -11 -12 -13 0 0 0 0 0 0 0 0
 0 0 0 0 0 0 0 0 0 0 0 0

STARTED READING SEMANTIC SET 1

SEMANTICS OK TYPE 11
 SEMANTICS OK TYPE 12
 SEMANTICS OK TYPE 13
 SEMANTICS OK TYPE 14
 SEMANTICS OK TYPE 15
 SEMANTICS OK TYPE 16

SEMANTICS ENTERED FOR SET NUMBER 1

STORES USED 90
 POINTERS FOR THIS SET:
 0 0 0 0 0 0 0 0 0 0 1 26
 30 35 39 76

CONTENTS OF THIS SEMANTIC SET
 -5 4 -9 37 46 -4 -3 -5 1 2 3 4
 -9 17 1 1 -3 -5 4 -9 37 47 -4 -3
 -13 37 2 -4 -13 30 37 41 -4 -13 37 1
 -4 -13 -5 1 -9 37 46 -4 17 1 1 37
 13 -4 26 2 37 11 -4 26 1 37 13 -4
 17 2 1 37 47 -4 -2 -5 2 -9 17 1
 1 -3 -13 17 1 1 -5 1 -9 37 11 -4
 17 2 1 -3 -13

READSEMANTICS COMPLETE TOTAL STORES OCCUPIED 89

NEW SOURCE STATEMENT
 TYPE 16

LENGTH 5

SOURCE STATEMENT :

4+X*Y

ANALYSED SUCCESSFULLY AS TYPE

16

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	16	1	1	5	2	5
2	15	2	1	1	3	0
3	11	3	1	1	4	0
4	14	5	1	1	0	0
5	16	2	3	5	6	0
6	15	1	3	5	7	9
7	11	1	3	3	8	0
8	12	1	3	3	0	0
9	15	2	5	5	10	0
10	11	2	5	5	11	0
11	13	1	5	5	0	0

TARGET STRING PRODUCED :

O+(1*Y+X*Y*)

****END OUTPUT THIS DOCUMENT

END OF RUN;

****END OUTPUT FOR THIS SET OF DOCUMENTS

APPENDIX B

A COMPLETE LISTING OF THE TRANSLATOR

A slightly modified version of King's translator is shown. This translator can perform algebraic differentiation of higher order, partial derivatives, derivatives of implicit functions and evaluation of the result.

```
'BEGIN'  
'COMMENT' THE FOLLOWING INPUT/OUTPUT ROUTINES HAVE BEEN ADDED TO THIS PROGRAM :
```

```
** INPUT/OUTPUT TRANSFER ROUTINES **
```

```
INPUT FUNCTION PROCEDURES : READR, READI, READB
```

```
OUTPUT ROUTINES :
```

```
BASIC PROCEDURES :
```

```
REAL FORMAT(BOOL, FW, D)   OUTREAL(CHANNEL, R)  
INTEGER FORMAT(FW)         OUTINTEGER(CHANNEL, I)  
BOOLEAN FORMAT(FW)         OUTBOOLEAN(CHANNEL, B)
```

```
FUNCTION PROCEDURES :
```

```
IOR(R, '('VAR NAME')', BOOL, FW, D)  
IOI(I, '('VAR NAME')', FW)  
IOB(B, '('VAR NAME')', FW)
```

```
PROCEDURES : IOA(A, L, U, '('ARRAY NAME')', BOOL, FW, D)
```

```
OTR(R, '('VAR NAME')', BOOL, FW, D)  
OTI(I, '('VAR NAME')', FW)  
OTB(B, '('VAR NAME')', FW)  
OTA(A, L, U, '('ARRAY NAME')', BOOL, FW, D)
```

```
OUTR(R, BOOL, FW, D), OUTI(I, FW), OUTB(B, FW)  
OUTA(A, L, U, BOOL, FW, D)  
OUTIA(IA, L, U, FW)  
OUTBA(BA, L, U, FW)
```

```
OUTST( STRING )
```

```
ECHO TEXT FROM INPUT TO OUTPUT DEVICE :  
COPYTEXT
```

```
** INPUT/OUTPUT CONTROL ROUTINES **
```

```
DATA SET CNTRL : OPEN(CHANNEL)   CLOSE(CHANNEL)  
                  LINESIZE(N)    PAGESIZE(N)
```

```
INPUT CNTRL : NEXTCARD
```

```
OUTPUT CNTRL : SPACE(N) TAB(N) NEWLINES(N) NEWPAGE ;
```

```

00339 *BEGIN
00339 *INTEGER*OB,CB,IN,OUT,ININ,OUTIN,OR,ORIN,TERM,TERMIN,
00339 ELSE,ELSEIN,ENDD,ENDDIN,STAR,EQUALS,OQ,CQ,COMMA,
00339 WW,KK,LL,CC,PP,SS,SP,NN,NL,AA,ZZ,ZERO,NINE,I,J,
00339 LILF,LDSYN,LDSEM,NS,NLINKS,LENGTH,RPOINT,SPOINT,
00339 WTERM,KTERM,KST,LSS,LTARG,NEXTIN,NEXTOUT,II,DD,EE,PT,EVAL,M,ORDER,
00339 LEVAL,LORDER;
00340 *REAL* RESULT,IX,R ;
00341 *BOOLEAN* TESTEVAL;
00342 *INTEGER* *PROCEDURE* READVALUE;
00343 *BEGIN* *INTEGER* I;
00344 START: INSYMBOL(0,(''ABCDEFGHIJKLMNPNRSTUVWXYZ&# '),'I);
00345 *IF* I=29 *THEN* *GOTO* START;
00346 READVALUE:=I
00346 *END*;
00347 *INTEGER* *PROCEDURE* READCH;
00348 *BEGIN*
00348 *INTEGER* I;
00349 START:
00349 INSYMBOL(0,(''0123456789+*/ABCDEFGHIJKLMNPNRSTUVWXYZ',.<>():=;|&-@ '),'I);
00349 *IF* I = 55 *THEN* *GOTO* START;
00350 READCH:=I
00351 *END*;
00352 *INTEGER* *PROCEDURE* READINT;
00353 *BEGIN* *INTEGER* I,J;
00354 *FOR* I:=READCH *WHILE* *NOT* DIGIT(I) *DO*;
00355 I:=I-ZERO;
00356 *FOR* J:=READCH *WHILE* DIGIT(J) *DO*
00356 I:=I*10 + J-ZERO;
00357 READINT:=I
00357 *END*;
00358 *PROCEDURE* PRINTCH(I); *VALUE* I; *INTEGER* I;
00361 *IF* I=-1 *THEN* SYSACT(1,14,1) *ELSE*
00361 OUTSYMBOL(1,(''0123456789+*/ABCDEFGHIJKLMNPNRSTUVWXYZ',.<>():=;|
00361 &-@'),'I);
00362 *PROCEDURE* MI(S,I); *VALUE* I; *INTEGER* I; *STRING* S;
00366 *BEGIN*
00366 SYSACT(1,14,1); OUTST(S); OUTI(I,10)
00368 *END*;
00369 *BOOLEAN* *PROCEDURE* LETTER(I); *VALUE* I; *INTEGER* I;
00372 LETTER:= AA<=I *AND* I<=ZZ;
00373 *BOOLEAN* *PROCEDURE* DIGIT(I); *VALUE* I; *INTEGER* I;
00376 DIGIT:=ZERO<=I *AND* I<=NINE ;
00377 OPEN(OUTCHANNEL);
00378 ZERO:=READCH; NINE:=READCH;
00380 OB:=READCH;
00381 CB:=READCH;

```



```
00382 ININ:=READCH; IN:=-10;
00384 OUTIN:=READCH; OUT:=-11;
00386 ORIN:=READCH; OR:=-12;
00388 EE:=READCH;
00389 TERMIN:=READCH; TERM:=-13;
00391 ELSEIN:=READCH; ELSE:=-2;
00393 ENDDIN:=READCH; ENDD:=-3;
00395 STAR:=READCH;
00396 EQUALS:=READCH;
00397 QQ:=READCH;
00398 CQ:=READCH;
00399 COMMA:=READCH;
00400 WW:=READCH;
00401 KK:=READCH;
00402 LL:=READCH;
00403 CC:=READCH;
00404 PP:=READCH;
00405 SS:=READCH; SP:=0;
00407 NN:=READCH; NL:=-1;
00409 AA:=READCH;
00410 ZZ:=READCH;
00411 II:=READCH;
00412 DD:=READCH;
00413 WTERM:=-4;
00414 KTERM:=-9;
00415 KST:=-5;
00416 LENGTH:=READINT;
00417 NLINKS:=READINT;
00418 LSS:=READINT;
00419 LTARG:=READINT;
00420 MI(('MAX LENGTH RECORD '),LENGTH);
00421 MI(('MAX NUMBER LINKS '),NLINKS);
00422 MI(('MAX LENGTH SOURCE STRING '),LSS);
00423 MI(('MAX LENGTH TARGET STRING '),LTARG);
00424 NEWDOC;
00424 NEWPAGE;
00425 'FOR' I:=READCH 'WHILE' I=STAR 'DO';
00426 PRINTCH(I);
00427 COPYTEXT;
00428 LILF:=READINT; 'IF' LILF=0 'THEN' 'GOTO' END;
00430 LDSYN:=READINT;
00431 LDSEM:=READINT;
00432 NS:=READINT;
00433 SYSACT(1,14,2);
00434 OUTST(('USER PARAMETERS FOR THIS DOCUMENT'));
00435 MI(('LENGTH OF ILIFFE VECTORS '),LILF);
00436 MI(('LENGTH OF SYNTAX VECTOR '),LDSYN);
00437 MI(('LENGTH OF SEMANTICS VECTOR '),LDSEM);
```

```

00438 MI('('NUMBER OF SEMANTIC SETS ')',NS);
00439   *BEGIN
00439     *INTEGER' ARRAY'
00439       ILFSYN(/1:LILF/), DSYN(/1:LDSYN/),
00439       ILFSEM(/1:NS,1:LILF/), DSEM(/1:LDSEM/),
00439       SOURCE(/1:LSS/), TARG(/1:LTARG/),
00439       TYPE,CAT,FROM,TO(/1:LENGTH/),
00439       LINK(/1:NLINKS,0:LENGTH/),SA(/1:10/),VAR(/1:2/);
00440     *REAL' ARRAY'
00440       A(/1:26/);
00441     *BOOLEAN' TAO,TMO;
00442   *PROCEDURE' GETVARIABLE;
00443     *BEGIN' *INTEGER' I,J,INDX;
00444     *BOOLEAN' TESTOR;
00445     *FOR' I:=1 *STEP' 1 *UNTIL' 2 *DO
00445       *BEGIN' *FOR' J:=READCH *WHILE' J<=52 *DO';
00446       VAR(/1/):=READCH
00446       *END';
00447       SYSACT(1,14,1);
00448       OUTST({' INDEPENDENT VARIABLE IS '});
00449       PRINTCH(VAR(/1/));
00450       OUTST({' DEPENDENT VARIABLE IS '});
00451       PRINTCH(VAR(/2/));
00452   TESTOR:='TRUE';
00453     INDX:=ILFSYN(/13/):=7;
00454     *FOR' I:=60 *STEP' 1 *UNTIL' 112 *DO
00454       *BEGIN' *IF' DSYN(/1/)=VAR(/1/) *THEN
00454         *BEGIN' ILFSYN(/11/):=1;
00455         DSYN(/1/):=VAR(/1/); DSYN(/2/):=-12; DSYN(/3/):=-13;
00458         TESTOR:='FALSE';
00459         *GOTO' ENDOLOOP
00459         *END';
00460         *IF' DSYN(/1/)=VAR(/2/) *THEN
00460         *BEGIN' ILFSYN(/12/):=4;
00461         DSYN(/4/):=VAR(/2/); DSYN(/5/):=-12; DSYN(/6/):=-13;
00464         TESTOR:='FALSE';
00465         *GOTO' ENDOLOOP
00465         *END';
00466         *IF' TESTOR *THEN' *BEGIN'
00466         DSYN(/INDX/):=DSYN(/1/);
00467         INDX:=INDX+1;
00468         *END'
00468       *ELSE' TESTOR:='TRUE';
00469     ENDOLOOP: *END';
00470     *END' GETVARIABLE;
00471   *PROCEDURE' GETVALUES;
00472     *BEGIN' *INTEGER' INDEX,MIN,MAX;
00473     *IF' VAR(/1/)<VAR(/2/) *THEN' *BEGIN'

```

```
00473         MAX:=VAR(/2/); MIN:=VAR(/1/)
00474         'END'
00474         'ELSE' 'BEGIN'
00474         MAX:=VAR(/1/); MIN:=VAR(/2/)
00475         'END';
00476         MIN:=MIN-14;
00477         MAX:=MAX-14;
00478 REREAD: INDEX:=READVALUE;
00479         'IF' INDEX=27 'THEN' 'GOTO' READCONST;
00480         'IF' INDEX=28 'THEN' 'GOTO' END;
00481         'GOTO' REREAD;
00482 READCONST: INDEX:=READVALUE;
00483         R:=READ R;
00484         SYSACT(1,14,1);
00485         OUTSYMBOL(1,('ABCDEFGHIJKLMNPOQRSTUVWXYZ&# '),INDEX);
00486         OUTST('(' = ')') ; OUT R(R,'TRUE',10,3);
00488         'IF' INDEX=VAR(/1/)-14 'THEN' IX:=R 'ELSE'
00488         A(/'IF' INDEX<MIN 'THEN' INDEX
00488         'ELSE' 'IF' INDEX<MAX 'THEN' INDEX-1
00488         'ELSE' INDEX-2/):=R;
00489         'GOTO' REREAD;
00490 END: 'END' GETVALUES;
00491 'PROCEDURE' FAULT(J,CAT); 'VALUE' J,CAT; 'BOOLEAN' CAT; 'INTEGER' J;
00495         'BEGIN'
00495         'IF' CAT 'THEN' SYSACT(1,14,4);
00496         MI('('*****FAILURE NUMBER ')',J);
00497         SYSACT(1,14,1);
00498         'IF' CAT 'THEN'
00498         'BEGIN'
00498         OUTST('('*****CATASTROPHIC FAILURE')');
00499         SYSACT(1,14,1);
00500         OUTST('('*****REMAINDER OF DOCUMENT IGNORED')');
00501         'FOR' I:=READCH 'WHILE'
00501         'NOT' (I=STAR 'AND' READCH=STAR) 'DO' ;
00502         'GOTO' NEWDOC
00502         'END';
00503         OUTST('('*****WARNING MESSAGE      DOCUMENT RESUMED')');
00504         SYSACT(1,14,1)
00504         'END';
00505 'PROCEDURE' READSYNTAX;
00506         'BEGIN' 'INTEGER' I,DP,TYPE,J,JJ;
00507         'PROCEDURE' ENTER(I); 'VALUE' I; 'INTEGER' I;
00510         'BEGIN'
00510         'IF' DP>LDSYN 'THEN' FAULT(1,'TRUE');
00511         DSYN(/DP/):=I; DP:=DP+1
00512         'END';
00513         'FOR' I:=1 'STEP' 1 'UNTIL' LILF 'DO' ILFSYN(/I/):=0;
00514         'FOR' I:=1 'STEP' 1 'UNTIL' LDSYN 'DO'
```

```

00514          DSYN(/I/):=0;
00515          DP:=60;
00516          SYSACT(1,14,1);
00517          OUTST(('READSYNTAX ENTERED'));
00518          SYSACT(1,14,1);
00519          'FOR' TYPE:=READINT 'WHILE' TYPE<=0 'DO'
00519          'IF' TYPE<=10 'OR' TYPE>LILF 'THEN' FAULT(2,'TRUE')
00519          'ELSE'
00519          'BEGIN'
00519          'IF' ILFSYN(/TYPE/) <= 0 'THEN' FAULT(3,'FALSE');
00520          ILFSYN(/TYPE/):=DP;
00521          'FOR' I:=READCH 'WHILE' I<=EQUALS 'DO';
00522          'FOR' I:=READCH 'WHILE' I<=TERMIN 'DO'
00522          'BEGIN'
00522          'IF' I=ORIN 'THEN' ENTER(OR)
00522          'ELSE' 'IF' I=ININ 'THEN' ENTER(IN)
00522          'ELSE' 'IF' I=OUTIN 'THEN'
00522          'BEGIN'
00522          ENTER(OR); ENTER(OUT)
00523          'END'
00523          'ELSE' 'IF' I=OB 'THEN'
00523          'BEGIN'
00523          I:=READCH;
00524          'IF' DIGIT(I) 'THEN'
00524          'BEGIN'
00524          ENTER(-1); I:=I-ZERO;
00526          'FOR' J:=READCH 'WHILE' DIGIT(J) 'DO'
00526          I:=I * 10 + J - ZERO;
00527          ENTER(I);
00528          'IF' J <= CB 'THEN' FAULT(4,'TRUE')
00528          'END'
00528          'ELSE' 'BEGIN'
00528          'IF' READCH <= CB 'THEN' FAULT(5,'TRUE');
00529          'IF' I=LL 'OR' I=DD 'OR' I=II 'OR' I=NN
00529          'OR' I=SS 'THEN' ENTER(-1);
00530          'IF' I <= EE 'THEN'
00530          ENTER('IF' I=LL 'THEN' 1
00530          'ELSE' 'IF' I=DD 'THEN' 2
00530          'ELSE' 'IF' I=II 'THEN' 4
00530          'ELSE' 'IF' I=NN 'THEN' 6
00530          'ELSE' 'IF' I=SS 'THEN' 8
00530          'ELSE' I)
00530          'END'
00530          'END'
00530          'ELSE' ENTER(I)
00530          'END';
00531          ENTER(OR); ENTER(TERM);
00533          MI(('SYNTAX OK TYPE'),TYPE);

```

```

SC      SOURCE STATEMENT
00534          'END';
00535      SYSACT(1,14,2);
00536      MI(('SYNTAX ENTERED  STORES OCCUPIED '),DP-1);
00537      SYSACT(1,14,1);
00538      OUTST(('POINTERS FOR SYNTAX'));
00539      SYSACT(1,14,1);
00540      'FOR' I:=1 'STEP' 1 'UNTIL' LILF 'DO' OUTI(ILFSYN(/I/),10);
00541      SYSACT(1,14,1);
00542      OUTST(('CONTENTS OF SYNTAX'));
00543      SYSACT(1,14,1);
00544      'FOR' I:=1 'STEP' 1 'UNTIL' LDSYN 'DO' OUTI(DSYN(/I/),10);
00545      SYSACT(1,14,2);
00546      'END' READSYNTAX ;
00547      'BOOLEAN' 'PROCEDURE' SEARCH(I); 'VALUE' I; 'INTEGER' I;
00550          'BEGIN'
00550          'INTEGER' THIS,DP,D,CATEG,DCATEG,T,SSTART,
00550          LP,DRSTART,DSSTART,DLPOINT;
00551          'BOOLEAN' DEEP ;
00552          DEEP:='FALSE';
00553          SSTART:=SPPOINT; THIS:=RPOINT;
00555          'IF' I<=10 'THEN' 'GOTO' BIPS;
00556          'IF' I>LILF 'THEN' FAULT(6,'TRUE');
00557          RPOINT:=RPOINT+1; DP:=ILFSYN(/I/);
00559          'IF' DP=0 'THEN'
00559              'BEGIN'
00559              SEARCH:='FALSE';
00560              FAULT(7,'FALSE');
00561              'GOTO' END
00561              'END';
00562          LP:=DCATEG:=CATEG:=1;
00563      NEXT:
00563          D:=DSYN(/DP/);
00564          'IF' D=OR 'THEN'
00564              'BEGIN'
00564              'IF' DEEP 'THEN'
00564                  'BEGIN'
00564                  'FOR' D:=DSYN(/DP/) 'WHILE' D=>OUT 'DO'
00564                      'BEGIN'
00564                      'IF' D=OR 'THEN' DCATEG:=DCATEG+1;
00565                      DP:=DP+1
00565                      'END';
00566                  DCATEG:=DCATEG-1;
00567                  DP:=DP+1;
00568                  DEEP:='FALSE';
00569                  'GOTO' NEXT
00569                  'END';
00570      SUCCESS: SEARCH:='TRUE';
00571          TYPE(/THIS/):=I;          CAT(/THIS/):=CATEG;

```

```

00573          FROM(/THIS/):=SSTART; TO(/THIS/):=SPOINT-1;
00575          'GOTO' END
00576          'END';
00576          'IF' D=TERM 'THEN'
00576          'BEGIN'
00576          SEARCH:='FALSE';
00577          RPOINT:=THIS;
00578          'GOTO' END
00578          'END';
00579          'IF' D=IN 'THEN'
00579          'BEGIN'
00579          DEEP:='TRUE'; DSSTART:=SPOINT;
00581          DRSTART:=RPOINT; DP:=DP+1;
00583          DLPOINT:=LP; 'GOTO' NEXT
00584          'END';
00585          'IF' D=OUT 'THEN'
00585          'BEGIN'
00585          CATEG:=CATEG-1;
00586          DEEP:='FALSE';
00587          'GOTO' TRYNEXTCATEGORY
00587          'END';
00588          'IF' D=-1 'THEN'
00588          'BEGIN'
00588          DP:=DP+1; T:=DSYN(/DP/);
00590          LINK(/LP,THIS/):=RPOINT;
00591          'IF' SEARCH(T) 'THEN'
00591          'BEGIN'
00591          LP:=LP+1; DP:=DP+1;
00593          'GOTO' NEXT
00593          'END';
00594          'GOTO' TRYNEXTCATEGORY
00594          'END';
00595          'IF' D=OUTSOURCE(SPOINT) 'THEN'
00595          'BEGIN'
00595          SPOINT:=SPOINT+1; DP:=DP+1;
00597          'GOTO' NEXT
00597          'END';
00598          TRYNEXTCATEGORY:
00598          'IF' DEEP 'THEN'
00598          'BEGIN'
00598          'FOR' D:=DSYN(/DP/) 'WHILE' D<=OR 'DO' DP:=DP+1;
00599          DP:=DP+1; CATEG:=CATEG+1;
00601          LP:=DLPOINT;
00602          SPOINT:=DSSTART;
00603          RPOINT:=DRSTART;
00604          'GOTO' NEXT
00604          'END';
00605          'FOR' DP:=DP+1 'WHILE' DSYN(/DP/)<=OR 'DO'

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```

00605      'IF' DSYN(/DP/)=IN 'THEN'
00605      'BEGIN'
00605      'FOR' DP:=DP+1 'WHILE' DSYN(/DP/) != OUT 'DO'
00605      'IF' DSYN(/DP/)=OR 'THEN' CATEG:=CATEG+1;
00606      CATEG:=CATEG-1
00606      'END';
00607      DP:=DP+1;
00608      CATEG:=( 'IF' CATEG>DCATEG 'THEN' CATEG 'ELSE' DCATEG)+1;
00609      RPOINT:=THIS+1; SPOINT:=SSTART;
00611      LP:=1; 'GOTO' NEXT;
00613  BIPS:
00613      'BEGIN'
00613      'SWITCH' S:=LETT, DIGIT, L, IDENT, L, INTEGER, L, STRING, L, L;
00614      D:=OUTSOURCE(SPOINT);
00615      'GOTO' S(/I/);
00616      L: FAULT(22, 'FALSE');
00617      SEARCH:='FALSE';
00618      'GOTO' END;
00619      LETTR: 'IF' LETTER(D) 'THEN' 'GOTO' BOTH;
00620      SEARCH:='FALSE'; 'GOTO' END;
00622  DIGITT: 'IF' DIGIT(D) 'THEN'
00622      BOTH: 'BEGIN'
00622      SPOINT:=SPOINT+1;
00623      CATEG:=D-( 'IF' I=2 'THEN' ZERO 'ELSE' AA);
00624      RPOINT:=RPOINT+1;
00625      'GOTO' SUCCESS
00625      'END';
00626      SEARCH:='FALSE'; 'GOTO' END;
00628  IDENT: 'IF' LETTER(D) 'THEN'
00628      'BEGIN' 'INTEGER' I;
00629      SPOINT:=SPOINT+1;
00630      CATEG :=D;
00631      I:=1;
00632      'FOR' D:=OUTSOURCE(SPOINT) 'WHILE' LETTER(D)
00632      'OR' DIGIT(D) 'DO'
00632      'BEGIN'
00632      SPOINT:=SPOINT+1;
00633      'IF' I<=6 'THEN' CATEG :=CATEG *ZZ+D;
00634      I:=I+1
00634      'END';
00635      RPOINT:=RPOINT+1;
00636      'GOTO' SUCCESS
00636      'END';
00637      SEARCH:='FALSE'; 'GOTO' END;
00639  INTEGER: 'IF' DIGIT(D) 'THEN'
00639      'BEGIN'
00639      SPOINT:=SPOINT+1; CATEG :=D-ZERO;
00641      'FOR' D:=OUTSOURCE(SPOINT) 'WHILE' DIGIT(D) 'DO'

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SC SOURCE STATEMENT

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00641          'BEGIN'
00641          SPOINT:=SPOINT+1;
00642          CATEG:=10*CATEG + D - ZERO
00642          'END';
00643          RPOINT:=RPOINT+1; 'GOTO' SUCCESS
00644          'END';
00645          SEARCH:='FALSE'; 'GOTO' END;
00647  STRING: 'IF' D=00 'THEN'
00647          'BEGIN' 'INTEGER' I,J;
00648          I:=1; SPOINT:=SPOINT+1; CATEG:=OUTSOURCE(SPOINT);
00651          'FOR' SPOINT:=SPOINT,SPOINT+1 'WHILE' I>0 'DO'
00651          'BEGIN' J:=OUTSOURCE(SPOINT);
00652          'IF' J=DQ 'THEN' I:=I+1 'ELSE' 'IF' J=CQ 'THEN'
00652          I:=I-1; J:=SPOINT
00653          'END';
00654          SPOINT:=J;
00655          RPOINT:=RPOINT+1;
00656          'GOTO' SUCCESS
00656          'END';
00657          SEARCH:='FALSE';
00658          'GOTO' END
00658          'END' BIPS;
00659  END: 'END' SEARCH ;
00660  'PROCEDURE' READSEMANTICS ;
00661          'BEGIN'
00661          'INTEGER' N,DP,T,SSN,I,J,X;
00662          'BOOLEAN' KF,ELF;
00663          'PROCEDURE' IGNORE TO(Z); 'VALUE' Z; 'INTEGER' Z;
00666          'BEGIN' 'INTEGER' I;
00667          'FOR' I:=READCH 'WHILE' I≠Z 'DO'
00667          'END';
00668          'PROCEDURE' ENTER(I); 'VALUE' I; 'INTEGER' I;
00671          'BEGIN'
00671          'IF' DP>LDSEM 'THEN' FAULT(8,'TRUE');
00672          DSEM(/DP/):=I; DP:=DP+1
00673          'END';
00674          'PROCEDURE' NUMBER(I,J); 'INTEGER' I,J;
00676          'BEGIN'
00676          I:=READCH - ZERO;
00677          'FOR' J:=READCH 'WHILE' DIGIT(J) 'DO' I:=I*10 + J - ZERO
00677          'END';
00678          'FOR' I:=1 'STEP' 1 'UNTIL' NS 'DO'
00678          'FOR' J:=1 'STEP' 1 'UNTIL' LILF 'DO' ILFSEM(/I,J/):=0;
00679          'FOR' I:=1 'STEP' 1 'UNTIL' LDSEM 'DO'
00679          DSEM(/I/):=0;
00680          DP:=1;
00681  START: SSN:=READINT; N:=DP-1;
00683          'IF' SSN=0 'THEN' 'GOTO' END;

```



```

00684 MI(('STARTED READING SEMANTIC SET'),SSN);
00685 'FOR' T:=READINT 'WHILE' T≠0 'DO'
00685 'BEGIN'
00685 'IF' T>LILF 'THEN' FAULT(23,'TRUE');
00686 'IF' ILFSEM(/SSN,T/)≠0 'THEN' FAULT(24,'FALSE');
00687 ILFSEM(/SSN,T/):=DP;
00688 KF:=ELF:='FALSE';
00689 IGNORETO(EQUALS);
00690 'FOR' X:=READCH 'WHILE' X≠TERMIN 'DO'
00690 'BEGIN'
00690 'IF' X=KK 'THEN'
00690 'BEGIN'
00690 'IF' KF 'THEN' FAULT(9,'TRUE');
00691 KF:='TRUE'; IGNORETO(OB);
00693 ENTER(KST); ELF:='FALSE';
00695 EKP: NUMBER(I,J); ENTER(I);
00697 'IF' J=COMMA 'THEN' 'GOTO' EKP;
00698 'IF' J=CB 'THEN' FAULT(10,'TRUE');
00699 ENTER(KTERM)
00699 'END'
00699 'ELSE' 'BEGIN'
00699 'IF' ELF 'THEN' FAULT(11,'TRUE');
00700 'IF' X=ELSEIN 'THEN'
00700 'BEGIN'
00700 'IF' 'NOT' KF 'THEN' FAULT(12,'TRUE');
00701 KF:='FALSE'; ELF:='TRUE';
00703 ENTER(ELSE)
00703 'END'
00703 'ELSE' 'IF' X=ENDDIN 'THEN'
00703 'BEGIN'
00703 'IF' 'NOT' KF 'THEN' FAULT(13,'TRUE');
00704 KF:='FALSE'; ENTER(ENDD)
00705 'END'
00705 'ELSE' 'IF' X=CC 'THEN'
00705 'BEGIN'
00705 IGNORETO(OB); ENTER(CC);
00707 NUMBER(I,J); ENTER(I);
00709 'IF' J=COMMA 'THEN' NUMBER(I,J)
00709 'ELSE' I:=SSN;
00710 'IF' J=CB 'THEN' FAULT(14,'TRUE');
00711 ENTER(I)
00711 'END'
00711 'ELSE' 'IF' X=LL 'THEN'
00711 'BEGIN'
00711 IGNORETO(OB); ENTER(LL);
00713 NUMBER(I,J); ENTER(I);
00715 'IF' J=CB 'THEN' FAULT(15,'TRUE')
00715 'END'

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```

00715          *ELSE* *IF* X=PP      *THEN* ENTER(X)
00715          *ELSE* *IF* X=WW      *THEN*
00715              *BEGIN*
00715                  IGNORETO(OB); ENTER(WW);
00717                  *FOR* X:=READCH *WHILE* X~CB *DO*
00717                      *BEGIN*
00717                          *IF* X=OB *THEN*
00717                              *BEGIN*
00717                                  X:=READCH;
00718                                  *IF* X=NN *THEN* X:=NL
00718                                  *ELSE* *IF* X=SS *THEN* X:=SP;
00719                                  *IF* READCH~CB *THEN*
00719                                      *FAULT*(16,'TRUE')
00719                              *END*;
00720                                  ENTER(X)
00720                              *END*;
00721                                  ENTER(WTERM)
00721                              *END*
00721                          *ELSE* *FAULT*(17,'TRUE')
00721                      *END*
00721                  *END*;
00722          *IF* KF *THEN*
00722              *BEGIN*
00722                  ENTER(ENDD); *FAULT*(18,'FALSE')
00723              *END*;
00724          ENTER(TERM);
00725          MI(*('SEMANTICS OK TYPE'),'T');
00726              *END*;
00727          SYSACT(1,14,2);
00728          MI(*('SEMANTICS ENTERED FOR SET NUMBER'),'SSN');
00729              SYSACT(1,14,1);
00730          MI(*('STORES USED'),'DP-N');
00731              SYSACT(1,14,1);
00732          OUTST(*('POINTERS FOR THIS SET: '));
00733              SYSACT(1,14,1);
00734          *FOR* I:=1 *STEP* 1 *UNTIL* LILF *DO*
00734              OUTI(ILFSEM(/SSN,I/),10);
00735              SYSACT(1,14,1);
00736          OUTST(*('CONTENTS OF THIS SEMANTIC SET'));
00737              SYSACT(1,14,1);
00738          *FOR* I:=N+1 *STEP* 1 *UNTIL* DP-1 *DO*
00738              OUTI(DSEM(/I/),10);
00739          *GOTO* START;
00740          END: SYSACT(1,14,2);
00741          MI(*('READSEMANTICS COMPLETE TOTAL STORES OCCUPIED'),'DP-1');
00742              SYSACT(1,14,2);
00743          *END* READSEMANTICS ;
00744          *PROCEDURE* COPY(I); *VALUE* I; *INTEGER* I;

```

```

00747      *BEGIN*  *INTEGER* P,Q,R;
00748      P:=FROM(/I/);  Q:=TO(/I/);
00750      *FOR* R:=P *STEP* 1 *UNTIL* Q *DO* INTARG(OUTSOURCE(R))
00750      *END*;
00751      *PROCEDURE* COPYLINK(I,J); *VALUE* I,J; *INTEGER* I,J;
00754      COPY(LINK(/I,J/));
00755      *PROCEDURE* COMPILE(I,J,P); *VALUE* I,J,P; *INTEGER* I,J,P;
00758      *BEGIN*
00758      *INTEGER* K,DP,D;
00759      *BOOLEAN* TC;
00760      J:=LINK(/I,J/); I:=TYPE(/J/); K:=CAT(/J/);
00763      TC:='TRUE';
00764      DP:=ILFSEM(/P,I/);
00765      *IF* DP=0 *THEN*
00765      *BEGIN*
00765      FAULT(19,'FALSE'); *GOTO* END
00766      *END*;
00767      NEXT: *IF* *NOT* TC *THEN*
00767      *BEGIN*
00767      *FOR* D:=DSEM(/DP/) *WHILE* D $\neq$ ELSE
00767      & D $\neq$ ENDD *DO* DP:=DP+1;
00768      *END*;
00769      D:=DSEM(/DP/); *IF* D=TERM *THEN* *GOTO* END;
00771      *IF* D=ELSE *THEN*
00771      *BEGIN*
00771      DP:=DP+1;
00772      *IF* TC *THEN*
00772      *BEGIN*
00772      *FOR* D:=DSEM(/DP/) *WHILE* D $\neq$ ENDD *DO* DP:=DP+1
00772      *END*
00772      *ELSE* TC:='TRUE'
00772      *END*
00772      *ELSE* *IF* D=ENDD *THEN*
00772      *BEGIN*
00772      TC:='TRUE'; DP:=DP+1
00773      *END*
00773      *ELSE* *IF* D=KST *THEN*
00773      *BEGIN*
00773      TC:='FALSE'; DP:=DP+1;
00775      *FOR* D:=DSEM(/DP/) *WHILE* D $\neq$ KTERM *DO*
00775      *BEGIN*
00775      *IF* D=K *THEN* TC:='TRUE';
00776      DP:=DP+1
00776      *END*;
00777      DP:=DP+1
00777      *END*
00777      *ELSE* *IF* D=PP *THEN*
00777      *BEGIN*

```

```

00777         DP:=DP+1; COPY(J)
00778         'END'
00778         'ELSE' 'IF' D=LL 'THEN'
00778         'BEGIN'
00778         COPYLINK(DSEM(/DP+1/),J);
00779         DP:=DP+2
00779         'END'
00779         'ELSE' 'IF' D=CC 'THEN'
00779         'BEGIN'
00779         COMPIL(DSEM(/DP+1/),J,DSEM(/DP+2/));
00780         DP:=DP+3
00780         'END'
00780         'ELSE' 'IF' D=WW 'THEN'
00780         'BEGIN'
00780         DP:=DP+1;
00781         'FOR' D:=DSEM(/DP/) 'WHILE' D-=-WTERM 'DO'
00781         'BEGIN'
00781         INTARG(D);
00782         DP:=DP+1
00782         'END';
00783         DP:=DP+1
00783         'END';
00784         'GOTO' NEXT;
00785         END: 'END' COMPIL ;
00786         'REAL' 'PROCEDURE' VALUE(I,J); 'VALUE' I,J; 'INTEGER' I,J;
00789         'BEGIN' 'INTEGER' K; 'REAL' L,Y;
00791         'SWITCH' S:=INDEPENDENT VARIABLE,DEP VARIABLE,CONSTANT,ALPHABETS,
00791         DUMMY1,QUOTE,DIGITS,INTEGERS,FRACTION PART,UNSIGNED DECIMAL NO,
00791         PRIMARY,SIG PRIMARY,FACTOR,TERM1,TERM2,SAE;
00792         'SWITCH' SS:=STATEMENT,FUNCTION NAME;
00793         'BOOLEAN' LTAO,LTMO;
00794         LTAO:=TAO;
00795         LTMO:=TMO;
00796         J:=LINK(/I,J/); I:=TYPE(/J/); K:=CAT(/J/);
00799         'IF' I>26 'THEN' 'GOTO' SS(/I-26/);         'GOTO' S(/I-10/);
00801         INDEPENDENT VARIABLE:
00801         VALUE:=IX; 'GOTO' END;
00803         DEP VARIABLE:
00803         PRINTCH(VAR(/2/)); 'GOTO' END;
00805         CONSTANT: VALUE:=A(/K/); 'GOTO' END;
00807         ALPHABETS: 'GOTO' END;
00808         DUMMY1: 'GOTO' END;
00809         QUOTE:
00809         QUTSYMBOL(1,(')'),1);
00810         'IF' K=1 'THEN' VALUE(1,J);
00811         'GOTO' END;
00812         DIGITS: VALUE:=SA(/K/);
00813         'GOTO' END;

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```

00814 INTEGERS: 'IF' K=1 'THEN' VALUE:=10*VALUE(1,J)+VALUE(2,J)
00814 'ELSE' VALUE:=VALUE(1,J);
00815 'GOTO' END;
00816 FRACTION PART:
00816 'IF' K=1 'THEN' VALUE:=VALUE(1,J)/10+VALUE(2,J)/100
00816 'ELSE' VALUE:=VALUE(1,J)/10;
00817 'GOTO' END;
00818 UNSIGNED DECIMAL NO:
00818 'IF' K=1 'THEN' VALUE:=VALUE(1,J)+VALUE(2,J)
00818 'ELSE' VALUE:=VALUE(1,J);
00819 'GOTO' END;
00820 PRIMARY: 'IF' K=1 'THEN' 'BEGIN'
00820 TAO:='TRUE';
00821 TMO:='TRUE';
00822 VALUE:=VALUE(1,J);
00823 TAO:=LTAO; TMO:=LTMO
00824 'END'
00824 'ELSE' VALUE:=VALUE(1,J);
00825 'IF' K=6 'THEN' VALUE(2,J);
00826 'GOTO' END;
00827 SIG PRIMARY:
00827 'IF' K=2 'THEN' VALUE:=-VALUE(1,J)
00827 'ELSE' VALUE:=VALUE(1,J);
00828 'GOTO' END;
00829 FACTOR: 'IF' K=1 'THEN' 'BEGIN' L:=VALUE(2,J);
00830 VALUE:=VALUE(1,J)**L
00830 'END'
00830 'ELSE' VALUE:=VALUE(1,J);
00831 'GOTO' END;
00832 TERM1:
00832 'IF' K=1 'THEN' 'BEGIN'
00832 'IF' LTMO 'THEN' 'BEGIN'
00832 Y:=VALUE(1,J);
00833 TMO:='FALSE';
00834 VALUE:=Y/VALUE(2,J)
00834 'END'
00834 'ELSE' VALUE:=VALUE(1,J)*VALUE(2,J)
00834 'END'
00834 'ELSE' VALUE:=VALUE(1,J);
00835 'GOTO' END;
00836 TERM2:
00836 'IF' K=1 'THEN' 'BEGIN'
00836 Y:=VALUE(1,J);
00837 TMO:='TRUE';
00838 VALUE:=Y*VALUE(2,J)
00838 'END'
00838 'ELSE' VALUE:=VALUE(1,J);
00839 'GOTO' END;

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SC      SOURCE STATEMENT
00840   SAE:      'IF' K=1 'THEN'
00840           'BEGIN'
00840           TMO:='TRUE';
00841           Y:=SIGN('IF' LTAO 'THEN' 1 'ELSE' -1)*VALUE(1,J);
00842           TAO:='TRUE' ;
00843           TMO:='TRUE';
00844           VALUE:=Y+VALUE(2,J);
00845           TMO:=LTMO
00845           'END'
00845           'ELSE' 'IF' K=2 'THEN'
00845           'BEGIN'
00845           TMO:='TRUE';
00846           Y:=SIGN('IF' LTAO 'THEN' 1 'ELSE' -1)*VALUE(1,J);
00847           TAO:='FALSE';
00848           TMO:='TRUE';
00849           VALUE:=Y+VALUE(2,J);
00850           TMO:=LTMO
00850           'END'
00850           'ELSE' VALUE:=SIGN('IF' LTAO 'THEN' 1 'ELSE'-1)*VALUE(1,J);
00851           'GOTO' END;
00852   STATEMENT:
00852           VALUE(1,J); OUTST('(*)'); VALUE:=VALUE(2,J);
00855           'GOTO' END;
00856   FUNCTION NAME:
00856           'BEGIN'
00856           'REAL' 'PROCEDURE' EXP1(X); 'REAL' X;
00858           EXP1:=EXP(X)-EXP(-X);
00859           'REAL' 'PROCEDURE' EXP2(X) ; 'REAL' X ;
00861           EXP2:=EXP(X)+EXP(-X);
00862           'SWITCH' S:=AEXP,ALN,SINH,COSH,TANH,COTH,SECH,CSCH,
00862   ARCSIN,ARCCOS,AARCTAN,ARCCOT,ARCSEC,ARCCSC,ASIN,ACOS;
00863           'SWITCH' SS:=TAN,COT,SEC,CSC,ARSINH,ARCOSH,ARTANH,ARCOth,ARSECH,
00863   ARCSCH;
00864           'REAL' X,PI;
00865           TAO:='TRUE'; TMO:='TRUE' ;
00867           X:=VALUE(1,J); PI:=3.141593;
00869           'IF' K<=16 'THEN' 'GOTO' S(/K/) 'ELSE' 'GOTO' SS(/K-16/);
00870           AEXP: VALUE:=EXP(X); 'GOTO' END;
00872           ALN: VALUE:=LN(X); 'GOTO' END;
00874           SINH: VALUE:=EXP1(X)/2; 'GOTO' END;
00876           COSH: VALUE:=EXP2(X)/2; 'GOTO' END;
00878           TANH: VALUE:=EXP1(X)/EXP2(X) ; 'GOTO' END;
00880           COTH: 'IF' EXP1(X)=0 'THEN' 'GOTO' OVERFLO;
00881           VALUE:=EXP2(X)/EXP1(X);
00882           'GOTO' END;
00883           SECH: VALUE:=2/EXP2(X) ; 'GOTO' END;
00885           CSCH: 'IF' EXP1(X)=0 'THEN' 'GOTO' OVERFLO;
00886           VALUE:=2/EXP1(X); 'GOTO' END;

```

```

00888 ARCSIN: 'IF' X=1 'THEN' 'GOTO' OVERFLO;
00889         VALUE:=ARCTAN(X/SQRT(1-X**2)); 'GOTO' END;
00891 ARCCOS: 'IF' X=1 'THEN' 'GOTO' OVERFLO;
00892         VALUE:=PI/2-ARCTAN(X/SQRT(1-X**2)); 'GOTO' END;
00894 AARCTAN: VALUE:=ARCTAN(X); 'GOTO' END;
00896 ARCCOT: VALUE:=PI/2-ARCTAN(X); 'GOTO' END;
00898 ARCSEC: 'IF' X>=0 'THEN' VALUE:=ARCTAN(SQRT(X**2-1))
00898         'ELSE' VALUE:=ARCTAN(SQRT(X**2-1)-PI); 'GOTO' END;
00900 ARCCSC: 'IF' X>=0 'THEN' VALUE:=ARCTAN(1/SQRT(X**2-1))
00900         'ELSE' VALUE:=ARCTAN(1/SQRT(X**2-1))-PI; 'GOTO' END;
00902 ASIN: VALUE:=SIN(X); 'GOTO' END;
00904 ACOS: VALUE:=COS(X); 'GOTO' END;
00906 TAN: 'IF' COS(X)=0 'THEN' 'GOTO' OVERFLO;
00907         VALUE:=SIN(X)/COS(X); 'GOTO' END;
00909 COT: 'IF' SIN(X)=0 'THEN' 'GOTO' OVERFLO;
00910         VALUE:=COS(X)/SIN(X); 'GOTO' END;
00912 SEC: 'IF' COS(X)=0 'THEN' 'GOTO' OVERFLO;
00913         VALUE:=1/COS(X); 'GOTO' END;
00915 CSC: 'IF' SIN(X)=0 'THEN' 'GOTO' OVERFLO;
00916         VALUE:=1/SIN(X); 'GOTO' END;
00918 ARSINH: VALUE:=LN(X+SQRT(X**2+1)); 'GOTO' END;
00920 ARCOSH: VALUE:=LN(X+SQRT(X**2-1)); 'GOTO' END;
00922 ARTANH: 'IF' ABS(X)>=1 'THEN' 'GOTO' OVERFLO;
00923         VALUE:=0.5*LN(ABS((1+X)/(1-X))); 'GOTO' END;
00925 ARCOTH: 'IF' ABS(X)<=1 'THEN' 'GOTO' OVERFLO;
00926         VALUE:=0.5*LN(ABS((1+X)/(X-1))); 'GOTO' END;
00928 ARSECH: VALUE:=LN(1/X+SQRT(1/X**2-1)); 'GOTO' END;
00930 ARCSCH: VALUE:=LN(1/X+SQRT(1/X**2+1));
00931         END: TAO:=LTAO; TMO:=LTMO;
00933         'END' ;
00934 END: 'END' VALUE;
00935 'PROCEDURE' INSOURCE(I); 'VALUE' I; 'INTEGER' I;
00938         'BEGIN'
00938         'IF' NEXTIN=LSS 'THEN'
00938         'BEGIN'
00938         FAULT(20,'FALSE'); 'GOTO' SSFAIL
00939         'END';
00940         NEXTIN:=NEXTIN+1; SOURCE(/NEXTIN/):=I
00941         'END';
00942 'INTEGER' 'PROCEDURE' OUTSOURCE(I); 'VALUE' I; 'INTEGER' I;
00945         OUTSOURCE:=SOURCE(/I/);
00946 'PROCEDURE' INTARG(I); 'VALUE' I; 'INTEGER' I;
00949         'BEGIN'
00949         'IF' NEXTOUT=LTARG 'THEN'
00949         'BEGIN'
00949         FAULT(21,'FALSE'); 'GOTO' TARGFAIL
00950         'END';
00951         NEXTOUT:=NEXTOUT+1; TARG(/NEXTOUT/):=I

```

```

00952         'END';
00953     'INTEGER' 'PROCEDURE' OUTTARG(I); 'VALUE' I; 'INTEGER' I;
00956         OUTTARG:=TARG(/I/);
00957     'PROCEDURE' READSS(TYPE); 'INTEGER' TYPE;
00959         'BEGIN' 'INTEGER' I;
00960             NEXTIN:=0;
00961             TYPE:=READINT; ORDER:=READINT; EVAL:=READINT;
00964             'IF' TYPE=0 'THEN' 'GOTO' DOCEND;
00965             'FOR' I:=READCH 'WHILE' I~TERMIN 'DO' INSOURCE(I);
00966             INSOURCE(TERM);
00967             NEXTIN:=NEXTIN-1;
00968             SYSACT(1,14,3);
00969             OUTST('('NEW SOURCE STATEMENT')');
00970             MI('('TYPE')',TYPE);
00971             MI('('LENGTH')',NEXTIN);
00972             MI('('HIGHEST ORDER OF DIFFERENTIATION ACHIEVED')',ORDER);
00973             MI('('CODE FOR EVALUATION IS 1 ELSE 0, CODE= ')',EVAL);
00974             SYSACT(1,14,2);
00975             OUTST('('SOURCE STATEMENT :')'); SYSACT(1,14,1);
00977             'FOR' I:=1 'STEP' 1 'UNTIL' NEXTIN 'DO' PRINTCH(OUTSOURCE(I))
00977         'END';
00978     'PROCEDURE' TARGPRINT;
00979         'FOR' I:=1 'STEP' 1 'UNTIL' NEXTOUT 'DO' PRINTCH(OUTTARG(I));
00980     READSYNTAX;
00981     READSEMANTICS;
00982     'FOR' I:=1 'STEP' 1 'UNTIL' 10 'DO' SA(/I/):=I-1;
00983     NEWS: REAOSS(PT);
00984     TESTEVAL:='FALSE'; LORDER:=ORDER; LEVAL:=EVAL;
00987     FINDVAL: 'FOR' M:=LORDER'STEP' -1 'UNTIL' 1 'DO'
00987         'BEGIN'
00987             'FOR' I:=1 'STEP' 1 'UNTIL' 26 'DO' A(/I/):=0;
00988             'FOR' I:=1 'STEP' 1 'UNTIL' LENGTH 'DO'
00988                 'BEGIN'
00988                     TYPE(/I/):=CAT(/I/):=FROM(/I/):=TO(/I/):=0;
00989                     'FOR' J:=1 'STEP' 1 'UNTIL' NLINKS 'DO' LINK(/J,I/):=0
00989                 'END';
00990             LINK(/1,0/):=1; RPOINT:=SPOINT:=1; NEXTOUT:=0;
00993     'IF' TESTEVAL 'THEN' 'BEGIN'
00993         TAO:='TRUE'; TMO:='TRUE';
00995         GETVALUES;
00996             'END'
00996         'ELSE' GETVARIABLE;
00997     'IF' SEARCH(PT) 'THEN'
00997         'BEGIN'
00997             MI('('ANALYSED SUCCESSFULLY AS TYPE')',PT); SYSACT(1,14,1);
00999             OUTST('('TABULAR ANALYSIS RECORD :')'); SYSACT(1,14,2);
01001             OUTST(
01001     ((' ENTRY TYPE CAT FROM TO LINKS'))

```



```
01001      ); SYSACT(1,14,2);
01002      'FOR' I:=1 'STEP' 1 'UNTIL' RPOINT-1 'DO'
01003      'BEGIN' OUTI(I,10); OUTI(TYPE(/I/),10);
01005      OUTI(CAT(/I/),10); OUTI(FROM(/I/),10); OUTI(TO(/I/),10);
01008      'FOR' J:=1 'STEP' 1 'UNTIL' NLINKS 'DO'
01008      OUTI(LINK(/J,I/),10);
01009      SYSACT(1,14,1)
01009      'END';
01010      SYSACT(1,14,2); OUTST('(*TARGET STRING PRODUCED *)');
01012      SYSACT(1,14,1);
01013      'IF' TESTEVAL 'THEN' 'BEGIN'
01013      RESULT:=VALUE(1,0);
01014      OUT R{RESULT,'TRUE',15,3};
01015      'END'
01015      'ELSE' 'BEGIN'
01015      COMPILE(1,0,1); SYSACT(1,14,1);
01017      TARGPRINT;
01018      'FOR' I:=1 'STEP' 1 'UNTIL' NEXTOUT 'DO'
01018      SOURCE(/I/):=TARG(/I/);
01019      NEXTIN:= NEXTOUT;
01020      INSOURCE(TERM);
01021      'END'
01021      'END'
01021      'ELSE'
01021      'BEGIN'
01021      MI('(*ABSENT AS TYPE*)',PT); SYSACT(1,14,1);
01023      OUTST('(*NO ATTEMPT AT COMPILATION POSSIBLE*)'); 'GOTO' NEWSS
01024      'END';
01025      'END';
01026      'IF' LEVAL=1 'THEN' 'BEGIN' TESTEVAL:=TRUE;
01027      LORDER:=1;LEVAL:=0; 'GOTO' FINDVAL
01029      'END';
01030      'GOTO' NEWSS;
01031      SSFAIL: MI('(****OVERLENGTH SS OF TYPE*)',PT); SYSACT(1,14,1);
01033      OUTST('(****PORTION OF SS READ *)'); SYSACT(1,14,1);
01035      'FOR' I:=1 'STEP' 1 'UNTIL' NEXTIN 'DO' PRINTCH(OUTSOURCE(I));
01036      SYSACT(1,14,1);
01037      OUTST('(****NO FURTHER OPERATION UPON THIS SS PERFORMED*)');
01038      'FOR' I:=READCH 'WHILE' I~='TERMIN' 'DO'; 'GOTO' NEWSS;
01040      TARGFAIL: SYSACT(1,14,1); OUTST('(****TARGET STRING OVERLENGTH*)');
01042      SYSACT(1,14,1); OUTST('(****TARGET STRING SO FAR PRODUCED*)');
01044      SYSACT(1,14,1); TARGPRINT; SYSACT(1,14,1);
01047      OUTST('(****NO FURTHER OPERATION UPON THIS SS PERFORMED*)');
01048      'GOTO' NEWSS;
01049      OVERFLO: SYSACT(1,14,1);
01050      OUTST('(***DENOMINATOR CONSISTS OF ZERO OR FUNCTION UNSOLABLE ***)');
01051      SYSACT(1,14,1);
01052      OUTST('(****NO FURTHER OPERATION UPON THIS SS PERFORMED ***)');
```

SC SOURCE STATEMENT

SOURCE PROGRAM

PAGE 035

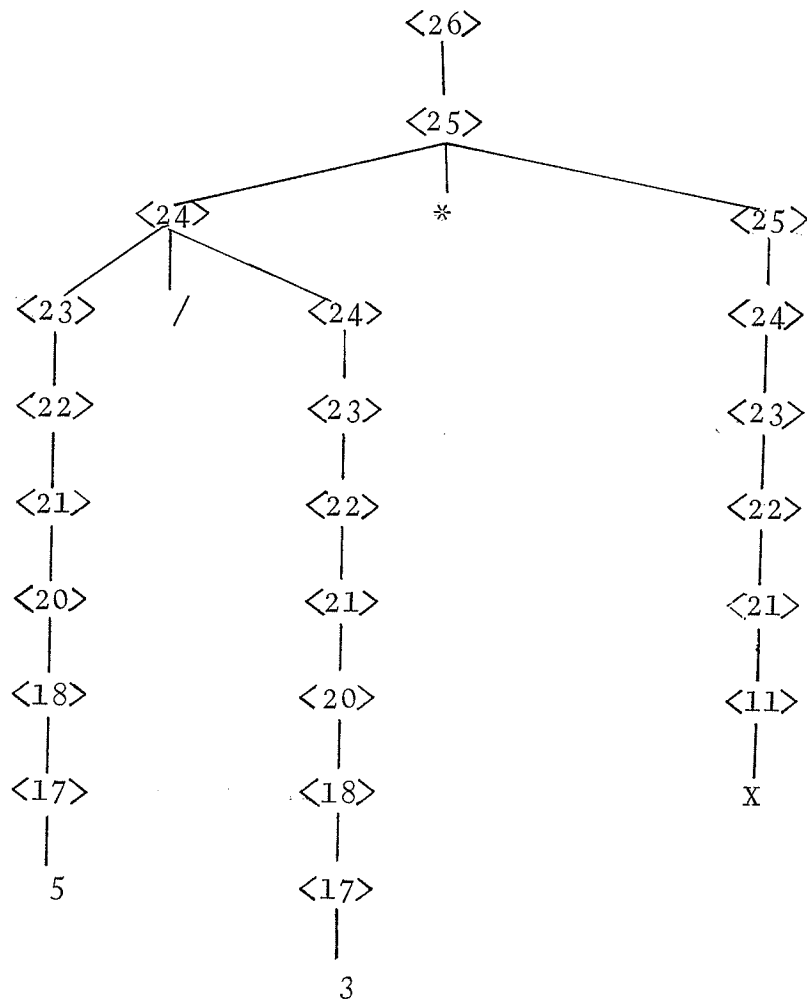
```
01053      *GOTO* NEWSS;
01054  DOCEND: SYSACT(1,14,2); OUTST('*****END OUTPUT THIS DOCUMENT');
01056      *FOR* I:=READCH *WHILE* I->=STAR *DO*
01056      *END*;
01057      *GOTO* NEWDOC;
01058  END: SYSACT(1,14,1);
01059      OUTST('*****END OUTPUT FOR THIS SET OF DOCUMENTS');
01059      *END*
01059      *END* OF ALGOL PROGRAM
```

APPENDIX C

PROBLEM ON PRECEDENCE OF / AND *

As we have mentioned the order of precedence for / is assumed to be higher than that of * in order to achieve correct results.

When using the syntax on page 36, the expression $5/3 * X$ (assumed to be type 26) would have the following syntax tree.



Again assumed the given semantic rules as on page 43, the target string produced would be

$$((0*3-5*0)/3@2*X+5/3*1)$$

On simplification the result is $5/3$ which is correct.

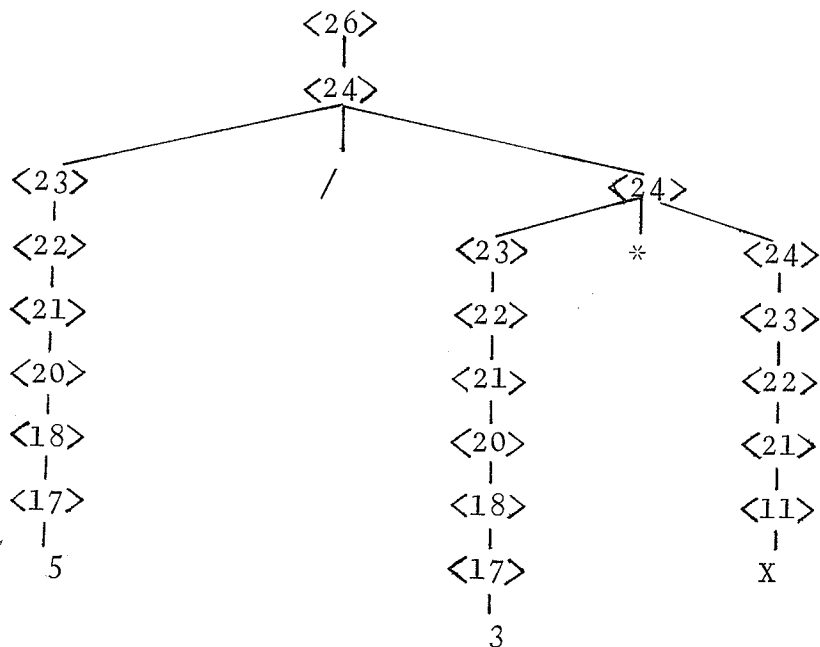
However, if the syntax are specified with * and / of equal precedence, so that the phrase types <term 1> and <term 2> of the syntax table on page 36 are combined as

$$24 \text{ <term> } = \text{ <23> } (* \text{ <24> } \mid / \text{ <24> } \mid \text{ <E> });$$

and phrase type <S.A.E.> is rewritten as

$$26 \text{ <S.A.E.> } = \text{ <24> } (+ \text{ <26> } \mid - \text{ <26> } \mid \text{ <E> });$$

the following syntax tree for $5/3*X$ would be formed.



With the appropriate semantic rules,
the target string produced would be

$$(0*3*X-5*(0*X+3*1))/3*X^2$$

On simplification it is corresponding
to $-5/3X^2$ which is actually the resulting differentiation
from $5/(3*X)$. The reason is that though the operators
/ and * are specified as of equal precedence, the working
order is from right to left in our translator, hence
the product rule is applied first before the quotient
rule, and the differentiation on $5/3*X$ would be as if
it were $5/(3*X)$. In order to avoid this either
parentheses are inserted as $(5/3)*X$ to achieve the
proper order for conventional arithmetic or redefine
the syntax so that / takes higher precedence than *
as in the given syntax on page 36. The latter approach
has been chosen since it has been found that the
resulting differentiation on a string with / and *
operators always gives the equivalent result when the
string is differentiated from left to right and with
equal precedence on * and /.

APPENDIX D

COMPLETE OUTPUT LISTING FOR
DIFFERENTIATIONS OF AN EXPRESSION

The following output listing showing the immediate symbolic and final numeric results of the expressions

$$\frac{\partial^2 x \cos(x-y)}{\partial x \partial y}, \quad \frac{\partial^2 x \cos(x-y)}{\partial y \partial x}$$

The complete input data set for this algebraic differentiation is shown on the Appendix E that follows, while program listing of the translator employed is shown in Appendix B.

MAX LENGTH RECORD	500	
MAX NUMBER LINKS	2	
MAX LENGTH SOURCE STRING		400
MAX LENGTH TARGET STRING		400

DIFFERENTIATION OF ALGEBRIC EXPRESSION;

USER PARAMETERS FOR THIS DOCUMENT
 LENGTH OF ILIFFE VECTORS 28
 LENGTH OF SYNTAX VECTOR 450
 LENGTH OF SEMANTICS VECTOR 1200
 NUMBER OF SEMANTIC SETS 1
 READSYNTAX ENTERED

SYNTAX OK TYPE 14
 SYNTAX OK TYPE 16
 SYNTAX OK TYPE 17
 SYNTAX OK TYPE 18
 SYNTAX OK TYPE 19
 SYNTAX OK TYPE 20
 SYNTAX OK TYPE 21
 SYNTAX OK TYPE 22
 SYNTAX OK TYPE 23
 SYNTAX OK TYPE 24
 SYNTAX OK TYPE 25
 SYNTAX OK TYPE 26
 SYNTAX OK TYPE 27
 SYNTAX OK TYPE 28

SYNTAX ENTERED STORES OCCUPIED 426
 POINTERS FOR SYNTAX

0	0	0	0	0	0	0	0	0	0	0	0	0
0	60	0	113	122	143	153	163	177	204	215	226	
237	248	263	274									
CONTENTS OF SYNTAX												
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
-12	16	-12	17	-12	18	-12	19	-12	20	-12	21	15
-12	22	-12	23	-12	24	-12	25	-12	26	-12	27	21
-12	28	-12	29	-12	30	-12	31	-12	32	-12	33	27
-12	34	-12	35	-12	36	-12	37	-12	38	-12	39	33
-12	40	-12	-13	41	-10	-1	16	-12	-12	-11	-12	39
-13	1	-12	2	-12	3	-12	4	-12	5	-12	6	-12
-12	7	-12	8	-12	9	-12	10	-12	-13	-1	17	6
-10	-1	-12	18	-12	-11	-12	-13	-1	17	-10	-1	17
19	-12	-12	-11	-12	-13	-1	18	43	-1	19	-12	-1
43	-1	19	-12	-1	18	-12	-13	46	-1	26	47	-12
-12	-1	28	-12	-1	20	-12	-1	13	-12	-1	11	47
-12	-1	12	-10	-1	16	-12	-12	-11	-12	-13	-10	11
11	-12	12	-12	-12	-11	-1	21	-12	-13	-1	22	-10
-10	54	-1	23	-12	-12	-11	-12	-13	-1	23	-10	22
14	-1	24	-12	-12	-11	-12	-13	-1	24	-10	13	-10
-1	25	-12	-12	-11	-12	-13	-1	25	-10	11	-1	13
26	-12	12	-1	26	-12	-12	-11	-12	-13	-1	26	-1
-10	49	-1	26	-12	-12	-11	-12	-13	-10	19	26	26
30	-12	26	28	-12	-12	-11	-12	-13	-10	19	38	38
33	22	-12	34	15	28	23	28	22	-12	17	29	29
-12	33	19	17	22	-12	17	33	17	22	-12	15	22
32	17	33	23	28	-12	15	32	17	17	29	33	15
-12	15	32	17	34	15	28	-12	15	32	17	17	33
29	34	-12	15	32	17	33	19	17	-12	15	32	17
17	17	33	17	-12	33	23	28	-12	17	29	33	32
-12	34	15	28	-12	17	29	34	-12	33	19	17	33
-12	17	33	17	-12	15	32	33	23	28	22	-12	17
15	32	17	29	33	22	-12	15	32	34	15	28	-12

22	-12	15	32	17	29	34	22	-12	15	32	33
19	17	22	-12	15	32	17	33	17	22	-12	-11
46	-1	26	47	-12	-13	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

STARTED READING SEMANTIC SET 1

SEMANTICS OK TYPE 11
SEMANTICS OK TYPE 12
SEMANTICS OK TYPE 13
SEMANTICS OK TYPE 16
SEMANTICS OK TYPE 20
SEMANTICS OK TYPE 21
SEMANTICS OK TYPE 22
SEMANTICS OK TYPE 23
SEMANTICS OK TYPE 24
SEMANTICS OK TYPE 25
SEMANTICS OK TYPE 26
SEMANTICS OK TYPE 27
SEMANTICS OK TYPE 28

SEMANTICS ENTERED FOR SET NUMBER 1

STORES USED 1067
POINTERS FOR THIS SET:

0	0	0	0	0	0	0	0	0	0	1	5
10	0	0	14	0	0	0	0	0	0	70	126
170	207	239	253				16	20	53		

CONTENTS OF THIS SEMANTIC SET

37	2	-4	-13	30	37	41	-4	-13	37	1	-4
-13	30	-13	37	1	-4	-13	-5	1	-9	37	46
-4	17	1	1	37	47	-4	-3	-5	2	3	4
5	6	7	-9	17	1	1	-3	-5	6	-9	17
2	1	-3	-13	-5	2	-9	37	12	-4	-3	-5
1	2	3	-9	17	1	1	-3	-13	-5	1	-9
26	1	37	54	-4	26	2	37	13	46	46	-4
26	2	37	13	-4	17	1	1	37	47	14	-4
26	1	37	11	-4	17	2	1	37	13	26	28
46	-4	26	1	37	47	47	-4	-2	-5	2	-9
17	1	1	-3	-13	-5	1	-9	37	46	-4	17
1	1	37	13	-4	26	2	37	12	-4	26	1
37	13	-4	17	2	1	37	47	14	-4	26	2
37	54	3	-4	-2	-5	2	-9	17	1	1	-3
-13	-5	1	-9	37	46	-4	37	1	1	37	13
-4	26	2	37	11	-4	26	1	37	13	-4	17
2	1	37	47	-4	-2	-5	2	-9	17	1	1
-3	-13	-5	1	2	3	-9	17	3	1	-3	-5
1	-9	37	11	-4	-3	-5	2	-9	37	12	-4
-3	-5	1	2	-9	17	2	1	-3	-13	17	1
1	-5	1	-9	37	46	-4	17	2	1	-3	-13
-5	1	-9	37	46	-4	17	1	1	37	47	13
19	38	30	46	-4	26	1	37	47	-4	-2	-5
2	-9	37	46	-4	17	1	1	37	47	14	46
-4	26	1	37	47	-4	-2	-5	3	-9	37	46
-4	17	1	1	37	47	13	17	29	33	22	46
-4	26	1	37	47	-4	-2	-5	4	-9	37	46
-4	17	1	1	37	47	13	33	23	28	22	46
-4	26	1	37	47	-4	-2	-5	5	-9	37	46
-4	17	1	1	37	47	13	46	33	19	17	22
46	-4	26	1	37	47	47	54	3	-4	-2	-5
6	-9	37	46	12	46	47	17	1	37	47	47
47	13	46	17	33	17	22	46	-4	26	1	37

47	47	54	3	-4	-2	-5	7	-9	37	46	12
46	-4	17	1	1	37	47	47	13	33	19	17
22	46	-4	26	1	37	47	13	34	15	28	22
46	-4	26	1	37	47	-4	-2	-5	8	-9	37
46	12	46	-4	17	1	1	37	47	47	13	17
33	17	22	46	-4	26	1	37	47	13	17	29
34	22	46	-4	26	1	37	47	-4	-2	-5	9
-9	37	46	-4	17	1	1	37	47	14	46	2
12	46	-4	26	1	37	47	54	3	47	54	46
2	14	3	47	-4	-2	-5	10	-9	37	46	12
46	-4	17	1	1	37	47	47	14	46	46	2
46	-4	26	1	37	47	54	3	47	54	46	2
14	3	47	-4	-2	-5	11	-9	37	46	-4	17
1	1	37	47	14	46	2	11	46	-4	26	1
37	47	54	3	47	-4	-2	-5	12	-9	37	46
12	46	-4	17	1	1	37	47	47	14	46	2
11	46	-4	26	1	37	47	54	3	47	-4	-2
-5	13	-9	37	46	-4	17	1	1	37	47	14
46	46	-4	26	1	37	47	13	46	46	-4	26
1	37	47	54	3	12	2	47	54	46	2	14
3	47	47	-4	-2	-5	14	-9	37	46	12	46
-4	17	1	17	37	47	47	14	46	46	-4	26
1	37	47	13	46	46	-4	26	1	37	47	54
3	12	2	47	54	46	2	14	3	47	47	-4
-2	-5	15	-9	37	46	-4	17	1	1	37	47
13	17	29	33	46	-4	26	1	37	47	-4	-2
-5	16	-9	37	46	12	46	-4	17	1	1	37
47	47	13	33	23	28	46	-4	26	1	37	47
-4	-2	-5	17	-9	37	46	-4	17	1	1	37
47	13	46	33	19	17	46	-4	26	1	37	47
47	54	3	-4	-2	-5	18	-9	37	46	12	46
-4	17	1	1	37	47	47	13	46	46	17	33
17	46	-4	26	1	37	47	47	54	3	-4	-2
-5	19	-9	37	46	-4	17	1	1	37	47	13
33	19	17	46	-4	26	1	37	47	13	34	15
28	46	-4	26	1	37	47	-4	-2	-5	20	-9
37	46	12	46	-4	17	1	1	37	47	47	13
17	33	17	46	-4	26	1	37	47	13	17	29
34	46	-4	26	1	37	47	-4	-2	-5	21	-9
37	46	-4	17	1	1	37	47	14	46	2	11
46	-4	26	1	37	47	54	3	47	54	46	2
14	3	47	-4	-2	-5	22	-9	37	46	-4	17
1	1	37	47	14	46	46	-4	26	1	37	47
54	3	12	2	47	54	46	2	14	3	47	-4
-2	-5	23	-9	37	46	-4	17	1	1	37	47
14	46	2	12	46	-4	26	1	37	47	54	3
47	-4	-2	-5	24	-9	37	46	-4	17	1	1
37	47	14	46	2	12	46	-4	26	1	37	47
54	3	47	-4	-2	-5	25	-9	37	46	12	46
-4	17	1	1	37	47	47	14	46	46	-4	26
1	37	47	13	46	2	12	46	-4	26	1	37
47	54	3	47	54	46	46	14	3	47	47	-4
-2	-5	26	-9	37	46	12	46	-4	17	1	1
37	47	14	46	46	46	-4	26	1	37	47	13
46	2	11	46	-4	26	1	37	47	54	3	47
54	46	2	14	3	47	47	-4	-3	-13		

READSEMANTICS COMPLETE

TOTAL STORES OCCUPIED

1066

NEW SOURCE STATEMENT

TYPE 27
 LENGTH 12
 HIGHEST ORDER OF DIFFERENTIATION ACHIEVED 2
 CODE FOR EVALUATION IS 1 ELSE 0, CODE= 1

SOURCE STATEMENT :

$W = X * \cos(X - Y)$

INDEPENDENT VARIABLE IS X DEPENDENT VARIABLE IS W
 ANALYSED SUCCESSFULLY AS TYPE 27

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	12	2	9
2	26	3	1	1	3	0
3	25	2	1	1	4	0
4	24	2	1	1	5	0
5	23	2	1	1	6	0
6	22	3	1	1	7	0
7	21	7	1	1	8	9
8	12	1	1	1	9	0
9	26	3	3	12	10	0
10	25	1	3	12	11	16
11	24	2	3	3	12	0
12	23	2	3	3	13	0
13	22	3	3	3	14	0
14	21	5	3	3	15	0
15	11	1	3	3	16	0
16	25	2	5	12	17	0
17	24	2	5	12	18	0
18	23	2	5	12	19	0
19	22	3	5	12	20	0
20	21	2	5	12	21	0
21	28	16	5	12	22	0
22	26	2	9	11	23	29
23	25	2	9	9	24	0
24	24	2	9	9	25	0
25	23	2	9	9	26	0
26	22	3	9	9	27	0
27	21	5	9	9	28	0
28	11	1	9	9	29	0
29	26	3	11	11	30	0
30	25	2	11	11	31	0
31	24	2	11	11	32	0
32	23	2	11	11	33	0
33	22	3	11	11	34	0
34	21	4	11	11	35	0
35	13	23	11	11	36	0

TARGET STRING PRODUCED :

$W' = (1 * \cos(X - Y) + X * (-(1 - 0)) * \sin(X - Y))$

INDEPENDENT VARIABLE IS Y DEPENDENT VARIABLE IS W
 ANALYSED SUCCESSFULLY AS TYPE 27

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	35	2	10
2	26	3	1	2	3	0
3	25	2	1	2	4	0
4	24	2	1	2	5	0
5	23	2	1	2	6	0

6	22	3	1	2	7	0
7	21	6	1	2	8	9
8	12	1	1	1	9	0
9	16	2	2	2	10	0
10	26	3	4	35	11	0
11	25	2	4	35	12	0
12	24	2	4	35	13	0
13	23	2	4	35	14	0
14	22	3	4	35	15	0
15	21	1	4	35	16	0
16	26	1	5	34	17	45
17	25	1	5	14	18	25
18	24	2	5	5	19	0
19	23	2	5	5	20	0
20	22	3	5	5	21	0
21	21	3	5	5	22	0
22	20	3	5	5	23	0
23	18	2	5	5	24	25
24	17	2	5	5	0	0
25	25	2	7	14	26	0
26	24	2	7	14	27	0
27	23	2	7	14	28	0
28	22	3	7	14	29	0
29	21	2	7	14	30	0
30	28	16	7	14	31	0
31	26	2	11	13	32	38
32	25	2	11	11	33	0
33	24	2	11	11	34	0
34	23	2	11	11	35	0
35	22	3	11	11	36	0
36	21	4	11	11	37	0
37	13	23	11	11	38	0
38	26	3	13	13	39	0
39	25	2	13	13	40	0
40	24	2	13	13	41	0
41	23	2	13	13	42	0
42	22	3	13	13	43	0
43	21	5	13	13	44	0
44	11	1	13	13	45	0
45	26	3	16	34	46	0
46	25	1	16	34	47	52
47	24	2	16	16	48	0
48	23	2	16	16	49	0
49	22	3	16	16	50	0
50	21	4	16	16	51	0
51	13	23	16	16	52	0
52	25	1	18	34	53	81
53	24	2	18	25	54	0
54	23	2	18	25	55	0
55	22	3	18	25	56	0
56	21	1	18	25	57	0
57	26	3	19	24	58	0
58	25	2	19	24	59	0
59	24	2	19	24	60	0
60	23	2	19	24	61	0
61	22	2	19	24	62	0
62	21	1	20	24	63	0
63	26	2	21	23	64	72
64	25	2	21	21	65	0
65	24	2	21	21	66	0
66	23	2	21	21	67	0
67	22	3	21	21	68	0
68	21	3	21	21	69	0
69	20	3	21	21	70	0

70	18	2	21	21	71	72
71	17	2	21	21	0	0
72	26	3	23	23	73	0
73	25	2	23	23	74	0
74	24	2	23	23	75	0
75	23	2	23	23	76	0
76	22	3	23	23	77	0
77	21	3	23	23	78	0
78	20	3	23	23	79	0
79	18	2	23	23	80	81
80	17	1	23	23	0	0
81	25	2	27	34	82	0
82	24	2	27	34	83	0
83	23	2	27	34	84	0
84	22	3	27	34	85	0
85	21	2	27	34	86	0
86	28	15	27	34	87	0
87	26	2	31	33	88	94
88	25	2	31	31	89	0
89	24	2	31	31	90	0
90	23	2	31	31	91	0
91	22	3	31	31	92	0
92	21	4	31	31	93	0
93	13	23	31	31	94	0
94	26	3	33	33	95	0
95	25	2	33	33	96	0
96	24	2	33	33	97	0
97	23	2	33	33	98	0
98	22	3	33	33	99	0
99	21	5	33	33	100	0
100	11	1	33	33	101	0

TARGET STRING PRODUCED :

$$W'' = ((0 * \cos(X-Y) + 1 * (-0-1)) * \sin(X-Y)) + (0 * (-1-0)) * \sin(X-Y) + X * ((-0-0)) * \sin(X-Y) + (-1-0) * (0-1) * \cos(X-Y))$$

$$X = 2.000$$

$$Y = 1.000$$

ANALYSED SUCCESSFULLY AS TYPE 27

TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS
1	27	1	1	106	2 11
2	26	3	1	3	3 0
3	25	2	1	3	4 0
4	24	2	1	3	5 0
5	23	2	1	3	6 0
6	22	3	1	3	7 0
7	21	6	1	3	8 9
8	12	1	1	1	9 0
9	16	1	2	3	10 0
10	16	2	3	3	11 0
11	26	3	5	106	12 0
12	25	2	5	106	13 0
13	24	2	5	106	14 0
14	23	2	5	106	15 0
15	22	3	5	106	16 0
16	21	1	5	106	17 0
17	26	1	6	105	18 110
18	25	2	6	37	19 0
19	24	2	6	37	20 0
20	23	2	6	37	21 0
21	22	3	6	37	22 0

22	21	1	6	37	23	0
23	26	1	7	36	24	52
24	25	1	7	16	25	32
25	24	2	7	7	26	0
26	23	2	7	7	27	0
27	22	3	7	7	28	0
28	21	3	7	7	29	0
29	20	3	7	7	30	0
30	18	2	7	7	31	32
31	17	1	7	7	0	0
32	25	2	9	16	33	0
33	24	2	9	16	34	0
34	23	2	9	16	35	0
35	22	3	9	16	36	0
36	21	2	9	16	37	0
37	28	16	9	16	38	0
38	26	2	13	15	39	45
39	25	2	13	13	40	0
40	24	2	13	13	41	0
41	23	2	13	13	42	0
42	22	3	13	13	43	0
43	21	4	13	13	44	0
44	13	23	13	13	45	0
45	26	3	15	15	46	0
46	25	2	15	15	47	0
47	24	2	15	15	48	0
48	23	2	15	15	49	0
49	22	3	15	15	50	0
50	21	5	15	15	51	0
51	11	1	15	15	52	0
52	26	3	18	36	53	0
53	25	1	18	36	54	61
54	24	2	18	18	55	0
55	23	2	18	18	56	0
56	22	3	18	18	57	0
57	21	3	18	18	58	0
58	20	3	18	18	59	0
59	18	2	18	18	60	61
60	17	2	18	18	0	0
61	25	1	20	36	62	90
62	24	2	20	27	63	0
63	23	2	20	27	64	0
64	22	3	20	27	65	0
65	21	1	20	27	66	0
66	26	3	21	26	67	0
67	25	2	21	26	68	0
68	24	2	21	26	69	0
69	23	2	21	26	70	0
70	22	2	21	26	71	0
71	21	1	22	26	72	0
72	26	2	23	25	73	81
73	25	2	23	23	74	0
74	24	2	23	23	75	0
75	23	2	23	23	76	0
76	22	3	23	23	77	0
77	21	3	23	23	78	0
78	20	3	23	23	79	0
79	18	2	23	23	80	81
80	17	1	23	23	0	0
81	26	3	25	25	82	0
82	25	2	25	25	83	0
83	24	2	25	25	84	0
84	23	2	25	25	85	0
85	22	3	25	25	86	0

86	21	3	25	25	87	0
87	20	3	25	25	88	0
88	18	2	25	25	89	90
89	17	2	25	25	90	0
90	25	2	29	36	91	0
91	24	2	29	36	92	0
92	23	2	29	36	93	0
93	22	3	29	36	94	0
94	21	2	29	36	95	0
95	28	15	29	36	96	0
96	26	2	33	35	97	103
97	25	2	33	33	98	0
98	24	2	33	33	99	0
99	23	2	33	33	100	0
100	22	3	33	33	101	0
101	21	4	33	33	102	0
102	13	23	33	33	103	0
103	26	3	35	35	104	0
104	25	2	35	35	105	0
105	24	2	35	35	106	0
106	23	2	35	35	107	0
107	22	3	35	35	108	0
108	21	5	35	35	109	0
109	11	1	35	35	110	0
110	26	3	39	105	111	0
111	25	2	39	105	112	0
112	24	2	39	105	113	0
113	23	2	39	105	114	0
114	22	3	39	105	115	0
115	21	1	39	105	116	0
116	26	1	40	104	117	174
117	25	1	40	58	118	125
118	24	2	40	40	119	0
119	23	2	40	40	120	0
120	22	3	40	40	121	0
121	21	3	40	40	122	0
122	20	3	40	40	123	0
123	18	2	40	40	124	125
124	17	1	40	40	0	0
125	25	1	42	58	126	154
126	24	2	42	49	127	0
127	23	2	42	49	128	0
128	22	3	42	49	129	0
129	21	1	42	49	130	0
130	26	3	43	48	131	0
131	25	2	43	48	132	0
132	24	2	43	48	133	0
133	23	2	43	48	134	0
134	22	2	43	48	135	0
135	21	1	44	48	136	0
136	26	2	45	47	137	145
137	25	2	45	45	138	0
138	24	2	45	45	139	0
139	23	2	45	45	140	0
140	22	3	45	45	141	0
141	21	3	45	45	142	0
142	20	3	45	45	143	0
143	18	2	45	45	144	145
144	17	2	45	45	0	0
145	26	3	47	47	146	0
146	25	2	47	47	147	0
147	24	2	47	47	148	0
148	23	2	47	47	149	0
149	22	3	47	47	150	0

150	21	3	47	47	151	0
151	20	3	47	47	152	0
152	18	2	47	47	153	154
153	17	1	47	47	0	0
154	25	2	51	58	155	0
155	24	2	51	58	156	0
156	23	2	51	58	157	0
157	22	3	51	58	158	0
158	21	2	51	58	159	0
159	28	15	51	58	160	0
160	26	2	55	57	161	167
161	25	2	55	55	162	0
162	24	2	55	55	163	0
163	23	2	55	55	164	0
164	22	3	55	55	165	0
165	21	4	55	55	166	0
166	13	23	55	55	167	0
167	26	3	57	57	168	0
168	25	2	57	57	169	0
169	24	2	57	57	170	0
170	23	2	57	57	171	0
171	22	3	57	57	172	0
172	21	5	57	57	173	0
173	11	1	57	57	174	0
174	26	3	60	104	175	0
175	25	1	60	104	176	181
176	24	2	60	60	177	0
177	23	2	60	60	178	0
178	22	3	60	60	179	0
179	21	4	60	60	180	0
180	13	23	60	60	181	0
181	25	2	62	104	182	0
182	24	2	62	104	183	0
183	23	2	62	104	184	0
184	22	3	62	104	185	0
185	21	1	62	104	186	0
186	26	1	63	103	187	236
187	25	1	63	79	188	216
188	24	2	63	70	189	0
189	23	2	63	70	190	0
190	22	3	63	70	191	0
191	21	1	63	70	192	0
192	26	3	64	69	193	0
193	25	2	64	69	194	0
194	24	2	64	69	195	0
195	23	2	64	69	196	0
196	22	2	64	69	197	0
197	21	1	65	69	198	0
198	26	2	66	68	199	207
199	25	2	66	66	200	0
200	24	2	66	66	201	0
201	23	2	66	66	202	0
202	22	3	66	66	203	0
203	21	3	66	66	204	0
204	20	3	66	66	205	0
205	18	2	66	66	206	207
206	17	1	66	66	0	0
207	26	3	68	68	208	0
208	25	2	68	68	209	0
209	24	2	68	68	210	0
210	23	2	68	68	211	0
211	22	3	68	68	212	0
212	21	3	68	68	213	0
213	20	3	68	68	214	0

214	18	2	68	68	215	216
215	17	1	68	68	0	0
216	25	2	72	79	217	0
217	24	2	72	79	218	0
218	23	2	72	79	219	0
219	22	3	72	79	220	0
220	21	2	72	79	221	0
221	28	15	72	79	222	0
222	26	2	76	78	223	229
223	25	2	76	76	224	0
224	24	2	76	76	225	0
225	23	2	76	76	226	0
226	22	3	76	76	227	0
227	21	4	76	76	228	0
228	13	23	76	76	229	0
229	26	3	78	78	230	0
230	25	2	78	78	231	0
231	24	2	78	78	232	0
232	23	2	78	78	233	0
233	22	3	78	78	234	0
234	21	5	78	78	235	0
235	11	1	78	78	236	0
236	26	3	81	103	237	0
237	25	1	81	103	238	266
238	24	2	81	88	239	0
239	23	2	81	88	240	0
240	22	3	81	88	241	0
241	21	1	81	88	242	0
242	26	3	82	87	243	0
243	25	2	82	87	244	0
244	24	2	82	87	245	0
245	23	2	82	87	246	0
246	22	2	82	87	247	0
247	21	1	83	87	248	0
248	26	2	84	86	249	257
249	25	2	84	84	250	0
250	24	2	84	84	251	0
251	23	2	84	84	252	0
252	22	3	84	84	253	0
253	21	3	84	84	254	0
254	20	3	84	84	255	0
255	18	2	84	84	256	257
256	17	2	84	84	0	0
257	26	3	86	86	258	0
258	25	2	86	86	259	0
259	24	2	86	86	260	0
260	23	2	86	86	261	0
261	22	3	86	86	262	0
262	21	3	86	86	263	0
263	20	3	86	86	264	0
264	18	2	86	86	265	266
265	17	1	86	86	0	0
266	25	1	90	103	267	289
267	24	2	90	94	268	0
268	23	2	90	94	269	0
269	22	3	90	94	270	0
270	21	1	90	94	271	0
271	26	2	91	93	272	280
272	25	2	91	91	273	0
273	24	2	91	91	274	0
274	23	2	91	91	275	0
275	22	3	91	91	276	0
276	21	3	91	91	277	0
277	20	3	91	91	278	0

278	18	2	91	91	279	280
279	17	1	91	91	0	0
280	26	3	93	93	281	0
281	25	2	93	93	282	0
282	24	2	93	93	283	0
283	23	2	93	93	284	0
284	22	3	93	93	285	0
285	21	3	93	93	286	0
286	20	3	93	93	287	0
287	18	2	93	93	288	289
288	17	2	93	93	0	0
289	25	2	96	103	290	0
290	24	2	96	103	291	0
291	23	2	96	103	292	0
292	22	3	96	103	293	0
293	21	2	96	103	294	0
294	28	16	96	103	295	0
295	26	2	100	102	296	302
296	25	2	100	100	297	0
297	24	2	100	100	298	0
298	23	2	100	100	299	0
299	22	3	100	100	300	0
300	21	4	100	100	301	0
301	13	23	100	100	302	0
302	26	3	102	102	303	0
303	25	2	102	102	304	0
304	24	2	102	102	305	0
305	23	2	102	102	306	0
306	22	3	102	102	307	0
307	21	5	102	102	308	0
308	11	1	102	102	309	0

TARGET STRING PRODUCED :
W**= 1.922

NEW SOURCE STATEMENT

TYPE 27
LENGTH 12
HIGHEST ORDER OF DIFFERENTIATION ACHIEVED 2
CODE FOR EVALUATION IS 1 ELSE 0, CODE= 1

SOURCE STATEMENT :

W=X* $\cos(X-Y)$
INDEPENDENT VARIABLE IS Y DEPENDENT VARIABLE IS W
ANALYSED SUCCESSFULLY AS TYPE 27
TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS	
1	27	1	1	12	2	9
2	26	3	1	1	3	0
3	25	2	1	1	4	0
4	24	2	1	1	5	0
5	23	2	1	1	6	0
6	22	3	1	1	7	0
7	21	7	1	1	8	9
8	12	1	1	1	9	0
9	26	3	3	12	10	0
10	25	1	3	12	11	16
11	24	2	3	3	12	0
12	23	2	3	3	13	0
13	22	3	3	3	14	0

14	21	4	3	3	15	0
15	13	23	3	3	16	0
16	25	2	5	12	17	0
17	24	2	5	12	18	0
18	23	2	5	12	19	0
19	22	3	5	12	20	0
20	21	2	5	12	21	0
21	28	16	5	12	22	0
22	26	2	9	11	23	29
23	25	2	9	9	24	0
24	24	2	9	9	25	0
25	23	2	9	9	26	0
26	22	3	9	9	27	0
27	21	4	9	9	28	0
28	13	23	9	9	29	0
29	26	3	11	11	30	0
30	25	2	11	11	31	0
31	24	2	11	11	32	0
32	23	2	11	11	33	0
33	22	3	11	11	34	0
34	21	5	11	11	35	0
35	11	1	11	11	36	0

TARGET STRING PRODUCED :

W*={0*cos(x-y)+x*(-(0-1))*sin(x-y)}
 INDEPENDENT VARIABLE IS X DEPENDENT VARIABLE IS W
 ANALYSED SUCCESSFULLY AS TYPE 27
 TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS
1	27	1	1	35	2
2	26	3	1	2	3
3	25	2	1	2	4
4	24	2	1	2	5
5	23	2	1	2	6
6	22	3	1	2	7
7	21	6	1	2	8
8	12	1	1	1	9
9	16	2	2	2	10
10	26	3	4	35	11
11	25	2	4	35	12
12	24	2	4	35	13
13	23	2	4	35	14
14	22	3	4	35	15
15	21	1	4	35	16
16	26	1	5	34	17
17	25	1	5	14	18
18	24	2	5	5	19
19	23	2	5	5	20
20	22	3	5	5	21
21	21	3	5	5	22
22	20	3	5	5	23
23	18	2	5	5	24
24	17	1	5	5	0
25	25	2	7	14	26
26	24	2	7	14	27
27	23	2	7	14	28
28	22	3	7	14	29
29	21	2	7	14	30
30	28	16	7	14	31
31	26	2	11	13	32

32	25	2	11	11	33	0
33	24	2	11	11	34	0
34	23	2	11	11	35	0
35	22	3	11	11	36	0
36	21	5	11	11	37	0
37	11	1	11	11	38	0
38	26	3	13	13	39	0
39	25	2	13	13	40	0
40	24	2	13	13	41	0
41	23	2	13	13	42	0
42	22	3	13	13	43	0
43	21	4	13	13	44	0
44	13	23	13	13	45	0
45	26	3	16	34	46	0
46	25	1	16	34	47	52
47	24	2	16	16	48	0
48	23	2	16	16	49	0
49	22	3	16	16	50	0
50	21	5	16	16	51	0
51	11	1	16	16	52	0
52	25	1	18	34	53	81
53	24	2	18	25	54	0
54	23	2	18	25	55	0
55	22	3	18	25	56	0
56	21	1	18	25	57	0
57	26	3	19	24	58	0
58	25	2	19	24	59	0
59	24	2	19	24	60	0
60	23	2	19	24	61	0
61	22	2	19	24	62	0
62	21	1	20	24	63	0
63	26	2	21	23	64	72
64	25	2	21	21	65	0
65	24	2	21	21	66	0
66	23	2	21	21	67	0
67	22	3	21	21	68	0
68	21	3	21	21	69	0
69	20	3	21	21	70	0
70	18	2	21	21	71	72
71	17	1	21	21	0	0
72	26	3	23	23	73	0
73	25	2	23	23	74	0
74	24	2	23	23	75	0
75	23	2	23	23	76	0
76	22	3	23	23	77	0
77	21	3	23	23	78	0
78	20	3	23	23	79	0
79	18	2	23	23	80	81
80	17	2	23	23	0	0
81	25	2	27	34	82	0
82	24	2	27	34	83	0
83	23	2	27	34	84	0
84	22	3	27	34	85	0
85	21	2	27	34	86	0
86	28	15	27	34	87	0
87	26	2	31	33	88	94
88	25	2	31	31	89	0
89	24	2	31	31	90	0
90	23	2	31	31	91	0
91	22	3	31	31	92	0
92	21	5	31	31	93	0
93	11	1	31	31	94	0
94	26	3	33	33	95	0
95	25	2	33	33	96	0

96	24	2	33	33	97	0
97	23	2	33	33	98	0
98	22	3	33	33	99	0
99	21	4	33	33	100	0
100	13	23	33	33	101	0

TARGET STRING PRODUCED :

W**=((0*COS(X-Y)+0*(-(1-0))*SIN(X-Y))+1*(-(0-1))*SIN(X-Y)+X*(-(0-0))*SIN(X-Y)+(-(0-1))*(1-0)*COS(X-Y)))
X = 2.000
Y = 1.000

ANALYSED SUCCESSFULLY AS TYPE 27
TABULAR ANALYSIS RECORD :

ENTRY	TYPE	CAT	FROM	TO	LINKS
1	27	1	1	106	2 11
2	26	3	1	3	3 0
3	25	2	1	3	4 0
4	24	2	1	3	5 0
5	23	2	1	3	6 0
6	22	3	1	3	7 0
7	21	6	1	3	8 9
8	12	1	1	1	9 0
9	16	1	2	3	10 0
10	16	2	3	3	11 0
11	26	3	5	106	12 0
12	25	2	5	106	13 0
13	24	2	5	106	14 0
14	23	2	5	106	15 0
15	22	3	5	106	16 0
16	21	1	5	106	17 0
17	26	1	6	105	18 110
18	25	2	6	37	19 0
19	24	2	6	37	20 0
20	23	2	6	37	21 0
21	22	3	6	37	22 0
22	21	1	6	37	23 0
23	26	1	7	36	24 52
24	25	1	7	16	25 32
25	24	2	7	7	26 0
26	23	2	7	7	27 0
27	22	3	7	7	28 0
28	21	3	7	7	29 0
29	20	3	7	7	30 0
30	18	2	7	7	31 32
31	17	1	7	7	0 0
32	25	2	9	16	33 0
33	24	2	9	16	34 0
34	23	2	9	16	35 0
35	22	3	9	16	36 0
36	21	2	9	16	37 0
37	28	16	9	16	38 0
38	26	2	13	15	39 45
39	25	2	13	13	40 0
40	24	2	13	13	41 0
41	23	2	13	13	42 0
42	22	3	13	13	43 0
43	21	5	13	13	44 0
44	11	1	13	13	45 0
45	26	3	15	15	46 0
46	25	2	15	15	47 0
47	24	2	15	15	48 0

48	23	2	15	15	49	0
49	22	3	15	15	50	0
50	21	4	15	15	51	0
51	13	23	15	15	52	0
52	26	3	18	36	53	0
53	25	1	18	36	54	61
54	24	2	18	18	55	0
55	23	2	18	18	56	0
56	22	3	18	18	57	0
57	21	3	18	18	58	0
58	20	3	18	18	59	0
59	18	2	18	18	60	61
60	17	1	18	18	0	0
61	25	1	20	36	62	90
62	24	2	20	27	63	0
63	23	2	20	27	64	0
64	22	3	20	27	65	0
65	21	1	20	27	66	0
66	26	3	21	26	67	0
67	25	2	21	26	68	0
68	24	2	21	26	69	0
69	23	2	21	26	70	0
70	22	2	21	26	71	0
71	21	1	22	26	72	0
72	26	2	23	25	73	81
73	25	2	23	23	74	0
74	24	2	23	23	75	0
75	23	2	23	23	76	0
76	22	3	23	23	77	0
77	21	3	23	23	78	0
78	20	3	23	23	79	0
79	18	2	23	23	80	81
80	17	2	23	23	0	0
81	26	3	25	25	82	0
82	25	2	25	25	83	0
83	24	2	25	25	84	0
84	23	2	25	25	85	0
85	22	3	25	25	86	0
86	21	3	25	25	87	0
87	20	3	25	25	88	0
88	18	2	25	25	89	90
89	17	1	25	25	0	0
90	25	2	29	36	91	0
91	24	2	29	36	92	0
92	23	2	29	36	93	0
93	22	3	29	36	94	0
94	21	2	29	36	95	0
95	28	15	29	36	96	0
96	26	2	33	35	97	103
97	25	2	33	33	98	0
98	24	2	33	33	99	0
99	23	2	33	33	100	0
100	22	3	33	33	101	0
101	21	5	33	33	102	0
102	11	1	33	33	103	0
103	26	3	35	35	104	0
104	25	2	35	35	105	0
105	24	2	35	35	106	0
106	23	2	35	35	107	0
107	22	3	35	35	108	0
108	21	4	35	35	109	0
109	13	23	35	35	110	0
110	26	3	39	105	111	0
111	25	2	39	105	112	0

112	24	2	39	105	113	0
113	23	2	39	105	114	0
114	22	3	39	105	115	0
115	21	1	39	105	116	0
116	26	1	40	104	117	174
117	25	1	40	58	118	125
118	24	2	40	40	119	0
119	23	2	40	40	120	0
120	22	3	40	40	121	0
121	21	3	40	40	122	0
122	20	3	40	40	123	0
123	18	2	40	40	124	125
124	17	2	40	40	0	0
125	25	1	42	58	126	154
126	24	2	42	49	127	0
127	23	2	42	49	128	0
128	22	3	42	49	129	0
129	21	1	42	49	130	0
130	26	3	43	48	131	0
131	25	2	43	48	132	0
132	24	2	43	48	133	0
133	23	2	43	48	134	0
134	22	2	43	48	135	0
135	21	1	44	48	136	0
136	26	2	45	47	137	145
137	25	2	45	45	138	0
138	24	2	45	45	139	0
139	23	2	45	45	140	0
140	22	3	45	45	141	0
141	21	3	45	45	142	0
142	20	3	45	45	143	0
143	18	2	45	45	144	145
144	17	1	45	45	0	0
145	26	3	47	47	146	0
146	25	2	47	47	147	0
147	24	2	47	47	148	0
148	23	2	47	47	149	0
149	22	3	47	47	150	0
150	21	3	47	47	151	0
151	20	3	47	47	152	0
152	18	2	47	47	153	154
153	17	2	47	47	0	0
154	25	2	51	58	155	0
155	24	2	51	58	156	0
156	23	2	51	58	157	0
157	22	3	51	58	158	0
158	21	2	51	58	159	0
159	28	15	51	58	160	0
160	26	2	55	57	161	167
161	25	2	55	55	162	0
162	24	2	55	55	163	0
163	23	2	55	55	164	0
164	22	3	55	55	165	0
165	21	5	55	55	166	0
166	11	1	55	55	167	0
167	26	3	57	57	168	0
168	25	2	57	57	169	0
169	24	2	57	57	170	0
170	23	2	57	57	171	0
171	22	3	57	57	172	0
172	21	4	57	57	173	0
173	13	23	57	57	174	0
174	26	3	60	104	175	0
175	25	1	60	104	176	181

176	24	2	60	60	177	0
177	23	2	60	60	178	0
178	22	3	60	60	179	0
179	21	5	60	60	180	0
180	11	1	60	60	181	0
181	25	2	62	104	182	0
182	24	2	62	104	183	0
183	23	2	62	104	184	0
184	22	3	62	104	185	0
185	21	1	62	104	186	0
186	26	1	63	103	187	236
187	25	1	63	79	188	216
188	24	2	63	70	189	0
189	23	2	63	70	190	0
190	22	3	63	70	191	0
191	21	1	63	70	192	0
192	26	3	64	69	193	0
193	25	2	64	69	194	0
194	24	2	64	69	195	0
195	23	2	64	69	196	0
196	22	2	64	69	197	0
197	21	1	65	69	198	0
198	26	2	66	68	199	207
199	25	2	66	66	200	0
200	24	2	66	66	201	0
201	23	2	66	66	202	0
202	22	3	66	66	203	0
203	21	3	66	66	204	0
204	20	3	66	66	205	0
205	18	2	66	66	206	207
206	17	1	66	66	0	0
207	26	3	68	68	208	0
208	25	2	68	68	209	0
209	24	2	68	68	210	0
210	23	2	68	68	211	0
211	22	3	68	68	212	0
212	21	3	68	68	213	0
213	20	3	68	68	214	0
214	18	2	68	68	215	216
215	17	1	68	68	0	0
216	25	2	72	79	217	0
217	24	2	72	79	218	0
218	23	2	72	79	219	0
219	22	3	72	79	220	0
220	21	2	72	79	221	0
221	28	15	72	79	222	0
222	26	2	76	78	223	229
223	25	2	76	76	224	0
224	24	2	76	76	225	0
225	23	2	76	76	226	0
226	22	3	76	76	227	0
227	21	5	76	76	228	0
228	11	1	76	76	229	0
229	26	3	78	78	230	0
230	25	2	78	78	231	0
231	24	2	78	78	232	0
232	23	2	78	78	233	0
233	22	3	78	78	234	0
234	21	4	78	78	235	0
235	13	23	78	78	236	0
236	26	3	81	103	237	0
237	25	1	81	103	238	266
238	24	2	81	88	239	0
239	23	2	81	88	240	0

240	22	3	81	88	241	0
241	21	1	81	88	242	0
242	26	3	82	87	243	0
243	25	2	82	87	244	0
244	24	2	82	87	245	0
245	23	2	82	87	246	0
246	22	2	82	87	247	0
247	21	1	83	87	248	0
248	26	2	84	86	249	257
249	25	2	84	84	250	0
250	24	2	84	84	251	0
251	23	2	84	84	252	0
252	22	3	84	84	253	0
253	21	3	84	84	254	0
254	20	3	84	84	255	0
255	18	2	84	84	256	257
256	17	1	84	84	0	0
257	26	3	86	86	258	0
258	25	2	86	86	259	0
259	24	2	86	86	260	0
260	23	2	86	86	261	0
261	22	3	86	86	262	0
262	21	3	86	86	263	0
263	20	3	86	86	264	0
264	18	2	86	86	265	266
265	17	2	86	86	0	0
266	25	1	90	103	267	289
267	24	2	90	94	268	0
268	23	2	90	94	269	0
269	22	3	90	94	270	0
270	21	1	90	94	271	0
271	26	2	91	93	272	280
272	25	2	91	91	273	0
273	24	2	91	91	274	0
274	23	2	91	91	275	0
275	22	3	91	91	276	0
276	21	3	91	91	277	0
277	20	3	91	91	278	0
278	18	2	91	91	279	280
279	17	2	91	91	0	0
280	26	3	93	93	281	0
281	25	2	93	93	282	0
282	24	2	93	93	283	0
283	23	2	93	93	284	0
284	22	3	93	93	285	0
285	21	3	93	93	286	0
286	20	3	93	93	287	0
287	18	2	93	93	288	289
288	17	1	93	93	0	0
289	25	2	96	103	290	0
290	24	2	96	103	291	0
291	23	2	96	103	292	0
292	22	3	96	103	293	0
293	21	2	96	103	294	0
294	28	16	96	103	295	0
295	26	2	100	102	296	302
296	25	2	100	100	297	0
297	24	2	100	100	298	0
298	23	2	100	100	299	0
299	22	3	100	100	300	0
300	21	5	100	100	301	0
301	11	1	100	100	302	0
302	26	3	102	102	303	0
303	25	2	102	102	304	0

304	24	2	102	102	305	0
305	23	2	102	102	306	0
306	22	3	102	102	307	0
307	21	4	102	102	308	0
308	13	23	102	102	309	0

TARGET STRING PRODUCED :
W**= 1.922

****END OUTPUT THIS DOCUMENT

END OF RUN:
***END OUTPUT FOR THIS SET OF DOCUMENTS

APPENDIX E
COMPLETE SYNTAX AND SEMANTICS
LISTING FOR DIFFERENTIATION

```

09
<>(|E; &. +=QU, WKLCPSNAZID
500;2;400;400;
+++
DIFFERENTIATION OF ALGEBRIC EXPRESSION;
28,450;1200,1;
14 <ALPHABETS>      =A|B|C|D|E|F|G|H|I|J|K|L|M|N|O|P|Q|R|S|T|U|V|W|X|Y|Z;
16 <QUOTE>          ='(<16>|<E>);
17 <DIGITS>         =0|1|2|3|4|5|6|7|8|9;
18 <INTEGERS>       =<17>(<18>|<E>);
19 <FRACTION>       =<17>(<19>|<E>);
20 <UNSIGNED DECIMAL NUMBER> =<18>.<19>|.<19>|<18>;
21 <PRIMARY>        =<(><26><)>|<28>|<20>|<13>|<11>|<12>(<16>|<E>);
22 <SIGNED PRIMARY>=<+|-|<E>|<21>;
23 <FACTOR>         =<22>(&<23>|<E>);
24 <TERM1>          =<23>(/<24>|<E>);
25 <TERM2>          =<24>(&<25>|<E>);
26 <S.A.E.>         =<25>(&<26>|-<26>|<E>);
27 <STATEMENT>     =<26>(<=><26>|<E>);
28 <FUNCTION>       =(EXP|LN|SINH|COSH|TANH|COth|SECH|CSCH|
                    ARCSIN|ARCCOS|ARCTAN|ARCCOT|ARCSEC|ARCCSC|
                    SIN|COS|TAN|COT|SEC|CSC|
                    ARSINH|ARCOSH|ARTANH|ARCOth|ARSECH|ARCSCH)<(><26><)>;

0;
1;
11 <INDEPENDENT VARIABLE>=W<1>;
12 <DEPENDENT VARIABLE>=PW<'>;
13 <CONSTANT>       =W<0>;
16 <QUOTE>          =P;
20 <UNSIGNED DECIMAL NUMBER>=W<0>;
21 <PRIMARY>        =K<1>W<(>C<1>W<)>.
                    K<2,3,4,5,6,7>C<1>.
                    K<6>C<2>. ;
22 <SIGNED PRIMARY>=K<2>W<->.
                    K<1,2,3>C<1>. ;
23 <FACTOR>         =K<1>L<1>W<@>L<2>W<*>C<1>W<)/L<1>W<+>C<2>
                    W<#LN(>L<1>W<)>)&K<2>C<1>. ;
24 <TERM1>          =K<1>W<(>C<1>W<*>L<2>W<->L<1>W<*>C<2>W<)/L<2>W<@2>&
                    K<2>C<1>. ;
25 <TERM2>          =K<1>W<(>C<1>W<*>L<2>W<+>L<1>W<*>C<2>W<)>&
                    K<2>C<1>. ;
26 <S.A.E.>         =K<1,2,3>C<1>.
                    K<1>W<+>.
                    K<2>W<->.
                    K<1,2>C<2>. ;
27 <STATEMENT>     =C<1>K<1>W<=>C<2>. ;
28 <FUNCTION>       =K<1>W<(>C<1>W<)>*EXP(>L<1>W<)>&
                    K<2>W<(>C<1>W<)/(>L<1>W<)>&
                    K<3>W<(>C<1>W<)*COSH(>L<1>W<)>&
                    K<4>W<(>C<1>W<)*SINH(>L<1>W<)>&
                    K<5>W<(>C<1>W<)*(SECH(>L<1>W<))@2>&
                    K<6>W<(-(>C<1>W<))*(CSCH(>L<1>W<))@2>&
                    K<7>W<(-(>C<1>W<))*SECH(>L<1>W<)*TANH(>L<1>W<)>&
                    K<8>W<(-(>C<1>W<))*CSCH(>L<1>W<)*COth(>L<1>W<)>&
                    K<9>W<(>C<1>W<)/(1-(>L<1>W<))@2)@(1/2)>&
                    K<10>W<(-(>C<1>W<))/(1-(>L<1>W<))@2)@(1/2)>&
                    K<11>W<(>C<1>W<)/(1+(>L<1>W<))@2)>&
                    K<12>W<(-(>C<1>W<))/(1+(>L<1>W<))@2)>&
                    K<13>W<(>C<1>W<)/((>L<1>W<)*((>L<1>W<))@2-1)@(1/2))>&

```

```

      K<14>W<(-(>C<1>W<)) / ((>L<1>W<)*(>L<1>W<)&2-1)&(1/2))>&
      K<15>W<(>C<1>W<)*COS(>L<1>W<)>&
      K<16>W<(-(>C<1>W<))*SIN(>L<1>W<)>&
      K<17>W<(>C<1>W<)*(SEC(>L<1>W<))&2>&
      K<18>W<(-(>C<1>W<))*((CSC(>L<1>W<))&2)>&
      K<19>W<(>C<1>W<)*SEC(>L<1>W<)*TAN(>L<1>W<)>&
      K<20>W<(-(>C<1>W<))*CSC(>L<1>W<)*COT(>L<1>W<)>&
      K<21>W<(>C<1>W<)/(1+(>L<1>W<)&2)&(1/2)>&
      K<22>W<(>C<1>W<)/((>L<1>W<)&2-1)&(1/2)>&
      K<23>W<(>C<1>W<)/(1-(>L<1>W<)&2)>&
      K<24>W<(>C<1>W<)/(1-(>L<1>W<)&2)>&
      K<25>W<(-(>C<1>W<))/((>L<1>W<)*(1-(>L<1>W<)&2)&(1/2))>&
      K<26>W<(-(>C<1>W<))/((>L<1>W<)*(1+(>L<1>W<)&2)&(1/2))>.;&

0;
0;
27;2;1;
W=X*COS(X-Y);
INDEPENDENT VARIABLE IS &X AND DEPENDENT VARIABLE IS &W
INDEPENDENT VARIABLE IS &Y AND DEPENDENT VARIABLE IS &W
&X=2.00; &Y=1.00;#
27;2;1;
W=X*COS(X-Y);
INDEPENDENT VARIABLE IS &Y AND DEPENDENT VARIABLE IS &W
INDEPENDENT VARIABLE IS &X AND DEPENDENT VARIABLE IS &W
&X=2.00; &Y=1.00;#
0;0;0;
+++
+++
END OF RUN;
0;

```

APPENDIX F

LISTINGS OF SIMPLIFIED EXPRESSION

The numeric results of the partial derivatives on Appendix D are done directly on the non-simplified resulting symbolic expressions. If simplified expressions are desired the following two approaches can be used.

1. The resulting non-simplified expressions are passed alternatively with SIM 1 and SIM 2 until the simplest expressions are obtained.
2. After each differentiation, the resulting expressions are simplified first with SIM 1 and SIM 2 before the next differentiation takes place, the results of which are then simplified similarly.

Both approaches are tried in the case of (a) $\frac{\partial^2 x \cos(x-y)}{\partial y \partial x}$ and (b) $\frac{\partial^2 x \cos(x-y)}{\partial x \partial y}$

I. First Approach.

solutions for (a) $\frac{\partial^2 x \cos(x-y)}{\partial y \partial x}$.

- (i) Original expression
 $W = X * \text{COS}(X - Y);$
- (ii) After first differentiation with respect to X.
 $W' = (1 * \text{COS}(X - Y) + X * (-(1 - 0)) * \text{SIN}(X - Y));$

- (iii) After elimination of 0 and 1 (i.e. pass through SIM 1)
- $$W' = (\cos(X-Y) + X * (-\sin(X-Y)));$$
- (iv) After elimination of brackets (i.e. pass through SIM 2)
- $$W' = (\cos(X-Y) + X * (-\sin(X-Y)));$$
- (v) After second differentiation with respect to Y.
- $$W'' = ((-(0-1)) * \sin(X-Y) + (0 * (-\sin(X-Y)) + X * (-(0-1) * \cos(X-Y))));$$
- (vi) After second pass through SIM 1
- $$W'' = (\sin(X-Y) + (X * (-(-\cos(X-Y))));$$
- (vii) After second pass through SIM 2
- $$W'' = (\sin(X-Y) + (X * (-(-\cos(X-Y))));$$
- Corresponding solutions for (b) $\frac{\partial^2 x \cos(x-y)}{\partial x \partial y}$:
- (i) $W = X * \cos(X-Y);$
- (ii) After first differentiation with respect to Y.
- $$W' = (0 * \cos(X-Y) + X * (-(0-1)) * \sin(X-Y));$$
- (iii) $W' = (X * \sin(X-Y));$
- (iv) $W' = (X * \sin(X-Y));$

(v) After second differentiation with respect to X.

$$W'' = (1 * \text{SIN}(X-Y) + X * (1-0) * \text{COS}(X-Y));$$

(vi) $W'' = ((\text{SIN}(X-Y) + X * \text{COS}(X-Y)));$

(vii) $W'' = (\text{SIN}(X-Y) + X * \text{COS}(X-Y));$

II. Second Approach.

solution for (a) $\frac{\partial^2 x \cos(x-y)}{\partial y \partial x} :$

(i) Original expression

$$W = X * \text{COS}(X-Y) ;$$

(ii) After two differentiations

$$W'' = ((0 * \text{COS}(X-Y) + 1 * (- (0-1)) * \text{SIN}(X-Y)) \\ + (0 * (- (1-0)) * \text{SIN}(X-Y) + X * ((- (0-0)) \\ * \text{SIN}(X-Y) + (- (1-0)) * (0-1) * \text{COS}(X-Y)))));$$

(iii) After first pass through SIM 1

$$W'' = ((\text{SIN}(X-Y)) + (X * ((-0) * \text{SIN}(X-Y) \\ + (- (-\text{COS}(X-Y))))));$$

(iv) After first pass through SIM 2

$$W'' = (\text{SIN}(X-Y) + (X * ((-0) * \text{SIN}(X-Y) \\ + (- (-\text{COS}(X-Y))))));$$

(v) After second pass through SIM 1

$$W'' = (\text{SIN}(X-Y) + (X * ((- (-\text{COS}(X-Y))))));$$

(vi) After second pass through SIM 2

$$W'' = (\text{SIN}(X-Y) + (X * (- (-\text{COS}(X-Y)))));$$

Corresponding solutions for (b) $\frac{\partial^2 x \cos(x-y)}{\partial x \partial y}$

- (i) $W = X * \text{COS}(X-Y);$
- (ii) $W'' = ((0 * \text{COS}(X-Y) + 0 * (- (1-0)) * \text{SIN}(X-Y)) + (1 * (- (0-1)) * \text{SIN}(X-Y) + X * ((- (0-0)) * \text{SIN}(X-Y) + (- (0-1)) * (1-0) * \text{COS}(X-Y)))));$
- (iii) $W'' = ((\text{SIN}(X-Y) + X * ((-0) * \text{SIN}(X-Y) + \text{COS}(X-Y)))));$
- (iv) $W'' = (\text{SIN}(X-Y) + X * ((-0) * \text{SIN}(X-Y) + \text{COS}(X-Y)));$
- (v) $W'' = (\text{SIN}(X-Y) + X * \text{COS}(X-Y));$
- (vi) $W'' = (\text{SIN}(X-Y) + X * \text{COS}(X-Y));$

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